

HOL-TestGenFW

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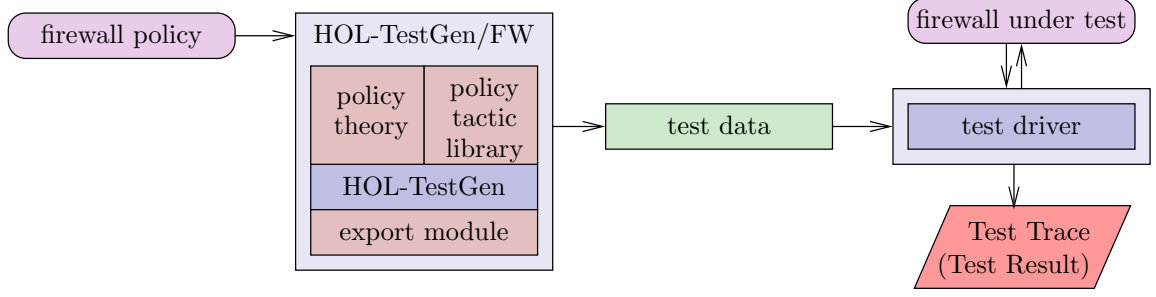


Figure 1: The HOL-TestGen/FW architecture.

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1. Introduction

As HOL-TestGen is built on the framework of Isabelle with a general plug-in mechanism, HOL-TestGen can be customized to implement domain-specific, model-based test tools in its own right. As an example for such a domain-specific test-tool, we developed HOL-TestGen/FW which extends HOL-TestGen by:

1. a theory (or library) formalising networks, protocols and firewall policies,
2. domain-specific extensions of the generic test-case procedures (tactics), and
3. support for an export format of test-data for external tools such as [4].

HOL-TestGen/FW is part of the HOL-TestGen distribution. It is located in the directory `add-ons/security`; see [3, 2] for more details.

Figure 1 shows the overall architecture of HOL-TestGen/FW.

In fact, [item 1](#) defines the formal semantics (in HOL) of a specification language for firewall policies; see [2] and the accompanying examples for details. On the technical level, this library also contains simplification rules together with the corresponding setup of the constraint resolution procedures.

With [item 2](#) we refer to domain-specific processing encapsulated into the general HOL-TestGen test-case generation. Since test specifications in our domain have a specific pattern consisting of a limited set of predicates and policy combinators, this can be exploited in specific pre-processing and post-processing of an optimised version of the procedure, now tuned for stateless firewall policies.

With [item 3](#), we refer to an own XML-like format for exchanging test-data for firewalls, i.e. a description of packets to be send together with the expected behavior of the firewall. This data can be imported in a test-driver for firewalls, for example [\[4\]](#). This completes our toolchain which, thus, supports the execution of test data on firewall implementations based on test cases derived from formal specifications.

2. Installing and using HOL-TestGen/FW

To install HOL-TestGen/FW you need a working installation of HOL-TestGen as described in the HOL-TestGen User Guide. To build the extension, go into the `add-ons/security/src/firewall` directory and build the HOL-TestGen/FW heap image for Isabelle by calling

```
isabelle make
```

Note that this requires that the HOL-UPF heap has been built before. HOL-TestGen/FW can now be started using the `isabelle` command:

```
isabelle jedit -l HOL-TestGenFW
```

or, if HOL-TestGen was built on top of HOLCF instead on HOL only:

```
isabelle jedit -l HOLCF-TestGenFW
```

3. Preliminaries

```
theory
  FWTesting
imports
  PacketFilter/PacketFilter
  NAT/NAT
  FWNormalisation/FWNormalisation
  StatefulFW/StatefulFW
begin
```

```
declare dom-eq-empty-conv [simp del]
```

This is the formalisation in Isabelle/HOL of firewall policies and corresponding networks and packets. It first contains the formalisation of stateless packet filters as described in [\[2\]](#), followed by a verified policy normalisation technique (described in [\[1\]](#)), and a formalisation of stateful protocols described in [\[3\]](#).

The following statement adjusts the pre-normalization step of the default test case generation algorithm. This turns out to be more efficient for the specific case of firewall policies.

Next, the Isar command *prepare-fw-spec* is specified. It can be used to turn test specifications of the form: " $C\ x \Longrightarrow FUT\ x = policy\ x$ " into the desired form for test case generation.

$\langle ML \rangle$

This theory contains a collection of (rather unsorted lemmas) which are of general use when processing test specification (including Normalization) of the Firewall Testing heap.

definition *policyID* **where**

policyID = $(\lambda\ x.\ (Some\ (allow\ (x))))$

lemma *setDistinct*: $\llbracket x \in a; x \notin b \rrbracket \Longrightarrow a \neq b$

$\langle proof \rangle$

lemma *setDistinct3*: $\llbracket x \neq y \rrbracket \Longrightarrow$

$\{\{(a, b, c). a \in x\}\} \neq \{\{(a, b, c). a \in y\}\}$

$\langle proof \rangle$

lemma *setDistinct4*: $\llbracket (a::int) \neq x; a < b; x < y \rrbracket$

$\Longrightarrow \{a..b\} \neq \{x..y\}$

$\langle proof \rangle$

lemma *setDistinct2*: $\llbracket x \neq y \rrbracket \Longrightarrow$

$\{(a, b). a \in x\} \neq \{(a, b). a \in y\}$

$\langle proof \rangle$

lemma *setDistinct1*: $\llbracket x \neq y \rrbracket \Longrightarrow$

$\{(a, b, c). a \in x\} \neq \{(a, b, c). a \in y\}$

$\langle proof \rangle$

lemma *AllowNMT*: $\llbracket x \neq \{\}; y \neq \{\} \rrbracket \Longrightarrow$

$dom\ (Cp\ (AllowPortFromTo\ \{\{(a, b, c). a \in x\}\}\ \{\{(a, b, c). a \in y\}\}\ (p,i))) \neq \{\}$

$\langle proof \rangle$

lemma *domDA1Cp*: $dom\ (Cp\ (DenyAll \oplus P)) = UNIV$

$\langle proof \rangle$

lemma *aux*: $\llbracket x \neq a; y \neq b; (x \neq y \wedge x \neq b) \vee (a \neq b \wedge a \neq y) \rrbracket \Longrightarrow \{x, a\} \neq \{y, b\}$

$\langle proof \rangle$

lemma *aux2*: $\{a,b\} = \{b,a\}$
 $\langle proof \rangle$

lemma *deny-comm[simp]*: $(deny() = X) = (X = deny())$
 $\langle proof \rangle$

lemma *allow-comm[simp]*: $(allow() = X) = (X = allow())$
 $\langle proof \rangle$

lemma *PO-add*: $PO\ P \implies P$
 $\langle proof \rangle$

lemma *dom-2-subset1*: $dom\ (r\ o\text{-}f\ ((A \otimes_2 B)\ o\ (\lambda x. (x,x)))) \subseteq dom\ A$
 $\langle proof \rangle$

lemma *dom-2-comm*: $dom\ (A \otimes_2 B) = dom\ (B \otimes_2 A)$
 $\langle proof \rangle$

lemma *dom-2-subset2*: $dom\ (r\ o\text{-}f\ ((A \otimes_2 B)\ o\ (\lambda x. (x,x)))) \subseteq dom\ B$
 $\langle proof \rangle$

lemma *dom-2*: $dom\ A = UNIV \implies$
 $dom\ (r\ o\text{-}f\ ((A \otimes_2 B)\ o\ (\lambda x. (x,x)))) = dom\ B$
 $\langle proof \rangle$

lemma *domDA1*: $dom\ (C\ (DenyAll \oplus P)) = UNIV$
 $\langle proof \rangle$

lemma *if-DA-simp*: $((if\ P\ then\ (deny\text{-}all\ x)\ else\ None) =$
 $None) = (\neg P)$
 $\langle proof \rangle$

lemma *if-DA-simp2*: $(\exists y. (if\ P\ then\ deny\text{-}all\ x\ else\ None) = Some\ y) = P$
 $\langle proof \rangle$

lemma *if-DA-simp3*: $((if\ P\ then\ deny\text{-}all\ x\ else\ None) =$
 $Some\ ae) = (P \wedge deny\text{-}all\ x = Some\ ae)$
 $\langle proof \rangle$

lemma *if-DA-simp4*: $((if\ P$

```

      then Some (allow x)
      else None) =
None) = ( $\neg$  P)
<proof>

```

```

lemma if-DA-simp5: ((if P
  then Some (deny x)
  else None) =
None) = ( $\neg$  P)
<proof>

```

```

lemma setDistinct6:  $\llbracket x \neq y \rrbracket \implies$ 
   $\{(a, b, c). a = x\} \neq \{(a, b, c). a = y\}$ 
<proof>

```

```

definition ip2int :: (int  $\times$  int  $\times$  int  $\times$  int)  $\Rightarrow$  int where
  ip2int x = (((fst x)*16777216) + ((fst (snd x))*65536) +
    ((fst (snd (snd x)))*256) + snd (snd (snd x)))

```

```

lemmas if-DA = if-DA-simp if-DA-simp2 if-DA-simp3 if-DA-simp4 if-DA-simp5

```

```

declare if-DA [simp]

```

```

end

```

4. Packets and Networks

```

theory NetworkCore
imports Main UPF
begin

```

In networks based e.g. on TCP/IP, a message from A to B is encapsulated in *packets*, which contain the content of the message and routing information. The routing information mainly contains its source and its destination address.

In the case of stateless packet filters, a firewall bases its decision upon this routing information and, in the stateful case, on the content. Thus, we model a packet as a four-tuple of the mentioned elements, together with an id field.

The ID is an integer:

```

type-synonym id = int

```

To enable different representations of addresses (e.g. IPv4 and IPv6, with or without

ports), we model them as an unconstrained type class and directly provide several instances:

```
class adr
```

```
type-synonym   ' $\alpha$  src = ' $\alpha$ 
```

```
type-synonym   ' $\alpha$  dest = ' $\alpha$ 
```

```
instance int :: adr <proof>
```

```
instance nat :: adr <proof>
```

```
instance fun :: (adr,adr) adr <proof>
```

```
instance prod :: (adr,adr) adr <proof>
```

The content is also specified with an unconstrained generic type:

```
type-synonym ' $\beta$  content = ' $\beta$ 
```

For applications where the concrete representation of the content field does not matter (usually the case for stateless packet filters), we provide a default type which can be used in those cases:

```
datatype DummyContent = data
```

Finally, a packet is:

```
type-synonym (' $\alpha$ , ' $\beta$ ) packet = id  $\times$  ' $\alpha$  src  $\times$  ' $\alpha$  dest  $\times$  ' $\beta$  content
```

Please note that protocols (e.g. http) are not modelled explicitly. In the case of stateless packet filters, they are only visible by the destination port of a packet, which are modelled as part of the address. Additionally, stateful firewalls often determine the protocol by the content of a packet.

```
definition src :: (' $\alpha$ ::adr, ' $\beta$ ) packet  $\Rightarrow$  ' $\alpha$ 
```

```
where src = fst  $\circ$  snd
```

Port numbers (which are part of an address) are also modelled in a generic way. The integers and the naturals are typical representations of port numbers.

```
class port
```

```
instance int :: port <proof>
```

```
instance nat :: port <proof>
```

```
instance fun :: (port,port) port <proof>
```

```
instance prod :: (port,port) port <proof>
```

A packet therefore has two parameters, the first being the address, the second the content. For the sake of simplicity, we do not allow to have a different address representation format for the source and the destination of a packet.

In order to access the different parts of a packet directly, we define a couple of projectors:

definition $id :: ('\alpha::adr, '\beta) packet \Rightarrow id$
where $id = fst$

definition $dest :: ('\alpha::adr, '\beta) packet \Rightarrow '\alpha dest$
where $dest = fst \circ snd \circ snd$

definition $content :: ('\alpha::adr, '\beta) packet \Rightarrow '\beta content$
where $content = snd \circ snd \circ snd$

datatype $protocol = tcp \mid udp$

lemma $either: \llbracket a \neq tcp; a \neq udp \rrbracket \Longrightarrow False$
 $\langle proof \rangle$

lemma $either2[simp]: (a \neq tcp) = (a = udp)$
 $\langle proof \rangle$

lemma $either3[simp]: (a \neq udp) = (a = tcp)$
 $\langle proof \rangle$

The following two constants give the source and destination port number of a packet. Address representations using port numbers need to provide a definition for these types.

consts $src-port :: ('\alpha::adr, '\beta) packet \Rightarrow '\gamma::port$
consts $dest-port :: ('\alpha::adr, '\beta) packet \Rightarrow '\gamma::port$
consts $src-protocol :: ('\alpha::adr, '\beta) packet \Rightarrow protocol$
consts $dest-protocol :: ('\alpha::adr, '\beta) packet \Rightarrow protocol$

A subnetwork (or simply a network) is a set of sets of addresses.

type-synonym $'\alpha net = '\alpha set set$

The relation $in_subnet (\sqsubset)$ checks if an address is in a specific network.

definition
 $in-subnet :: '\alpha::adr \Rightarrow '\alpha net \Rightarrow bool$ (**infixl** $\sqsubset 100$) **where**
 $in-subnet a S = (\exists s \in S. a \in s)$

The following lemmas will be useful later.

lemma $in-subnet:$
 $(a, e) \sqsubset \{\{(x1, y). P x1 y\}\} = P a e$
 $\langle proof \rangle$

lemma $src-in-subnet:$
 $src(q, (a, e), r, t) \sqsubset \{\{(x1, y). P x1 y\}\} = P a e$
 $\langle proof \rangle$

lemma *dest-in-subnet*:
 $\text{dest } (q, r, ((a), e), t) \sqsubset \{\{(x1, y). P \ x1 \ y\}\} = P \ a \ e$
<proof>

Address models should provide a definition for the following constant, returning a network consisting of the input address only.

consts *subnet-of* :: $'\alpha::\text{adr} \Rightarrow '\alpha \ \text{net}$

lemmas *packet-defs* = *in-subnet-def id-def content-def src-def dest-def*

end

5. Address Representations

theory
NetworkModels
imports

DatatypeAddress
DatatypePort

IntegerAddress
IntegerPort
IntegerPort-TCPUDP

IPv4
IPv4-TCPUDP

begin

One can think of many different possible address representations. In this distribution, we include seven different variants:

- *DatatypeAddress*: Three explicitly named addresses, which build up a network consisting of three disjunct subnetworks. I.e. there are no overlaps and there is no way to distinguish between individual hosts within a network.
- *DatatypePort*: An address is a pair, with the first element being the same as above, and the second being a port number modelled as an Integer¹.
- *adr_i*: An address in an Integer.
- *adr_ip*: An address is a pair of an Integer and a port (which is again an Integer).

¹For technical reasons, we always use Integers instead of Naturals. As a consequence, the test specifications have to be adjusted to eliminate negative numbers.

- `adr_ipp`: An address is a triple consisting of two Integers modelling the IP address and the port number, and the specification of the network protocol
- `IPv4`: An address is a pair. The first element is a four-tuple of Integers, modelling an IPv4 address, the second element is an Integer denoting the port number.
- `IPv4_TCPUDP`: The same as above, but including additionally the specification of the network protocol.

The theories of each of the networks are relatively small. It suffices to provide the required types, a couple of lemmas, and - if required - a definition for the source and destination ports of a packet.

end

5.1. Datatype Addresses

```
theory DatatypeAddress
imports NetworkCore
begin
```

A theory describing a network consisting of three subnetworks. Hosts within a network are not distinguished.

```
datatype DatatypeAddress = dmz-adr | intranet-adr | internet-adr
```

definition

```
dmz::DatatypeAddress net where
dmz = {{dmz-adr}}
```

definition

```
intranet::DatatypeAddress net where
intranet = {{intranet-adr}}
```

definition

```
internet::DatatypeAddress net where
internet = {{internet-adr}}
```

end

5.2. Datatype Addresses with Ports

```
theory DatatypePort
imports NetworkCore
begin
```

A theory describing a network consisting of three subnetworks, including port numbers modelled as Integers. Hosts within a network are not distinguished.

datatype *DatatypeAddress* = *dmz-adr* | *intranet-adr* | *internet-adr*

type-synonym

port = *int*

type-synonym

DatatypePort = (*DatatypeAddress* × *port*)

instance *DatatypeAddress* :: *adr* *<proof>*

definition

dmz::*DatatypePort net* **where**

dmz = { {(a,b). a = *dmz-adr*}}

definition

intranet::*DatatypePort net* **where**

intranet = { {(a,b). a = *intranet-adr*}}

definition

internet::*DatatypePort net* **where**

internet = { {(a,b). a = *internet-adr*}}

defs (overloaded)

src-port-def: *src-port* (*x*::(*DatatypePort*,*'β*) *packet*) ≡ (*snd o fst o snd*) *x*

dest-port-def: *dest-port* (*x*::(*DatatypePort*,*'β*) *packet*) ≡ (*snd o fst o snd o snd*) *x*

subnet-of-def: *subnet-of* (*x*::*DatatypePort*) ≡ { {(a,b). a = *fst x*}}

lemma *src-port* : *src-port* ((a,x,d,e)::(*DatatypePort*,*'β*) *packet*) = *snd x*
<proof>

lemma *dest-port* : *dest-port* ((a,d,x,e)::(*DatatypePort*,*'β*) *packet*) = *snd x*
<proof>

lemmas *DatatypePortLemmas* = *src-port dest-port src-port-def dest-port-def*

end

5.3. Integer Addresses

theory *IntegerAddress*

imports *NetworkCore*

begin

A theory where addresses are modelled as Integers.

type-synonym

$adr_i = int$

end

5.4. Integer Addresses with Ports

```
theory IntegerPort
imports NetworkCore
begin
```

A theory describing addresses which are modelled as a pair of Integers - the first being the host address, the second the port number.

```
type-synonym
address = int
```

```
type-synonym
port = int
```

```
type-synonym
 $adr_{ip} = address \times port$ 
```

```
defs (overloaded)
```

```
src-port-def:  $src\_port\ (x::(adr_{ip},'\beta)\ packet) \equiv (snd\ o\ fst\ o\ snd)\ x$ 
dest-port-def:  $dest\_port\ (x::(adr_{ip},'\beta)\ packet) \equiv (snd\ o\ fst\ o\ snd\ o\ snd)\ x$ 
subnet-of-def:  $subnet\_of\ (x::(adr_{ip})) \equiv \{\{(a,b). a = fst\ x\}\}$ 
```

```
lemma src-port:  $src\_port\ (a,x::adr_{ip},d,e) = snd\ x$ 
<proof>
```

```
lemma dest-port:  $dest\_port\ (a,d,x::adr_{ip},e) = snd\ x$ 
<proof>
```

```
lemmas  $adr_{ip} Lemmas = src\_port\ dest\_port\ src\_port\_def\ dest\_port\_def$ 
```

end

5.5. Integer Addresses with Ports and Protocols

```
theory IntegerPort-TCPUDP
imports NetworkCore
```

begin

A theory describing addresses which are modelled as a pair of Integers - the first being the host address, the second the port number.

type-synonym

address = int

type-synonym

port = int

type-synonym

adr_{ipp} = address × port × protocol

instance *protocol* :: *adr* <proof>

defs (overloaded)

src-port-def: *src-port* (*x*::(*adr_{ipp}*,*'β*) *packet*) ≡ (*fst o snd o fst o snd*) *x*

dest-port-def: *dest-port* (*x*::(*adr_{ipp}*,*'β*) *packet*) ≡ (*fst o snd o fst o snd o snd*) *x*

subnet-of-def: *subnet-of* (*x*::(*adr_{ipp}*)) ≡ { {(*a,b,c*). *a* = *fst x*} }

src-protocol-def: *src-protocol* (*x*::(*adr_{ipp}*,*'β*) *packet*) ≡ (*snd o snd o fst o snd*) *x*

dest-protocol-def: *dest-protocol* (*x*::(*adr_{ipp}*,*'β*) *packet*) ≡ (*snd o snd o fst o snd o snd*) *x*

lemma *src-port*: *src-port* (*a,x*::*adr_{ipp}*,*d,e*) = *fst* (*snd x*)

<proof>

lemma *dest-port*: *dest-port* (*a,d,x*::*adr_{ipp}*,*e*) = *fst* (*snd x*)

<proof>

Common test constraints:

definition *port-positive* :: (*adr_{ipp}*,*'b*) *packet* ⇒ *bool* **where**

port-positive x = (*dest-port x* > (*0*::*port*))

definition *fix-values* :: (*adr_{ipp}*,*DummyContent*) *packet* ⇒ *bool* **where**

fix-values x = (*src-port x* = (*1*::*port*) ∧ *src-protocol x* = *udp* ∧ *content x* = *data* ∧ *id x* = *1*)

lemmas *adr_{ipp}Lemmas* = *src-port dest-port src-port-def dest-port-def src-protocol-def dest-protocol-def subnet-of-def*

lemmas *adr_{ipp}TestConstraints* = *port-positive-def fix-values-def*

end

5.6. IPv4 Addresses

```
theory IPv4
imports NetworkCore
begin
```

A theory describing IPv4 addresses with ports. The host address is a four-tuple of Integers, the port number is a single Integer.

```
type-synonym
  ipv4-ip = (int × int × int × int)
```

```
type-synonym
  port    = int
```

```
type-synonym
  ipv4     = (ipv4-ip × port)
```

```
defs (overloaded)
  src-port-def: src-port (x::(ipv4,'β) packet) ≡ (snd o fst o snd) x
defs (overloaded)
  dest-port-def: dest-port (x::(ipv4,'β) packet) ≡ (snd o fst o snd o snd) x
defs (overloaded)
  subnet-of-def: subnet-of (x::ipv4) ≡ { {(a,b). a = fst x} }
```

```
definition subnet-of-ip :: ipv4-ip ⇒ ipv4 net
where subnet-of-ip ip = { {(a,b). (a = ip)} }
```

```
lemma src-port: src-port (a,(x::ipv4),d,e) = snd x
  ⟨proof⟩
```

```
lemma dest-port: dest-port (a,d,(x::ipv4),e) = snd x
  ⟨proof⟩
```

```
lemmas IPv4Lemmas = src-port dest-port src-port-def dest-port-def
```

```
end
```

6. Policies

6.1. Policy Core

```

theory PolicyCore
imports NetworkCore UPF
begin

```

A policy is seen as a partial mapping from packet to packet out.

```

type-synonym ('α, 'β) FWPolicy = ('α, 'β) packet  $\mapsto$  unit

```

When combining several rules, the firewall is supposed to apply the first matching one. In our setting this means the first rule which maps the packet in question to *Some* (*packet out*). This is exactly what happens when using the map-add operator (*rule1* ++ *rule2*). The only difference is that the rules must be given in reverse order.

The constant *p-accept* is *True* iff the policy accepts the packet.

definition

```

  p-accept :: ('α, 'β) packet  $\Rightarrow$  ('α, 'β) FWPolicy  $\Rightarrow$  bool where
  p-accept p pol = (pol p =  $\lfloor$ allow () $\rfloor$ )

```

```

end

```

6.2. Policy Combinators

```

theory PolicyCombinators
imports
  PolicyCore
begin

```

In order to ease the specification of a concrete policy, we define some combinators. Using these combinators, the specification of a policy gets very easy, and can be done similarly as in tools like IPTables.

definition

```

  allow-all-from :: 'α::adr net  $\Rightarrow$  (('α, 'β) packet  $\mapsto$  unit) where
  allow-all-from src-net = {pa. src pa  $\sqsubset$  src-net}  $\triangleleft$  AU

```

definition

```

  deny-all-from :: 'α::adr net  $\Rightarrow$  (('α, 'β) packet  $\mapsto$  unit) where
  deny-all-from src-net = {pa. src pa  $\sqsubset$  src-net}  $\triangleleft$  DU

```

definition

```

  allow-all-to :: 'α::adr net  $\Rightarrow$  (('α, 'β) packet  $\mapsto$  unit) where
  allow-all-to dest-net = {pa. dest pa  $\sqsubset$  dest-net}  $\triangleleft$  AU

```

definition

deny-all-to :: $'\alpha::\text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**
deny-all-to dest-net = $\{pa. \text{dest pa} \sqsubset \text{dest-net}\} \triangleleft D_U$

definition

allow-all-from-to :: $'\alpha::\text{adr net} \Rightarrow ' \alpha::\text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**
allow-all-from-to src-net dest-net =
 $\{pa. \text{src pa} \sqsubset \text{src-net} \wedge \text{dest pa} \sqsubset \text{dest-net}\} \triangleleft A_U$

definition

deny-all-from-to :: $'\alpha::\text{adr net} \Rightarrow ' \alpha::\text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**
deny-all-from-to src-net dest-net =
 $\{pa. \text{src pa} \sqsubset \text{src-net} \wedge \text{dest pa} \sqsubset \text{dest-net}\} \triangleleft D_U$

All these combinators and the default rules are put into one single lemma called *PolicyCombinators* to facilitate proving over policies.

lemmas *PolicyCombinators* = *allow-all-from-def deny-all-from-def*
allow-all-to-def deny-all-to-def allow-all-from-to-def
deny-all-from-to-def UPFDefs

end

6.3. Policy Combinators with Ports

theory *PortCombinators*
imports *PolicyCombinators*
begin

This theory defines policy combinators for those network models which have ports. They are provided in addition to the the ones defined in the *PolicyCombinators* theory.

This theory requires from the network models a definition for the two following constants:

- *src_port* :: $(' \alpha, ' \beta) \text{ packet} \Rightarrow (' \gamma :: \text{port})$
- *dest_port* :: $(' \alpha, ' \beta) \text{ packet} \Rightarrow (' \gamma :: \text{port})$

definition

allow-all-from-port :: $'\alpha::\text{adr net} \Rightarrow (' \gamma::\text{port}) \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**
allow-all-from-port src-net s-port =
 $\{pa. \text{src-port pa} = \text{s-port}\} \triangleleft \text{allow-all-from src-net}$

definition

deny-all-from-port :: $'\alpha::\text{adr net} \Rightarrow ' \gamma::\text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**
deny-all-from-port src-net s-port =
 $\{pa. \text{src-port pa} = \text{s-port}\} \triangleleft \text{deny-all-from src-net}$

definition

$allow-all-to-port :: 'α::adr\ net \Rightarrow 'γ::port \Rightarrow (('α, 'β)\ packet \mapsto unit)$ **where**
 $allow-all-to-port\ dest-net\ d-port =$
 $\{pa.\ dest-port\ pa = d-port\} \triangleleft allow-all-to\ dest-net$

definition

$deny-all-to-port :: 'α::adr\ net \Rightarrow 'γ::port \Rightarrow (('α, 'β)\ packet \mapsto unit)$ **where**
 $deny-all-to-port\ dest-net\ d-port =$
 $\{pa.\ dest-port\ pa = d-port\} \triangleleft deny-all-to\ dest-net$

definition

$allow-all-from-port-to :: 'α::adr\ net \Rightarrow 'γ::port \Rightarrow 'α::adr\ net \Rightarrow (('α, 'β)\ packet \mapsto unit)$
where
 $allow-all-from-port-to\ src-net\ s-port\ dest-net$
 $= \{pa.\ src-port\ pa = s-port\} \triangleleft allow-all-from-to\ src-net\ dest-net$

definition

$deny-all-from-port-to :: 'α::adr\ net \Rightarrow 'γ::port \Rightarrow 'α::adr\ net \Rightarrow (('α, 'β)\ packet \mapsto unit)$
where
 $deny-all-from-port-to\ src-net\ s-port\ dest-net$
 $= \{pa.\ src-port\ pa = s-port\} \triangleleft deny-all-from-to\ src-net\ dest-net$

definition

$allow-all-from-port-to-port :: 'α::adr\ net \Rightarrow 'γ::port \Rightarrow 'α::adr\ net \Rightarrow 'γ::port \Rightarrow$
 $(('α, 'β)\ packet \mapsto unit)$ **where**
 $allow-all-from-port-to-port\ src-net\ s-port\ dest-net\ d-port =$
 $\{pa.\ dest-port\ pa = d-port\} \triangleleft allow-all-from-port-to\ src-net\ s-port\ dest-net$

definition

$deny-all-from-port-to-port :: 'α::adr\ net \Rightarrow 'γ::port \Rightarrow 'α::adr\ net \Rightarrow$
 $'γ::port \Rightarrow (('α, 'β)\ packet \mapsto unit)$ **where**
 $deny-all-from-port-to-port\ src-net\ s-port\ dest-net\ d-port =$
 $\{pa.\ dest-port\ pa = d-port\} \triangleleft deny-all-from-port-to\ src-net\ s-port\ dest-net$

definition

$allow-all-from-to-port :: 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow$
 $'γ::port \Rightarrow (('α, 'β)\ packet \mapsto unit)$ **where**
 $allow-all-from-to-port\ src-net\ dest-net\ d-port =$
 $\{pa.\ dest-port\ pa = d-port\} \triangleleft allow-all-from-to\ src-net\ dest-net$

definition

$deny-all-from-to-port :: 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow 'γ::port \Rightarrow$
 $(('α, 'β)\ packet \mapsto unit)$ **where**
 $deny-all-from-to-port\ src-net\ dest-net\ d-port = \{pa.\ dest-port\ pa = d-port\} \triangleleft deny-all-from-to$
 $src-net\ dest-net$

definition

$allow-from-port-to :: 'γ::port \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow (('α, 'β)\ packet \mapsto unit)$
where

allow-from-port-to port src-net dest-net =
 $\{pa. \text{src-port } pa = \text{port}\} \triangleleft \text{allow-all-from-to src-net dest-net}$

definition

deny-from-port-to :: $'\gamma::\text{port} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$

where

deny-from-port-to port src-net dest-net =
 $\{pa. \text{src-port } pa = \text{port}\} \triangleleft \text{deny-all-from-to src-net dest-net}$

definition

allow-from-to-port :: $'\gamma::\text{port} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$

where

allow-from-to-port port src-net dest-net =
 $\{pa. \text{dest-port } pa = \text{port}\} \triangleleft \text{allow-all-from-to src-net dest-net}$

definition

deny-from-to-port :: $'\gamma::\text{port} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$

where

deny-from-to-port port src-net dest-net =
 $\{pa. \text{dest-port } pa = \text{port}\} \triangleleft \text{deny-all-from-to src-net dest-net}$

definition

allow-from-ports-to :: $'\gamma::\text{port set} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow$
 $((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**

allow-from-ports-to ports src-net dest-net =
 $\{pa. \text{src-port } pa \in \text{ports}\} \triangleleft \text{allow-all-from-to src-net dest-net}$

definition

allow-from-to-ports :: $'\gamma::\text{port set} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow$
 $((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**

allow-from-to-ports ports src-net dest-net =
 $\{pa. \text{dest-port } pa \in \text{ports}\} \triangleleft \text{allow-all-from-to src-net dest-net}$

definition

deny-from-ports-to :: $'\gamma::\text{port set} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow$
 $((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**

deny-from-ports-to ports src-net dest-net =
 $\{pa. \text{src-port } pa \in \text{ports}\} \triangleleft \text{deny-all-from-to src-net dest-net}$

definition

deny-from-to-ports :: $'\gamma::\text{port set} \Rightarrow '\alpha::\text{adr net} \Rightarrow '\alpha::\text{adr net} \Rightarrow$
 $((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$ **where**

deny-from-to-ports ports src-net dest-net =
 $\{pa. \text{dest-port } pa \in \text{ports}\} \triangleleft \text{deny-all-from-to src-net dest-net}$

definition

allow-all-from-port-tos :: $'\alpha :: \text{adr net} \Rightarrow (' \gamma :: \text{port}) \text{ set} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$

where

allow-all-from-port-tos src-net s-port dest-net
 $= \{pa. \text{ dest-port } pa \in s\text{-port}\} \triangleleft \text{allow-all-from-to src-net dest-net}$

As before, we put all the rules into one lemma called PortCombinators to ease writing later.

lemmas *PortCombinatorsCore* =

allow-all-from-port-def deny-all-from-port-def allow-all-to-port-def
deny-all-to-port-def allow-all-from-to-port-def
deny-all-from-to-port-def
allow-from-ports-to-def allow-from-to-ports-def
deny-from-ports-to-def deny-from-to-ports-def
allow-all-from-port-to-def deny-all-from-port-to-def
allow-from-port-to-def allow-from-to-port-def deny-from-to-port-def
deny-from-port-to-def allow-all-from-port-tos-def

lemmas *PortCombinators* =

PortCombinatorsCore PolicyCombinators

end

6.4. Policy Combinators with Ports and Protocols

theory *ProtocolPortCombinators*

imports *PortCombinators*

begin

This theory defines policy combinators for those network models which have ports. They are provided in addition to the the ones defined in the PolicyCombinators theory.

This theory requires from the network models a definition for the two following constants:

- *src_port* :: $(' \alpha, ' \beta) \text{ packet} \Rightarrow (' \gamma :: \text{port})$
- *dest_port* :: $(' \alpha, ' \beta) \text{ packet} \Rightarrow (' \gamma :: \text{port})$

definition

allow-all-from-port-prot :: $\text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow (' \gamma :: \text{port}) \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$

where

allow-all-from-port-prot p src-net s-port =
 $\{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{allow-all-from-port src-net s-port}$

definition

$$\text{deny-all-from-port-prot} \quad :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$$
where

$$\begin{aligned} \text{deny-all-from-port-prot } p \text{ src-net s-port} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{deny-all-from-port src-net s-port} \end{aligned}$$
definition

$$\text{allow-all-to-port-prot} \quad :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit}) \text{ where}$$

$$\begin{aligned} \text{allow-all-to-port-prot } p \text{ dest-net d-port} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{allow-all-to-port dest-net d-port} \end{aligned}$$
definition

$$\text{deny-all-to-port-prot} \quad :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit}) \text{ where}$$

$$\begin{aligned} \text{deny-all-to-port-prot } p \text{ dest-net d-port} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{deny-all-to-port dest-net d-port} \end{aligned}$$
definition

$$\text{allow-all-from-port-to-prot} :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$$
where

$$\begin{aligned} \text{allow-all-from-port-to-prot } p \text{ src-net s-port dest-net} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{allow-all-from-port-to src-net s-port dest-net} \end{aligned}$$
definition

$$\text{deny-all-from-port-to-prot} :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit})$$
where

$$\begin{aligned} \text{deny-all-from-port-to-prot } p \text{ src-net s-port dest-net} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{deny-all-from-port-to src-net s-port dest-net} \end{aligned}$$
definition

$$\text{allow-all-from-port-to-port-prot} :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit}) \text{ where}$$

$$\begin{aligned} \text{allow-all-from-port-to-port-prot } p \text{ src-net s-port dest-net d-port} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{allow-all-from-port-to-port src-net s-port dest-net d-port} \end{aligned}$$
definition

$$\text{deny-all-from-port-to-port-prot} :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit}) \text{ where}$$

$$\begin{aligned} \text{deny-all-from-port-to-port-prot } p \text{ src-net s-port dest-net d-port} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{deny-all-from-port-to-port src-net s-port dest-net d-port} \end{aligned}$$
definition

$$\text{allow-all-from-to-port-prot} \quad :: \text{protocol} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \alpha :: \text{adr net} \Rightarrow ' \gamma :: \text{port} \Rightarrow ((' \alpha, ' \beta) \text{ packet} \mapsto \text{unit}) \text{ where}$$

$$\begin{aligned} \text{allow-all-from-to-port-prot } p \text{ src-net dest-net d-port} = \\ \{pa. \text{ dest-protocol } pa = p\} \triangleleft \text{allow-all-from-to-port src-net dest-net d-port} \end{aligned}$$
definition

$deny-all-from-to-port-prot :: protocol => 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow 'γ::port \Rightarrow$
 $((\alpha, \beta)\ packet \mapsto unit) \textbf{ where}$
 $deny-all-from-to-port-prot\ p\ src-net\ dest-net\ d-port =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft deny-all-from-to-port\ src-net\ dest-net\ d-port$

definition

$allow-from-port-to-prot :: protocol => 'γ::port \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow ((\alpha, \beta)\ packet$
 $\mapsto unit)$
where
 $allow-from-port-to-prot\ p\ port\ src-net\ dest-net =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft allow-from-port-to\ port\ src-net\ dest-net$

definition

$deny-from-port-to-prot :: protocol => 'γ::port \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow ((\alpha, \beta)\ packet$
 $\mapsto unit)$
where
 $deny-from-port-to-prot\ p\ port\ src-net\ dest-net =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft deny-from-port-to\ port\ src-net\ dest-net$

definition

$allow-from-to-port-prot :: protocol => 'γ::port \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow ((\alpha, \beta)\ packet$
 $\mapsto unit)$
where
 $allow-from-to-port-prot\ p\ port\ src-net\ dest-net =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft allow-from-to-port\ port\ src-net\ dest-net$

definition

$deny-from-to-port-prot :: protocol => 'γ::port \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow ((\alpha, \beta)\ packet$
 $\mapsto unit)$
where
 $deny-from-to-port-prot\ p\ port\ src-net\ dest-net =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft deny-from-to-port\ port\ src-net\ dest-net$

definition

$allow-from-ports-to-prot :: protocol => 'γ::port\ set \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow$
 $((\alpha, \beta)\ packet \mapsto unit) \textbf{ where}$
 $allow-from-ports-to-prot\ p\ ports\ src-net\ dest-net =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft allow-from-ports-to\ ports\ src-net\ dest-net$

definition

$allow-from-to-ports-prot :: protocol => 'γ::port\ set \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow$
 $((\alpha, \beta)\ packet \mapsto unit) \textbf{ where}$
 $allow-from-to-ports-prot\ p\ ports\ src-net\ dest-net =$
 $\{pa.\ dest-protocol\ pa = p\} \triangleleft allow-from-to-ports\ ports\ src-net\ dest-net$

definition

$deny-from-ports-to-prot :: protocol => 'γ::port\ set \Rightarrow 'α::adr\ net \Rightarrow 'α::adr\ net \Rightarrow$
 $((\alpha, \beta)\ packet \mapsto unit) \textbf{ where}$

deny-from-ports-to-prot p *ports* *src-net* *dest-net* =
 $\{pa. \text{dest-protocol } pa = p\} \triangleleft \text{deny-from-ports-to } ports \text{ src-net dest-net}$

definition

deny-from-to-ports-prot :: *protocol* $\Rightarrow \gamma :: port \text{ set} \Rightarrow \alpha :: adr \text{ net} \Rightarrow \alpha :: adr \text{ net} \Rightarrow$
 $((\alpha, \beta) \text{ packet} \mapsto unit)$ **where**
deny-from-to-ports-prot p *ports* *src-net* *dest-net* =
 $\{pa. \text{dest-protocol } pa = p\} \triangleleft \text{deny-from-to-ports } ports \text{ src-net dest-net}$

As before, we put all the rules into one lemma to ease writing later.

lemmas *ProtocolCombinatorsCore* =

allow-all-from-port-prot-def deny-all-from-port-prot-def allow-all-to-port-prot-def
deny-all-to-port-prot-def allow-all-from-to-port-prot-def
deny-all-from-to-port-prot-def
allow-from-ports-to-prot-def allow-from-to-ports-prot-def
deny-from-ports-to-prot-def deny-from-to-ports-prot-def
allow-all-from-port-to-prot-def deny-all-from-port-to-prot-def
allow-from-port-to-prot-def allow-from-to-port-prot-def deny-from-to-port-prot-def
deny-from-port-to-prot-def

lemmas *ProtocolCombinators* = *PortCombinators.PortCombinators*
ProtocolCombinatorsCore

end

6.5. Ports

theory *Ports*

imports *Main*

begin

This theory can be used if we want to specify the port numbers by names denoting their default Integer values. If you want to use them, please add *Ports* to the simplifier before test data generation.

definition *http::int* **where** *http* = 80

lemma *http1*: $x \neq 80 \implies x \neq http$
 $\langle proof \rangle$

lemma *http2*: $x \neq 80 \implies http \neq x$
 $\langle proof \rangle$

definition *smtp::int* **where** *smtp* = 25

lemma *smtp1*: $x \neq 25 \implies x \neq \text{smtp}$
 $\langle \text{proof} \rangle$

lemma *smtp2*: $x \neq 25 \implies \text{smtp} \neq x$
 $\langle \text{proof} \rangle$

definition *ftp::int* **where** *ftp* = 21

lemma *ftp1*: $x \neq 21 \implies x \neq \text{ftp}$
 $\langle \text{proof} \rangle$

lemma *ftp2*: $x \neq 21 \implies \text{ftp} \neq x$
 $\langle \text{proof} \rangle$

And so on for all desired port numbers.

lemmas *Ports* = *http1 http2 ftp1 ftp2 smtp1 smtp2*

end

7. Network Address Translation

theory *NAT*

imports *../PacketFilter/PacketFilter*

begin

definition *src2pool* :: $'\alpha \text{ set} \Rightarrow (' \alpha :: \text{adr}, ' \beta) \text{ packet} \Rightarrow (' \alpha, ' \beta) \text{ packet set}$ **where**
 $\text{src2pool } t = (\lambda p. (\{(i, s, d, da). (i = \text{id } p \wedge s \in t \wedge d = \text{dest } p \wedge da = \text{content } p)\}))$

definition *src2poolAP* **where**
 $\text{src2poolAP } t = A_f (\text{src2pool } t)$

definition *srcNat2pool* :: $' \alpha \text{ set} \Rightarrow ' \alpha \text{ set} \Rightarrow (' \alpha :: \text{adr}, ' \beta) \text{ packet} \mapsto (' \alpha, ' \beta) \text{ packet set}$ **where**
 $\text{srcNat2pool } \text{srcs } \text{transl} = \{x. \text{src } x \in \text{srcs}\} \triangleleft (\text{src2poolAP } \text{transl})$

definition *src2poolPort* :: $\text{int set} \Rightarrow (\text{adr}_{ip}, ' \beta) \text{ packet} \Rightarrow (\text{adr}_{ip}, ' \beta) \text{ packet set}$ **where**
 $\text{src2poolPort } t = (\lambda p. (\{(i, (s1, s2), (d1, d2), da). (i = \text{id } p \wedge s1 \in t \wedge s2 = (\text{snd } (\text{src } p)) \wedge d1 = (\text{fst } (\text{dest } p)) \wedge d2 = \text{snd } (\text{dest } p) \wedge da = \text{content } p)\}))$

definition *src2poolPort-Protocol* :: $\text{int set} \Rightarrow (\text{adr}_{ipp}, ' \beta) \text{ packet} \Rightarrow (\text{adr}_{ipp}, ' \beta) \text{ packet set}$ **where**

$$\begin{aligned} \text{src2poolPort-Protocol } t &= (\lambda p. (\{(i, (s1, s2, s3), (d1, d2, d3), da). \\ &\quad (i = id\ p \wedge s1 \in t \wedge s2 = (fst\ (snd\ (src\ p))) \wedge s3 = snd\ (snd\ (src\ p)) \wedge \\ &\quad (d1, d2, d3) = dest\ p \wedge da = content\ p\}\})) \end{aligned}$$

definition *srcNat2pool-IntPort* :: *address set* \Rightarrow *address set* \Rightarrow
(adr_{ip}, 'β) *packet* \mapsto *(adr_{ip}, 'β)* *packet set* **where**
srcNat2pool-IntPort *srcs* *transl* =
 $\{x. fst\ (src\ x) \in srcs\} \triangleleft (A_f\ (src2poolPort\ transl))$

definition *srcNat2pool-IntProtocolPort* :: *int set* \Rightarrow *int set* \Rightarrow
(adr_{ipp}, 'β) *packet* \mapsto *(adr_{ipp}, 'β)* *packet set* **where**
srcNat2pool-IntProtocolPort *srcs* *transl* =
 $\{x. (fst\ ((src\ x))) \in srcs\} \triangleleft (A_f\ (src2poolPort-Protocol\ transl))$

definition *srcPat2poolPort-t* :: *int set* \Rightarrow *(adr_{ip}, 'β)* *packet* \Rightarrow *(adr_{ip}, 'β)* *packet set* **where**
srcPat2poolPort-t *t* = $(\lambda p. (\{(i, (s1, s2), (d1, d2), da). \\ (i = id\ p \wedge s1 \in t \wedge d1 = (fst\ (dest\ p)) \wedge d2 = snd\ (dest\ p) \wedge da = content\ p)\}))$

definition *srcPat2poolPort-Protocol-t* :: *int set* \Rightarrow *(adr_{ipp}, 'β)* *packet* \Rightarrow *(adr_{ipp}, 'β)* *packet set*
where
srcPat2poolPort-Protocol-t *t* = $(\lambda p. (\{(i, (s1, s2, s3), (d1, d2, d3), da). \\ (i = id\ p \wedge s1 \in t \wedge s3 = src\text{-protocol}\ p \wedge (d1, d2, d3) = dest\ p \wedge da = content\ p)\}))$

definition *srcPat2pool-IntPort* :: *int set* \Rightarrow *int set* \Rightarrow *(adr_{ip}, 'β)* *packet* \mapsto
(adr_{ip}, 'β) *packet set* **where**
srcPat2pool-IntPort *srcs* *transl* =
 $\{x. (fst\ (src\ x)) \in srcs\} \triangleleft (A_f\ (srcPat2poolPort-t\ transl))$

definition *srcPat2pool-IntProtocol* ::
int set \Rightarrow *int set* \Rightarrow *(adr_{ipp}, 'β)* *packet* \mapsto *(adr_{ipp}, 'β)* *packet set* **where**
srcPat2pool-IntProtocol *srcs* *transl* =
 $\{x. (fst\ (src\ x)) \in srcs\} \triangleleft (A_f\ (srcPat2poolPort-Protocol-t\ transl))$

The following lemmas are used for achieving a "standard" output format of NATted packets for use in a firewall execution tool.

lemma *datasimp*: $\{(i, (s1, s2, s3), aba).$
 $\forall a\ aa\ b\ ba. aba = ((a, aa, b), ba) \longrightarrow i = i1 \wedge s1 = i101 \wedge$
 $s3 = iudp \wedge a = i110 \wedge aa = X606X3 \wedge b = X607X4 \wedge ba = data\} =$
 $\{(i, (s1, s2, s3), aba).$
 $i = i1 \wedge s1 = i101 \wedge s3 = iudp \wedge (\lambda ((a, aa, b), ba). a = i110 \wedge aa = X606X3 \wedge$
 $b = X607X4 \wedge ba = data)\ aba\}$
<proof>

lemma *datasimp2*: $\{(i, (s1, s2, s3), aba).$

$$\begin{aligned} & \forall a \text{ aa } b \text{ ba}. aba = ((a, aa, b), ba) \longrightarrow i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge \\ & s2 = i1 \wedge a = i110 \wedge aa = i4 \wedge b = iudp \wedge ba = data\} \\ = & \{(i, (s1, s2, s3), aba). \\ & i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge s2 = i1 \wedge (\lambda ((a, aa, b), ba). a = i110 \wedge \\ & aa = i4 \wedge b = iudp \wedge ba = data) aba\} \\ & \langle proof \rangle \end{aligned}$$

lemma *datasimp3*: $\{(i, (s1, s2, s3), aba).$

$$\begin{aligned} & ALL a \text{ aa } b \text{ ba}. aba = ((a, aa, b), ba) \dashrightarrow i = i1 \ \& \ i115 < s1 \ \& \ s1 < i124 \ \& \\ & s3 = iudp \ \& \ s2 = ii1 \ \& \ a = i110 \ \& \ aa = i3 \ \& \ b = itcp \ \& \ ba = data\} = \\ & \{(i, (s1, s2, s3), aba). \\ & i = i1 \ \& \ i115 < s1 \ \& \ s1 < i124 \ \& \ s3 = iudp \ \& \ s2 = ii1 \ \& \ (\lambda ((a, aa, b), ba). \\ & a = i110 \ \& \ aa = i3 \ \& \ b = itcp \ \& \ ba = data) aba\} \\ & \langle proof \rangle \end{aligned}$$

lemma *datasimp4*: $\{(i, (s1, s2, s3), aba).$

$$\begin{aligned} & \forall a \text{ aa } b \text{ ba}. aba = ((a, aa, b), ba) \longrightarrow i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge \\ & s2 = ii1 \wedge a = i110 \wedge aa = i7 \wedge b = itcp \wedge ba = data\} = \\ & \{(i, (s1, s2, s3), aba). \\ & i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge s2 = ii1 \wedge (\lambda ((a, aa, b), ba). a = i110 \wedge \\ & aa = i7 \wedge b = itcp \wedge ba = data) aba\} \\ & \langle proof \rangle \end{aligned}$$

lemma *datasimp5*: $\{(i, (s1, s2, s3), aba).$

$$\begin{aligned} & i = i1 \wedge s1 = i101 \wedge s3 = iudp \wedge (\lambda ((a, aa, b), ba). a = i110 \wedge aa = X606X3 \wedge \\ & b = X607X4 \wedge ba = data) aba\} \\ = & \{(i, (s1, s2, s3), (a, aa, b), ba). \\ & i = i1 \wedge s1 = i101 \wedge s3 = iudp \wedge a = i110 \wedge aa = X606X3 \wedge \\ & b = X607X4 \wedge ba = data\} \\ & \langle proof \rangle \end{aligned}$$

lemma *datasimp6*: $\{(i, (s1, s2, s3), aba).$

$$\begin{aligned} & i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge s2 = i1 \wedge (\lambda ((a, aa, b), ba). a = i110 \wedge \\ & aa = i4 \wedge b = iudp \wedge ba = data) aba\} = \\ & \{(i, (s1, s2, s3), (a, aa, b), ba). \\ & i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge s2 = i1 \wedge a = i110 \wedge aa = i4 \wedge b = iudp \wedge ba \\ & = data\} \\ & \langle proof \rangle \end{aligned}$$

lemma *datasimp7*: $\{(i, (s1, s2, s3), aba).$

$$\begin{aligned} & i = i1 \ \& \ i115 < s1 \ \& \ s1 < i124 \ \& \ s3 = iudp \ \& \ s2 = ii1 \ \& \ (\lambda ((a, aa, b), ba). a = \\ & i110 \ \& \\ & aa = i3 \ \& \ b = itcp \ \& \ ba = data) aba\} = \\ & \{(i, (s1, s2, s3), (a, aa, b), ba). \\ & i = i1 \ \& \ i115 < s1 \ \& \ s1 < i124 \ \& \ s3 = iudp \ \& \ s2 = ii1 \ \& \ a = i110 \ \& \ aa = i3 \ \& \ b = itcp \\ & \ \& \ ba = data\} \\ & \langle proof \rangle \end{aligned}$$

```

lemma datasimp8:  $\{(i, (s1, s2, s3), aba).$ 
 $i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge s2 = ii1 \wedge (\lambda ((a,aa,b),ba). a = i110 \wedge aa =$ 
 $i7 \wedge$ 
 $b = itcp \wedge ba = data) aba\}$ 
 $= \{(i, (s1, s2, s3), (a,aa,b),ba).$ 
 $i = i1 \wedge s1 = i132 \wedge s3 = iudp \wedge s2 = ii1 \wedge a = i110 \wedge aa = i7 \wedge b = itcp \wedge$ 
 $ba = data\}$ 
 $\langle proof \rangle$ 

```

```

lemmas datasimps = datasimp datasimp2 datasimp3 datasimp4
 $datasimp5 datasimp6 datasimp7 datasimp8$ 

```

```

lemmas NATLemmas = src2pool-def src2poolPort-def
 $src2poolPort-Protocol-def src2poolAP-def srcNat2pool-def$ 
 $srcNat2pool-IntProtocolPort-def srcNat2pool-IntPort-def$ 
 $srcPat2poolPort-t-def srcPat2poolPort-Protocol-t-def$ 
 $srcPat2pool-IntPort-def srcPat2pool-IntProtocol-def$ 

```

end

8. Policy Normalisation

```

theory
 $FWNormalisationCore$ 
imports
 $../PacketFilter/PacketFilter$ 
begin

```

This theory contains all the definitions used for policy normalisation as described in [1]. The normalisation procedure transforms policies into semantically equivalent ones which are "easier" to test. It is organized into nine phases. We impose the following two restrictions on the input policies:

- Each policy must contain a **DenyAll** rule. If this restriction were to be lifted, the

`insertDenies` phase would have to be adjusted accordingly.

- For each pair of networks n_1 and n_2 , the networks are either disjoint or equal. If this restriction were to be lifted, we would need some additional phases before the start of the normalisation procedure presented below. This rule would split single rules into several by splitting up the networks such that they are all pairwise disjoint or equal. Such a transformation is clearly semantics-preserving and the condition would hold after these phases.

As a result, the procedure generates a list of policies, in which:

- each element of the list contains a policy which completely specifies the blocking behavior between two networks, and
- there are no shadowed rules.

This result is desirable since the test case generation for rules between networks A and B is independent of the rules that specify the behavior for traffic flowing between networks C and D . Thus, the different segments of the policy can be processed individually. The normalization procedure does not aim to minimize the number of rules. While it does remove unnecessary ones, it also adds new ones, enabling a policy to be split into several independent parts.

Policy transformations are functions that map policies to policies. We decided to represent policy transformations as *syntactic rules*; this choice paves the way for expressing the entire normalisation process inside HOL by functions manipulating abstract policy syntax.

8.1. Basics

We define a very simple policy language:

```
datatype (' $\alpha$ , ' $\beta$ ) Combinators =
  DenyAll
| DenyAllFromTo ' $\alpha$  ' $\alpha$ 
| AllowPortFromTo ' $\alpha$  ' $\alpha$  ' $\beta$ 
| Conc ((' $\alpha$ , ' $\beta$ ) Combinators) ((' $\alpha$ , ' $\beta$ ) Combinators) (infixr  $\oplus$  80)
```

And define the semantic interpretation of it. For technical reasons, we fix here the type to policies over IntegerPort addresses. However, we could easily provide definitions for other address types as well, using a generic `consts` for the type definition and a `primrec` definition for each desired address model.

8.2. Auxiliary definitions and functions.

This subsection defines several functions which are useful later for the combinators, invariants, and proofs.

```
fun srcNet where
  srcNet (DenyAllFromTo  $x$   $y$ ) =  $x$ 
| srcNet (AllowPortFromTo  $x$   $y$   $p$ ) =  $x$ 
```

```
|srcNet DenyAll = undefined
|srcNet (v ⊕ va) = undefined
```

```
fun destNet where
  destNet (DenyAllFromTo x y) = y
|destNet (AllowPortFromTo x y p) = y
|destNet DenyAll = undefined
|destNet (v ⊕ va) = undefined
```

```
fun srcnets where
  srcnets DenyAll = []
|srcnets (DenyAllFromTo x y) = [x]
|srcnets (AllowPortFromTo x y p) = [x]
|(srcnets (x ⊕ y)) = (srcnets x)@(srcnets y)
```

```
fun destnets where
  destnets DenyAll = []
|destnets (DenyAllFromTo x y) = [y]
|destnets (AllowPortFromTo x y p) = [y]
|(destnets (x ⊕ y)) = (destnets x)@(destnets y)
```

```
fun (sequential) net-list-aux where
  net-list-aux [] = []
|net-list-aux (DenyAll#xs) = net-list-aux xs
|net-list-aux ((DenyAllFromTo x y)#xs) = x#y#(net-list-aux xs)
|net-list-aux ((AllowPortFromTo x y p)#xs) = x#y#(net-list-aux xs)
|net-list-aux ((x⊕y)#xs) = (net-list-aux [x])@(net-list-aux [y])@(net-list-aux xs)
```

```
fun net-list where net-list p = remdups (net-list-aux p)
```

```
definition bothNets where bothNets x = (zip (srcnets x) (destnets x))
```

```
fun (sequential) normBothNets where
  normBothNets ((a,b)#xs) = (if ((b,a) ∈ set xs) ∨ (a,b) ∈ set xs)
    then (normBothNets xs)
    else (a,b)#(normBothNets xs)
|normBothNets x = x
```

```
fun makeSets where
  makeSets ((a,b)#xs) = ({a,b}#(makeSets xs))
|makeSets [] = []
```

```
fun bothNet where
  bothNet DenyAll = {}
|bothNet (DenyAllFromTo a b) = {a,b}
|bothNet (AllowPortFromTo a b p) = {a,b}
|bothNet (v ⊕ va) = undefined
```

Nets_List provides from a list of rules a list where the entries are the appearing sets of

source and destination network of each rule.

definition *Nets-List* **where** *Nets-List* $x = \text{makeSets } (\text{normBothNets } (\text{bothNets } x))$

fun (*sequential*) *first-srcNet* **where**
 $\text{first-srcNet } (x \oplus y) = \text{first-srcNet } x$
 $|\text{first-srcNet } x = \text{srcNet } x$

fun (*sequential*) *first-destNet* **where**
 $\text{first-destNet } (x \oplus y) = \text{first-destNet } x$
 $|\text{first-destNet } x = \text{destNet } x$

fun (*sequential*) *first-bothNet* **where**
 $\text{first-bothNet } (x \oplus y) = \text{first-bothNet } x$
 $|\text{first-bothNet } x = \text{bothNet } x$

fun (*sequential*) *in-list* **where**
 $\text{in-list } \text{DenyAll } l = \text{True}$
 $|\text{in-list } x \ l = (\text{bothNet } x \in \text{set } l)$

fun *all-in-list* **where**
 $\text{all-in-list } [] \ l = \text{True}$
 $|\text{all-in-list } (x \# xs) \ l = (\text{in-list } x \ l \wedge \text{all-in-list } xs \ l)$

fun (*sequential*) *member* **where**
 $\text{member } a \ (x \oplus xs) = ((\text{member } a \ x) \vee (\text{member } a \ xs))$
 $|\text{member } a \ x = (a = x)$

fun *sdnets* **where**
 $\text{sdnets } \text{DenyAll} = \{\}$
 $|\text{sdnets } (\text{DenyAllFromTo } a \ b) = \{(a, b)\}$
 $|\text{sdnets } (\text{AllowPortFromTo } a \ b \ c) = \{(a, b)\}$
 $|\text{sdnets } (a \oplus b) = \text{sdnets } a \cup \text{sdnets } b$

definition *packet-Nets* **where** *packet-Nets* $x \ a \ b = ((\text{src } x \sqsubset a \wedge \text{dest } x \sqsubset b) \vee (\text{src } x \sqsubset b \wedge \text{dest } x \sqsubset a))$

definition *subnetsOfAdr* **where** *subnetsOfAdr* $a = \{x. a \sqsubset x\}$

definition *fst-set* **where** *fst-set* $s = \{a. \exists \ b. (a, b) \in s\}$

definition *snd-set* **where** *snd-set* $s = \{a. \exists \ b. (b, a) \in s\}$

fun *memberP* **where**
 $\text{memberP } r \ (x \# xs) = (\text{member } r \ x \vee \text{memberP } r \ xs)$
 $|\text{memberP } r \ [] = \text{False}$

fun *firstList* **where**
 $\text{firstList } (x \# xs) = (\text{first-bothNet } x)$

$|firstList \ [] = \{\}$

8.3. Invariants

If there is a DenyAll, it is at the first position

```
fun wellformed-policy1:: (('α, 'β) Combinators) list ⇒ bool where
  wellformed-policy1 [] = True
| wellformed-policy1 (x#xs) = (DenyAll ∉ (set xs))
```

There is a DenyAll at the first position

```
fun wellformed-policy1-strong:: (('α, 'β) Combinators) list ⇒ bool
where
  wellformed-policy1-strong [] = False
| wellformed-policy1-strong (x#xs) = (x=DenyAll ∧ (DenyAll ∉ (set xs)))
```

All two networks are either disjoint or equal.

definition netsDistinct **where** $netsDistinct\ a\ b = (\neg (\exists\ x. x \sqsubset a \wedge x \sqsubset b))$

definition twoNetsDistinct **where**
 $twoNetsDistinct\ a\ b\ c\ d = (netsDistinct\ a\ c \vee netsDistinct\ b\ d)$

definition allNetsDistinct **where**
 $allNetsDistinct\ p = (\forall\ a\ b. (a \neq b \wedge a \in set\ (net-list\ p) \wedge b \in set\ (net-list\ p)) \longrightarrow netsDistinct\ a\ b)$

definition disjSD-2 **where**
 $disjSD-2\ x\ y = (\forall\ a\ b\ c\ d. ((a,b) \in sdnets\ x \wedge (c,d) \in sdnets\ y \longrightarrow (twoNetsDistinct\ a\ b\ c\ d \wedge twoNetsDistinct\ a\ b\ d\ c)))$

The policy is given as a list of single rules.

```
fun singleCombinators where
  singleCombinators [] = True
| singleCombinators ((x⊕y)#xs) = False
| singleCombinators (x#xs) = singleCombinators xs
```

definition onlyTwoNets **where**
 $onlyTwoNets\ x = ((\exists\ a\ b. (sdnets\ x = \{(a,b)\})) \vee (\exists\ a\ b. sdnets\ x = \{(a,b),(b,a)\}))$

Each entry of the list contains rules between two networks only.

```
fun OnlyTwoNets where
  OnlyTwoNets (DenyAll#xs) = OnlyTwoNets xs
| OnlyTwoNets (x#xs) = (onlyTwoNets x ∧ OnlyTwoNets xs)
| OnlyTwoNets [] = True
```

```
fun noDenyAll where
  noDenyAll (x#xs) = ((¬ member DenyAll x) ∧ noDenyAll xs)
| noDenyAll [] = True
```

```

fun noDenyAll1 where
  noDenyAll1 (DenyAll#xs) = noDenyAll xs
| noDenyAll1 xs = noDenyAll xs

fun separated where
  separated (x#xs) = (( $\forall$  s. s  $\in$  set xs  $\longrightarrow$  disjSD-2 x s)  $\wedge$  separated xs)
| separated [] = True

fun NetsCollected where
  NetsCollected (x#xs) = (((first-bothNet x  $\neq$  firstList xs)  $\longrightarrow$ 
    ( $\forall$  a $\in$ set xs. first-bothNet x  $\neq$  first-bothNet a))  $\wedge$  NetsCollected (xs))
| NetsCollected [] = True

fun NetsCollected2 where
  NetsCollected2 (x#xs) = (xs = []  $\vee$  (first-bothNet x  $\neq$  firstList xs  $\wedge$ 
    NetsCollected2 xs))
| NetsCollected2 [] = True

```

8.4. Transformations

The following two functions transform a policy into a list of single rules and vice-versa - by staying on the combinator level.

```

fun policy2list::('α, 'β) Combinators  $\Rightarrow$ 
  (('α, 'β) Combinators) list where
  policy2list (x  $\oplus$  y) = (concat [(policy2list x), (policy2list y)])
| policy2list x = [x]

fun list2FWpolicy::(('α, 'β) Combinators) list  $\Rightarrow$ 
  (('α, 'β) Combinators) where
  list2FWpolicy [] = undefined
| list2FWpolicy (x#[]) = x
| list2FWpolicy (x#y) = x  $\oplus$  (list2FWpolicy y)

```

Remove all the rules appearing before a DenyAll. There are two alternative versions.

```

fun removeShadowRules1 where
  removeShadowRules1 (x#xs) = (if (DenyAll  $\in$  set xs)
    then ((removeShadowRules1 xs))
    else x#xs)
| removeShadowRules1 [] = []

fun removeShadowRules1-alternative-rev where
  removeShadowRules1-alternative-rev [] = []
| removeShadowRules1-alternative-rev (DenyAll#xs) = [DenyAll]
| removeShadowRules1-alternative-rev [x] = [x]
| removeShadowRules1-alternative-rev (x#xs) =
  x#(removeShadowRules1-alternative-rev xs)

```

definition removeShadowRules1-alternative **where**

```

removeShadowRules1-alternative p =
    rev (removeShadowRules1-alternative-rev (rev p))

```

Remove all the rules which allow a port, but are shadowed by a deny between these subnets

```

fun removeShadowRules2:: (('α, 'β) Combinators) list ⇒
    (('α, 'β) Combinators) list
where
    (removeShadowRules2 ((AllowPortFromTo x y p)#z)) =
        (if (((DenyAllFromTo x y) ∈ set z))
            then ((removeShadowRules2 z))
            else (((AllowPortFromTo x y p)#(removeShadowRules2 z))))
| removeShadowRules2 (x#y) = x#(removeShadowRules2 y)
| removeShadowRules2 [] = []

```

Sorting a pocliy. We first need to define an ordering on rules. This ordering depends on the *Nets_List* of a policy.

```

fun smaller :: ('α, 'β) Combinators ⇒
    ('α, 'β) Combinators ⇒
    (('α) set) list ⇒ bool
where
    smaller DenyAll x l = True
| smaller x DenyAll l = False
| smaller x y l =
    ((x = y) ∨
        (if (bothNet x) = (bothNet y) then
            (case y of (DenyAllFromTo a b) ⇒ (x = DenyAllFromTo b a)
            | - ⇒ True)
        else
            (position (bothNet x) l <= position (bothNet y) l)))

```

We provide two different sorting algorithms: Quick Sort (qsort) and Insertion Sort (sort)

```

fun qsort where
    qsort [] l = []
| qsort (x#xs) l = (qsort [y←xs. ¬ (smaller x y l)] l) @ [x] @ (qsort [y←xs. smaller x y l] l)

```

lemma qsort-permutes:
 set (qsort xs l) = set xs
 ⟨proof⟩

lemma set-qsort [simp]: set (qsort xs l) = set xs
 ⟨proof⟩

```

fun insert where
    insert a [] l = [a]
| insert a (x#xs) l = (if (smaller a x l) then a#x#xs else x#(insert a xs l))

```

fun sort **where**


```

    sort [] l = []
| sort (x#xs) l = insert x (sort xs l) l

```

```

fun sorted where
    sorted [] l  $\longleftrightarrow$  True |
    sorted [x] l  $\longleftrightarrow$  True |
    sorted (x#y#zs) l  $\longleftrightarrow$  smaller x y l  $\wedge$  sorted (y#zs) l

```

```

fun separate where
    separate (DenyAll#x) = DenyAll#(separate x)
| separate (x#y#z) = (if (first-bothNet x = first-bothNet y)
    then (separate ((x $\oplus$ y)#z))
    else (x#(separate(y#z))))
| separate x = x

```

Insert the DenyAllFromTo rules, such that traffic between two networks can be tested individually

```

fun insertDenies where
    insertDenies (x#xs) = (case x of DenyAll  $\Rightarrow$  (DenyAll#(insertDenies xs))
    | -  $\Rightarrow$  (DenyAllFromTo (first-srcNet x) (first-destNet x)  $\oplus$ 
    (DenyAllFromTo (first-destNet x) (first-srcNet x)  $\oplus$  x)#
    (insertDenies xs))
| insertDenies [] = []

```

Remove duplicate rules. This is especially necessary as insertDenies might have inserted duplicate rules.

The second function is supposed to work on a list of policies. Only rules which are duplicated within the same policy are removed.

```

fun removeDuplicates where
    removeDuplicates (x $\oplus$ xs) = (if member x xs then (removeDuplicates xs)
    else x $\oplus$ (removeDuplicates xs))
| removeDuplicates x = x

```

```

fun removeAllDuplicates where
    removeAllDuplicates (x#xs) = ((removeDuplicates (x))#(removeAllDuplicates xs))
| removeAllDuplicates x = x

```

Insert a DenyAll at the beginning of a policy.

```

fun insertDeny where
    insertDeny (DenyAll#xs) = DenyAll#xs
| insertDeny xs = DenyAll#xs

```

definition sort' p l = sort l p

definition qsort' p l = qsort l p

declare *dom-eq-empty-conv* [*simp del*]

fun *list2policyR*::(('α, 'β) *Combinators*) *list* ⇒
 (('α, 'β) *Combinators*) **where**
 list2policyR (*x*#[]) = *x*
 list2policyR (*x*#*y*) = (*list2policyR* *y*) ⊕ *x*
 list2policyR [] = *undefined*

We provide the definitions for two address representations.

8.5. IntPort

fun *C* :: (*adr_{ip}* *net*, *port*) *Combinators* ⇒ (*adr_{ip}*, *DummyContent*) *packet* ↦ *unit*
where
 C *DenyAll* = *deny-all*
 C (*DenyAllFromTo* *x y*) = *deny-all-from-to* *x y*
 C (*AllowPortFromTo* *x y p*) = *allow-from-to-port* *p x y*
 C (*x* ⊕ *y*) = *C* *x* ++ *C* *y*

fun *CRotate* :: (*adr_{ip}* *net*, *port*) *Combinators* ⇒ (*adr_{ip}*, *DummyContent*) *packet* ↦ *unit*
where
 CRotate *DenyAll* = *C* *DenyAll*
 CRotate (*DenyAllFromTo* *x y*) = *C* (*DenyAllFromTo* *x y*)
 CRotate (*AllowPortFromTo* *x y p*) = *C* (*AllowPortFromTo* *x y p*)
 CRotate (*x* ⊕ *y*) = ((*CRotate* *y*) ++ ((*CRotate* *x*)))

fun *rotatePolicy* **where**
 rotatePolicy *DenyAll* = *DenyAll*
 rotatePolicy (*DenyAllFromTo* *a b*) = *DenyAllFromTo* *a b*
 rotatePolicy (*AllowPortFromTo* *a b p*) = *AllowPortFromTo* *a b p*
 rotatePolicy (*a*⊕*b*) = (*rotatePolicy* *b*) ⊕ (*rotatePolicy* *a*)

lemma *check*:: *rev* (*policy2list* (*rotatePolicy* *p*)) = *policy2list* *p*
 ⟨*proof*⟩

All rules appearing at the left of a *DenyAllFromTo*, have disjunct domains from it (except *DenyAll*)

fun (*sequential*) *wellformed-policy2* **where**
 wellformed-policy2 [] = *True*
 wellformed-policy2 (*DenyAll*#*xs*) = *wellformed-policy2* *xs*
 wellformed-policy2 (*x*#*xs*) = ((∀ *c a b*. *c* = *DenyAllFromTo* *a b* ∧ *c* ∈ *set xs* →
 Map.dom (*C* *x*) ∩ *Map.dom* (*C* *c*) = {}) ∧ *wellformed-policy2* *xs*)

An allow rule is disjunct with all rules appearing at the right of it. This invariant is not necessary as it is a consequence from others, but facilitates some proofs.

```
fun (sequential) wellformed-policy3::((adrip net,port) Combinators) list  $\Rightarrow$  bool where
  wellformed-policy3 [] = True
| wellformed-policy3 ((AllowPortFromTo a b p)#xs) = (( $\forall$  r. r  $\in$  set xs  $\longrightarrow$ 
  dom (C r)  $\cap$  dom (C (AllowPortFromTo a b p)) = {})  $\wedge$  wellformed-policy3 xs)
| wellformed-policy3 (x#xs) = wellformed-policy3 xs
```

definition

```
normalize' p = (removeAllDuplicates o insertDenies o separate o
  (sort' (Nets-List p)) o removeShadowRules2 o remdups o
  (removeShadowRules3 C) o insertDeny o removeShadowRules1 o
  policy2list) p
```

definition

```
normalizeQ' p = (removeAllDuplicates o insertDenies o separate o
  (qsort' (Nets-List p)) o removeShadowRules2 o remdups o
  (removeShadowRules3 C) o insertDeny o removeShadowRules1 o
  policy2list) p
```

definition normalize ::

```
(adrip net, port) Combinators  $\Rightarrow$ 
(adrip net, port) Combinators list
```

where

```
normalize p = (removeAllDuplicates (insertDenies (separate (sort
  (removeShadowRules2 (remdups ((removeShadowRules3 C) (insertDeny
  (removeShadowRules1 (policy2list p)))))) ((Nets-List p))))))
```

definition

```
normalize-manual-order p l = removeAllDuplicates (insertDenies (separate
  (sort (removeShadowRules2 (remdups ((removeShadowRules3 C) (insertDeny
  (removeShadowRules1 (policy2list p)))))) ((l))))
```

definition normalizeQ ::

```
(adrip net, port) Combinators  $\Rightarrow$ 
(adrip net, port) Combinators list
```

where

```
normalizeQ p = (removeAllDuplicates (insertDenies (separate (qsort
  (removeShadowRules2 (remdups ((removeShadowRules3 C) (insertDeny
  (removeShadowRules1 (policy2list p)))))) ((Nets-List p))))))
```

definition

```
normalize-manual-orderQ p l = removeAllDuplicates (insertDenies (separate
  (qsort (removeShadowRules2 (remdups ((removeShadowRules3 C) (insertDeny
```

$(\text{removeShadowRules1 } (\text{policy2list } p)))))) ((l))))$

Of course, `normalize` is equal to `normalize'`, the latter looks nicer though.

lemma `normalize = normalize'`

<proof>

declare `C.simps [simp del]`

8.6. TCP_UDP_IntegerPort

fun `Cp` :: $(\text{adr}_{ipp} \text{ net}, \text{protocol} \times \text{port}) \text{ Combinators} \Rightarrow$
 $(\text{adr}_{ipp}, \text{DummyContent}) \text{ packet} \mapsto \text{unit}$

where

$Cp \text{ DenyAll} = \text{deny-all}$
 $Cp (\text{DenyAllFromTo } x \ y) = \text{deny-all-from-to } x \ y$
 $Cp (\text{AllowPortFromTo } x \ y \ p) = \text{allow-from-to-port-prot } (\text{fst } p) (\text{snd } p) \ x \ y$
 $Cp (x \oplus y) = Cp \ x \ ++ \ Cp \ y$

fun `Dp` :: $(\text{adr}_{ipp} \text{ net}, \text{protocol} \times \text{port}) \text{ Combinators} \Rightarrow$
 $(\text{adr}_{ipp}, \text{DummyContent}) \text{ packet} \mapsto \text{unit}$

where

$Dp \text{ DenyAll} = Cp \text{ DenyAll}$
 $Dp (\text{DenyAllFromTo } x \ y) = Cp (\text{DenyAllFromTo } x \ y)$
 $Dp (\text{AllowPortFromTo } x \ y \ p) = Cp (\text{AllowPortFromTo } x \ y \ p)$
 $Dp (x \oplus y) = Cp (y \oplus x)$

All rules appearing at the left of a `DenyAllFromTo`, have disjunct domains from it (except `DenyAll`)

fun $(\text{sequential}) \text{ wellformed-policy2Pr}$ **where**

$\text{wellformed-policy2Pr } [] = \text{True}$
 $| \text{wellformed-policy2Pr } (\text{DenyAll} \# xs) = \text{wellformed-policy2Pr } xs$
 $| \text{wellformed-policy2Pr } (x \# xs) = ((\forall \ c \ a \ b. \ c = \text{DenyAllFromTo } a \ b \wedge c \in \text{set } xs \longrightarrow$
 $\text{Map.dom } (Cp \ x) \cap \text{Map.dom } (Cp \ c) = \{\}) \wedge \text{wellformed-policy2Pr } xs)$

An allow rule is disjunct with all rules appearing at the right of it. This invariant is not necessary as it is a consequence from others, but facilitates some proofs.

fun $(\text{sequential}) \text{ wellformed-policy3Pr} :: ((\text{adr}_{ipp} \text{ net}, \text{protocol} \times \text{port}) \text{ Combinators}) \text{ list} \Rightarrow \text{bool}$

where

$\text{wellformed-policy3Pr } [] = \text{True}$
 $| \text{wellformed-policy3Pr } ((\text{AllowPortFromTo } a \ b \ p) \# xs) = ((\forall \ r. \ r \in \text{set } xs \longrightarrow$
 $\text{dom } (Cp \ r) \cap \text{dom } (Cp (\text{AllowPortFromTo } a \ b \ p)) = \{\}) \wedge \text{wellformed-policy3Pr } xs)$
 $| \text{wellformed-policy3Pr } (x \# xs) = \text{wellformed-policy3Pr } xs$

definition

$normalizePr' :: (adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators}$
 $\Rightarrow (adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators list where}$
 $normalizePr' p = (removeAllDuplicates \circ insertDenies \circ separate \circ$
 $\quad (sort' (Nets-List p)) \circ removeShadowRules2 \circ remdups \circ$
 $\quad (removeShadowRules3 Cp) \circ insertDeny \circ removeShadowRules1 \circ$
 $\quad policy2list) p$

definition $normalizePr ::$

$(adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators}$
 $\Rightarrow (adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators list where}$
 $normalizePr p = (removeAllDuplicates (insertDenies (separate (sort$
 $\quad (removeShadowRules2 (remdups ((removeShadowRules3 Cp) (insertDeny$
 $\quad (removeShadowRules1 (policy2list p)))))) ((Nets-List p))))))$

definition

$normalize-manual-orderPr p l = removeAllDuplicates (insertDenies (separate$
 $\quad (sort (removeShadowRules2 (remdups ((removeShadowRules3 Cp) (insertDeny$
 $\quad (removeShadowRules1 (policy2list p)))))) ((l))))$

definition

$normalizePrQ' :: (adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators}$
 $\Rightarrow (adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators list where}$
 $normalizePrQ' p = (removeAllDuplicates \circ insertDenies \circ separate \circ$
 $\quad (qsort' (Nets-List p)) \circ removeShadowRules2 \circ remdups \circ$
 $\quad (removeShadowRules3 Cp) \circ insertDeny \circ removeShadowRules1 \circ$
 $\quad policy2list) p$

definition $normalizePrQ ::$

$(adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators}$
 $\Rightarrow (adr_{ipp} \text{ net}, protocol \times port) \text{ Combinators list where}$
 $normalizePrQ p = (removeAllDuplicates (insertDenies (separate (qsort$
 $\quad (removeShadowRules2 (remdups ((removeShadowRules3 Cp) (insertDeny$
 $\quad (removeShadowRules1 (policy2list p)))))) ((Nets-List p))))))$

definition

$normalize-manual-orderPrQ p l = removeAllDuplicates (insertDenies (separate$
 $\quad (qsort (removeShadowRules2 (remdups ((removeShadowRules3 Cp) (insertDeny$
 $\quad (removeShadowRules1 (policy2list p)))))) ((l))))$

Of course, $normalize$ is equal to $normalize'$, the latter looks nicer though.

lemma $normalizePr = normalizePr'$

$\langle proof \rangle$

The following definition helps in creating the test specification for the individual parts of a normalized policy.

possible dynamic ports or not to allow that protocol at all. Neither of these options is satisfactory. In the first case, all ports above 1024 would have to be opened which introduces a big security hole in the system, in the second case users wouldn't be very happy. A firewall which tracks the state of the TCP connections on a system doesn't help here either, as the opening and closing of the ports takes place on the application layer. Therefore, a firewall needs to have some knowledge of the application protocols being run and track the states of these protocols. We next model this behaviour.

The key point of our model is the idea that a policy remains the same as before: a mapping from packet to packet out. We still specify for every packet, based on its source and destination address, the expected action. The only thing that changes now is that this mapping is allowed to change over time. This indicates that our test data will not consist of single packets but rather of sequences thereof.

At first we hence need a state. It is a tuple from some memory to be refined later and the current policy.

type-synonym $(\alpha, \beta, \gamma) \text{ FWState} = \alpha \times ((\beta, \gamma) \text{ packet} \mapsto \text{unit})$

Having a state, we need of course some state transitions. Such a transition can happen every time a new packet arrives. State transitions can be modelled using a state-exception monad. We provide two types of firewall monads: one

type-synonym $(\alpha, \beta, \gamma) \text{ FWStateTransitionP} = (\beta, \gamma) \text{ packet} \Rightarrow ((\beta, \gamma) \text{ packet} \mapsto \text{unit}) \text{ decision}, (\alpha, \beta, \gamma) \text{ FWState} \text{ MON}_{SE}$

type-synonym $(\alpha, \beta, \gamma) \text{ FWStateTransition} = ((\beta, \gamma) \text{ packet} \times (\alpha, \beta, \gamma) \text{ FWState}) \rightarrow (\alpha, \beta, \gamma) \text{ FWState}$

The memory could be modelled as a list of accepted packets.

type-synonym $(\beta, \gamma) \text{ history} = (\beta, \gamma) \text{ packet list}$

fun *packet-with-id* **where**

packet-with-id $[] = []$
packet-with-id $(x \# xs) \ i = (\text{if } id \ x = i \text{ then } (x \# (\text{packet-with-id } xs \ i)) \text{ else } (\text{packet-with-id } xs \ i))$

fun *ids1* **where**

ids1 $i \ (x \# xs) = (id \ x = i \wedge ids1 \ i \ xs)$
ids1 $i \ [] = \text{True}$

fun *ids* **where**

ids $a \ (x \# xs) = (\text{NetworkCore.id } x \in a \wedge ids \ a \ xs)$
ids $a \ [] = \text{True}$

definition *applyPolicy*:: $(i \times (i \mapsto o)) \mapsto o$

where *applyPolicy* $= (\lambda (x, z). z \ x)$

end

9.2. FTP Protocol

```
theory FTP
imports
  Stateful
begin
```

9.2.1. The protocol syntax

The File Transfer Protocol FTP is a well known example of a protocol which uses dynamic ports and is therefore a natural choice to use as an example for our model.

We model only a simplified version of the FTP protocol over IntegerPort addresses, still containing all messages that matter for our purposes. It consists of the following four messages:

1. *init*: The client contacts the server indicating his wish to get some data.
2. *ftp-port-request p*: The client, usually after having received an acknowledgement of the server, indicates a port number on which he wants to receive the data.
3. *ftp-ftp-data*: The server sends the requested data over the new channel. There might be an arbitrary number of such messages, including zero.
4. *ftp-close*: The client closes the connection. The dynamic port gets closed again.

The content field of a packet therefore now consists of either one of those four messages or a default one.

datatype *msg* = *ftp-init* | *ftp-port-request port* | *ftp-data* | *ftp-close* | *ftp-other*

We now also make use of the ID field of a packet. It is used as session ID and we make the assumption that they are all unique among different protocol runs.

At first, we need some predicates which check if a packet is a specific FTP message and has the correct session ID.

definition

is-init :: *id* \Rightarrow (*adr_{ip}*, *msg*)*packet* \Rightarrow *bool* **where**
is-init = (λ *i p*. (*id p* = *i* \wedge *content p* = *ftp-init*))

definition

is-ftp-port-request :: *id* \Rightarrow *port* \Rightarrow (*adr_{ip}*, *msg*) *packet* \Rightarrow *bool* **where**
is-ftp-port-request = (λ *i port p*. (*id p* = *i* \wedge *content p* = *ftp-port-request port*))

definition

$is\text{-}ftp\text{-}data :: id \Rightarrow (adr_{ip}, msg) \text{ packet} \Rightarrow bool$ **where**
 $is\text{-}ftp\text{-}data = (\lambda i p. (id\ p = i \wedge content\ p = ftp\text{-}data))$

definition

$is\text{-}ftp\text{-}close :: id \Rightarrow (adr_{ip}, msg) \text{ packet} \Rightarrow bool$ **where**
 $is\text{-}ftp\text{-}close = (\lambda i p. (id\ p = i \wedge content\ p = ftp\text{-}close))$

definition

$port\text{-}open :: (adr_{ip}, msg) \text{ history} \Rightarrow id \Rightarrow port \Rightarrow bool$ **where**
 $port\text{-}open = (\lambda L a p. (not\text{-}before\ (is\text{-}ftp\text{-}close\ a)\ (is\text{-}ftp\text{-}port\text{-}request\ a\ p)\ L))$

definition

$is\text{-}ftp\text{-}other :: id \Rightarrow (adr_{ip}, msg) \text{ packet} \Rightarrow bool$ **where**
 $is\text{-}ftp\text{-}other = (\lambda i p. (id\ p = i \wedge content\ p = ftp\text{-}other))$

fun are-ftp-other where

$are\text{-}ftp\text{-}other\ i\ (x\#xs) = (is\text{-}ftp\text{-}other\ i\ x \wedge are\text{-}ftp\text{-}other\ i\ xs)$
 $| are\text{-}ftp\text{-}other\ i\ [] = True$

9.2.2. The protocol policy specification

We now have to model the respective state transitions. It is important to note that state transitions themselves allow all packets which are allowed by the policy, not only those which are allowed by the protocol. Their only task is to change the policy. As an alternative, we could have decided that they only allow packets which follow the protocol (e.g. come on the correct ports), but this should in our view rather be reflected in the policy itself.

Of course, not every message changes the policy. In such cases, we do not have to model different cases, one is enough. In our example, only messages 2 and 4 need special transitions. The default says that if the policy accepts the packet, it is added to the history, otherwise it is simply dropped. The policy remains the same in both cases.

fun last-opened-port where

$last\text{-}opened\text{-}port\ i\ ((j,s,d,ftp\text{-}port\text{-}request\ p)\#xs) = (if\ i=j\ then\ p\ else\ last\text{-}opened\text{-}port\ i\ xs)$
 $| last\text{-}opened\text{-}port\ i\ (x\#xs) = last\text{-}opened\text{-}port\ i\ xs$
 $| last\text{-}opened\text{-}port\ x\ [] = undefined$

fun FTP-STA :: ((adr_{ip},msg) history, adr_{ip}, msg) FWStateTransition where

$FTP\text{-}STA\ ((i,s,d,ftp\text{-}port\text{-}request\ pr), (log\ ,\ pol)) =$
 $(if\ before\ (Not\ o\ is\text{-}ftp\text{-}close\ i)\ (is\text{-}init\ i)\ log\ \wedge$
 $\quad dest\text{-}port\ (i,s,d,ftp\text{-}port\text{-}request\ pr) = (21::port)$
 $\quad then\ Some\ (((i,s,d,ftp\text{-}port\text{-}request\ pr)\#log\ ,$
 $\quad \quad (allow\text{-}from\text{-}to\text{-}port\ pr\ (subnet\text{-}of\ d)\ (subnet\text{-}of\ s)) \oplus pol))$
 $\quad else\ Some\ (((i,s,d,ftp\text{-}port\text{-}request\ pr)\#log\ ,pol)))$

$$\begin{aligned}
|FTP\text{-}STA \ ((i,s,d,ftp\text{-}close), (log, pol)) = \\
& (if \ (\exists \ p. \ port\text{-}open \ log \ i \ p) \wedge \ dest\text{-}port \ (i,s,d,ftp\text{-}close) = (21::port) \\
& \quad then \ Some \ ((i,s,d,ftp\text{-}close)\#log, \\
& \quad \quad \quad deny\text{-}from\text{-}to\text{-}port \ (last\text{-}opened\text{-}port \ i \ log) \ (subnet\text{-}of \ d)(subnet\text{-}of \ s) \oplus \ pol) \\
& \quad else \ Some \ (((i,s,d,ftp\text{-}close)\#log, pol)))
\end{aligned}$$

$$|FTP\text{-}STA \ (p, s) = Some \ (p\#(fst \ s), snd \ s)$$

fun *FTP-STD* ::
 ((*adr_{ip}*,*msg*) *history*, *adr_{ip}*, *msg*) *FWStateTransition*
where
FTP-STD (*p*,*s*) = *Some s*

definition *TRPolicy* :: (*adr_{ip}*,*msg*)*packet* × (*adr_{ip}*,*msg*)*history* × ((*adr_{ip}*,*msg*)*packet* ↦ *unit*)
 ↦ (*unit* × (*adr_{ip}*,*msg*)*history* × ((*adr_{ip}*,*msg*)*packet* ↦ *unit*))
where *TRPolicy* = ((*FTP-STA*,*FTP-STD*) ⊗_∇ *applyPolicy*) o (λ(*x*,(*y*,*z*)).((*x*,*z*),(*x*,(*y*,*z*))))

definition *TRPolicy_{Mon}*
where *TRPolicy_{Mon}* = *policy2MON*(*TRPolicy*)

If required to contain the policy in the output

definition *TRPolicy_{Mon}'*
where *TRPolicy_{Mon}'* = *policy2MON* (((λ(*x*,*y*,*z*). (*z*,(*y*,*z*))) o-f *TRPolicy*))

Now we specify our test scenario in more detail. We could test:

- one correct FTP-Protocol run,
- several runs after another,
- several runs interleaved,
- an illegal protocol run, or
- several illegal protocol runs.

We only do the the simplest case here: one correct protocol run.

There are four different states which are modelled as a datatype.

datatype *ftp-states* = *S0* | *S1* | *S2* | *S3*

The following constant is *True* for all sets which are correct FTP runs for a given source and destination address, ID, and data-port number.

fun
is-ftp :: *ftp-states* ⇒ *adr_{ip}* ⇒ *adr_{ip}* ⇒ *id* ⇒ *port* ⇒
 (*adr_{ip}*,*msg*) *history* ⇒ *bool*

where

$$\begin{aligned}
& \text{is-ftp } H \text{ } c \text{ } s \text{ } i \text{ } p \text{ } [] = (H=S3) \\
& | \text{is-ftp } H \text{ } c \text{ } s \text{ } i \text{ } p \text{ } (x\#InL) = (snd \text{ } s = 21 \wedge ((\lambda (id, sr, de, co). (((id = i \wedge (\\
& \quad (H=ftp-states.S2 \wedge sr = c \wedge de = s \wedge co = ftp-init \wedge \text{is-ftp } S3 \text{ } c \text{ } s \text{ } i \text{ } p \text{ } InL) \vee \\
& \quad (H=ftp-states.S1 \wedge sr = c \wedge de = s \wedge co = ftp-port-request \text{ } p \wedge \text{is-ftp } S2 \text{ } c \text{ } s \text{ } i \text{ } p \text{ } InL) \vee \\
& \quad (H=ftp-states.S1 \wedge sr = s \wedge de = (fst \text{ } c, p) \wedge co = ftp-data \wedge \text{is-ftp } S1 \text{ } c \text{ } s \text{ } i \text{ } p \text{ } InL) \vee \\
& \quad (H=ftp-states.S0 \wedge sr = c \wedge de = s \wedge co = ftp-close \wedge \text{is-ftp } S1 \text{ } c \text{ } s \text{ } i \text{ } p \text{ } InL)))))) x))
\end{aligned}$$

definition *is-single-ftp-run* :: $adr_{ip} \text{ } src \Rightarrow adr_{ip} \text{ } dest \Rightarrow id \Rightarrow port \Rightarrow (adr_{ip}, msg) \text{ history set}$

where *is-single-ftp-run* $s \text{ } d \text{ } i \text{ } p = \{x. (\text{is-ftp } S0 \text{ } s \text{ } d \text{ } i \text{ } p \text{ } x)\}$

The following constant then returns a set of all the historys which denote such a normal behaviour FTP run, again for a given source and destination address, ID, and data-port.

The following definition returns the set of all possible interleaving of two correct FTP protocol runs.

definition

$$\begin{aligned}
& \text{ftp-2-interleaved} :: adr_{ip} \text{ } src \Rightarrow adr_{ip} \text{ } dest \Rightarrow id \Rightarrow port \Rightarrow \\
& \quad adr_{ip} \text{ } src \Rightarrow adr_{ip} \text{ } dest \Rightarrow id \Rightarrow port \Rightarrow \\
& \quad (adr_{ip}, msg) \text{ history set } \mathbf{where} \\
& \text{ftp-2-interleaved } s1 \text{ } d1 \text{ } i1 \text{ } p1 \text{ } s2 \text{ } d2 \text{ } i2 \text{ } p2 = \\
& \quad \{x. (\text{is-ftp } S0 \text{ } s1 \text{ } d1 \text{ } i1 \text{ } p1 \text{ } (\text{packet-with-id } x \text{ } i1)) \wedge \\
& \quad (\text{is-ftp } S0 \text{ } s2 \text{ } d2 \text{ } i2 \text{ } p2 \text{ } (\text{packet-with-id } x \text{ } i2)))\}
\end{aligned}$$

lemma *subnetOf-lemma*: $(a::int) \neq (c::int) \implies \forall x \in \text{subnet-of } (a, b::port). (c, d) \notin x$
 <proof>

lemma *subnetOf-lemma2*: $\forall x \in \text{subnet-of } (a::int, b::port). (a, b) \in x$
 <proof>

lemma *subnetOf-lemma3*: $(\exists x. x \in \text{subnet-of } (a::int, b::port))$
 <proof>

lemma *subnetOf-lemma4*: $\exists x \in \text{subnet-of } (a::int, b::port). (a, c::port) \in x$
 <proof>

lemma *port-open-lemma*: $\neg (Ex (\text{port-open } [] (x::port)))$
 <proof>

lemmas *FTPLemmas* = *TRPolicy-def applyPolicy-def policy2MON-def*
Let-def in-subnet-def src-def
dest-def subnet-of-def
is-init-def p-accept-def port-open-def is-ftp-data-def is-ftp-close-def

```

is-ftp-port-request-def
content-def PortCombinators
exI subnetOf-lemma
subnetOf-lemma2 subnetOf-lemma3 subnetOf-lemma4
NetworkCore.id-def adrip Lemmas port-open-lemma
bind-SE-def unit-SE-def valid-SE-def

```

end

```

theory FTP-WithPolicy
imports
  FTP
begin

```

FTP where the policy is part of the output.

```

definition POL :: 'a ⇒ 'a
where POL x = x

```

Variant 2 takes the policy into the output

```

fun FTP-STP ::
  ((id ⇒ port), adrip, msg) FWStateTransitionP
where

```

```

  FTP-STP (i,s,d,ftp-port-request pr) (ports, policy) =
    (if p-accept (i,s,d,ftp-port-request pr) policy then
      Some (allow (POL ((allow-from-to-port pr (subnet-of d) (subnet-of s)) ⊕ policy)),
        ( (ports(i→pr)),(allow-from-to-port pr (subnet-of d) (subnet-of s))
          ⊕ policy))
    else (Some (deny (POL policy),(ports,policy))))

```

```

  | FTP-STP (i,s,d,ftp-close) (ports,policy) =
    (if (p-accept (i,s,d,ftp-close) policy) then
      case ports i of
        Some pr ⇒
          Some(allow (POL (deny-from-to-port pr (subnet-of d) (subnet-of s) ⊕ policy)),
            ports(i:=None),
            deny-from-to-port pr (subnet-of d) (subnet-of s) ⊕ policy)
        | None ⇒ Some(allow (POL policy), ports, policy)
      else Some (deny (POL policy), ports, policy))

```

```

|FTP-STP p x = (if p-accept p (snd x)
                 then Some (allow (POL (snd x)),((fst x),snd x))
                 else Some (deny (POL (snd x)),(fst x,snd x)))

```

end

9.3. VoIP Protocol

```

theory VOIP
imports Stateful
begin

```

After the FTP-Protocol which was rather simple we show the strength of the model with a more current and especially much more complicated example, namely Voice over IP (VoIP). VoIP is standardized by the ITU-T under the name H.323, which can be seen as an "umbrella standard" which aggregates standards for multimedia conferencing over packet-based networks (for a good overview of the protocol suite, see [?]). H.323 poses many problems to firewalls. These problems include (taken from [?]):

- An H.323 call is made up of many different simultaneous connections.
- Most connections are made to dynamic ports.
- The addresses and port numbers are exchanged within the data stream of the next higher connection.
- Calls can be initiated from outside the firewall.

Again we only consider a simplified VoIP scenario with the following seven messages which are grouped into four subprotocols (see Figure ??):

- Registration and Admission (H.225, port 1719): The caller contacts its gatekeeper with a call request. The gatekeeper either rejects or confirms the request, returning the address of the callee in the latter case.
 - Admission Request (ARQ)
 - Admission Reject (ARJ)
 - Admission Confirm (ACF) 'a
- Call Signaling (Q.931, port 1720) The caller and the callee agree on the dynamic ports over which the call will take place.
 - Setup port
 - Connect port

- Stream (dynamic ports). The call itself. In reality, several connections are used here.
- Fin (port 1720).

The two main differences to FTP are:

- In VoIP, we deal with three different entities: the caller, the callee, and the gatekeeper.
- We do not know in advance which entity will close the connection.

We model the protocol as seen from a firewall at the caller, namely we are not interested in the messages from the callee to its gatekeeper. Incoming calls are not modelled either, they would require a different set of state transitions.

The content of a packet now consists of one of the seven messages or a default one. It is parameterized with the type of the address that the gatekeeper returns.

```
datatype 'a voip-msg = ARQ
    | ACF 'a
    | ARJ
    | Setup port
    | Connect port
    | Stream
    | Fin
    | other
```

As before, we need operators which check if a packet contains a specific content and ID, respectively if such a packet has appeared in the trace.

definition

```
is-arq :: NetworkCore.id ⇒ ('a::adr, 'b voip-msg) packet ⇒ bool where
is-arq i p = (NetworkCore.id p = i ∧ content p = ARQ)
```

definition

```
is-fin :: id ⇒ ('a::adr, 'b voip-msg) packet ⇒ bool where
is-fin i p = (id p = i ∧ content p = Fin)
```

definition

```
is-connect :: id ⇒ port ⇒ ('a::adr, 'b voip-msg) packet ⇒ bool where
is-connect i port p = (id p = i ∧ content p = Connect port)
```

definition

```
is-setup :: id ⇒ port ⇒ ('a::adr, 'b voip-msg) packet ⇒ bool where
is-setup i port p = (id p = i ∧ content p = Setup port)
```

We need also an operator *ports-open* to get access to the two dynamic ports.

definition

```
ports-open :: id ⇒ port × port ⇒ (adrip, 'a voip-msg) history ⇒ bool where
ports-open i p L = ((not-before (is-fin i) (is-setup i (fst p)) L) ∧
    not-before (is-fin i) (is-connect i (snd p)) L)
```

As we do not know which entity closes the connection, we define an operator which checks if the closer is the caller.

fun

```
src-is-initiator :: id ⇒ adrip ⇒ (adrip, 'b voip-msg) history ⇒ bool where
src-is-initiator i a [] = False
|src-is-initiator i a (p#S) = (((id p = i) ∧
                               (∃ port. content p = Setup port) ∧
                               ((fst (src p) = fst a))) ∨
                               (src-is-initiator i a S))
```

The first state transition is for those messages which do not change the policy. In this scenario, this only happens for the Stream messages.

definition *subnet-of-adr* **where**

subnet-of-adr $x = \{(a,b). a = x\}$

fun *VOIP-STA* ::

```
((adrip, address voip-msg) history, adrip, address voip-msg) FWStateTransition
where
```

```
VOIP-STA ((a,c,d,ARQ), (InL, policy)) =
  Some (((a,c,d, ARQ)#InL,
    (allow-from-to-port (1719::port)(subnet-of d) (subnet-of c)) ⊕ policy))
```

```
| VOIP-STA ((a,c,d,ARJ), (InL, policy)) =
  (if (not-before (is-fin a) (is-arq a) InL)
    then Some (((a,c,d,ARJ)#InL,
      deny-from-to-port (14::port) (subnet-of c) (subnet-of d) ⊕ policy))
    else Some (((a,c,d,ARJ)#InL,policy)))
```

```
| VOIP-STA ((a,c,d,ACF callee), (InL, policy)) =
  Some (((a,c,d,ACF callee)#InL,
    allow-from-to-port (1720::port) (subnet-of-adr callee) (subnet-of d) ⊕
    allow-from-to-port (1720::port) (subnet-of d) (subnet-of-adr callee) ⊕
    deny-from-to-port (1719::port) (subnet-of d) (subnet-of c) ⊕
    policy))
```

```
| VOIP-STA ((a,c,d, Setup port), (InL, policy)) =
  Some (((a,c,d,Setup port)#InL,
    allow-from-to-port port (subnet-of d) (subnet-of c) ⊕ policy))
```

```
| VOIP-STA ((a,c,d, Connect port), (InL, policy)) =
  Some (((a,c,d,Connect port)#InL,
    allow-from-to-port port (subnet-of d) (subnet-of c) ⊕ policy))
```

$| VOIP-STA ((a,c,d,Fin), (InL,policy)) =$
 $(if \exists p1 p2. ports-open a (p1,p2) InL then ($
 $(if src-is-initiator a c InL$
 $then (Some (((a,c,d,Fin)\#InL,$
 $(deny-from-to-port (1720::int) (subnet-of c) (subnet-of d)) \oplus$
 $(deny-from-to-port (snd (SOME p. ports-open a p InL))$
 $(subnet-of c) (subnet-of d)) \oplus$
 $(deny-from-to-port (fst (SOME p. ports-open a p InL))$
 $(subnet-of d) (subnet-of c)) \oplus policy)))$
 $else (Some (((a,c,d,Fin)\#InL,$
 $(deny-from-to-port (1720::int) (subnet-of c) (subnet-of d)) \oplus$
 $(deny-from-to-port (fst (SOME p. ports-open a p InL))$
 $(subnet-of c) (subnet-of d)) \oplus$
 $(deny-from-to-port (snd (SOME p. ports-open a p InL))$
 $(subnet-of d) (subnet-of c)) \oplus policy))))$
 $else$
 $(Some (((a,c,d,Fin)\#InL,policy))))$

$| VOIP-STA (p, (InL, policy)) =$
 $Some ((p\#InL,policy))$

fun *VOIP-STD* **where**
 $VOIP-STD (p,s) = Some s$

definition *VOIP-TRPolicy* **where**
 $VOIP-TRPolicy = policy2MON ($
 $((VOIP-STA, VOIP-STD) \otimes_{\nabla} applyPolicy) o (\lambda (x,(y,z)). ((x,z),(x,(y,z))))$

For a full protocol run, six states are needed.

datatype *voip-states* = *S0* | *S1* | *S2* | *S3* | *S4* | *S5*

The constant *is-voip* checks if a trace corresponds to a legal VoIP protocol, given the IP-addresses of the three entities, the ID, and the two dynamic ports.

fun *is-voip* :: *voip-states* \Rightarrow *address* \Rightarrow *address* \Rightarrow *address* \Rightarrow *id* \Rightarrow *port* \Rightarrow
 $port \Rightarrow (adr_{ip}, address\ voip-msg)\ history \Rightarrow bool$

where

$is-voip\ H\ s\ d\ g\ i\ p1\ p2\ [] = (H = S5)$
 $| is-voip\ H\ s\ d\ g\ i\ p1\ p2\ (x\#InL) =$
 $(((\lambda (id,sr,de,co).$
 $((id = i \wedge$
 $(H = S4 \wedge ((sr = (s,1719) \wedge de = (g,1719) \wedge co = ARQ \wedge$
 $is-voip\ S5\ s\ d\ g\ i\ p1\ p2\ InL)))) \vee$
 $(H = S0 \wedge sr = (g,1719) \wedge de = (s,1719) \wedge co = ARJ \wedge$

$$\begin{aligned}
& \text{is-voip } S4 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S3 \wedge sr = (g, 1719) \wedge de = (s, 1719) \wedge co = ACF \ d \wedge \\
& \quad \text{is-voip } S4 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S2 \wedge sr = (s, 1720) \wedge de = (d, 1720) \wedge co = Setup \ p1 \wedge \\
& \quad \text{is-voip } S3 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S1 \wedge sr = (d, 1720) \wedge de = (s, 1720) \wedge co = Connect \ p2 \wedge \\
& \quad \text{is-voip } S2 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S1 \wedge sr = (s, p1) \wedge de = (d, p2) \wedge co = Stream \wedge \\
& \quad \text{is-voip } S1 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S1 \wedge sr = (d, p2) \wedge de = (s, p1) \wedge co = Stream \wedge \\
& \quad \text{is-voip } S1 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S0 \wedge sr = (d, 1720) \wedge de = (s, 1720) \wedge co = Fin \wedge \\
& \quad \text{is-voip } S1 \ s \ d \ g \ i \ p1 \ p2 \ InL) \vee \\
& (H = S0 \wedge sr = (s, 1720) \wedge de = (d, 1720) \wedge co = Fin \wedge \\
& \quad \text{is-voip } S1 \ s \ d \ g \ i \ p1 \ p2 \ InL)))))) \ x)
\end{aligned}$$

Finally, *NB-voip* returns the set of protocol traces which correspond to a correct protocol run given the three addresses, the ID, and the two dynamic ports.

definition

NB-voip :: *address* \Rightarrow *address* \Rightarrow *address* \Rightarrow *id* \Rightarrow *port* \Rightarrow *port* \Rightarrow
 (*adr_{ip}*, *address voip-msg*) *history set* **where**
NB-voip *s d g i p1 p2* = {*x*. (*is-voip* *S0 s d g i p1 p2 x*)}

end

9.4. FTP and VoIP Protocol Interleaved

theory *FTPVOIP*

imports

FTP-WithPolicy VOIP

begin

datatype *ftpvoip* = *ARQ*
 | *ACF int*
 | *ARJ*
 | *Setup port*
 | *Connect port*
 | *Stream*
 | *Fin*
 | *ftp-init*
 | *ftp-port-request port*
 | *ftp-data*
 | *ftp-close*
 | *other*

We now also make use of the ID field of a packet. It is used as session ID and we make

the assumption that they are all unique among different protocol runs.

At first, we need some predicates which check if a packet is a specific FTP message and has the correct session ID.

definition

$FTPVOIP\text{-}is\text{-}init :: id \Rightarrow (adr_{ip}, ftpvoip) \text{ packet} \Rightarrow bool$ **where**
 $FTPVOIP\text{-}is\text{-}init = (\lambda i p. (id\ p = i \wedge content\ p = ftp\text{-}init))$

definition

$FTPVOIP\text{-}is\text{-}port\text{-}request :: id \Rightarrow port \Rightarrow (adr_{ip}, ftpvoip) \text{ packet} \Rightarrow bool$ **where**
 $FTPVOIP\text{-}is\text{-}port\text{-}request = (\lambda i port p. (id\ p = i \wedge content\ p = ftp\text{-}port\text{-}request\ port))$

definition

$FTPVOIP\text{-}is\text{-}data :: id \Rightarrow (adr_{ip}, ftpvoip) \text{ packet} \Rightarrow bool$ **where**
 $FTPVOIP\text{-}is\text{-}data = (\lambda i p. (id\ p = i \wedge content\ p = ftp\text{-}data))$

definition

$FTPVOIP\text{-}is\text{-}close :: id \Rightarrow (adr_{ip}, ftpvoip) \text{ packet} \Rightarrow bool$ **where**
 $FTPVOIP\text{-}is\text{-}close = (\lambda i p. (id\ p = i \wedge content\ p = ftp\text{-}close))$

definition

$FTPVOIP\text{-}port\text{-}open :: (adr_{ip}, ftpvoip) \text{ history} \Rightarrow id \Rightarrow port \Rightarrow bool$ **where**
 $FTPVOIP\text{-}port\text{-}open = (\lambda L a p. (not\text{-}before\ (FTPVOIP\text{-}is\text{-}close\ a)\ (FTPVOIP\text{-}is\text{-}port\text{-}request\ a\ p)\ L))$

definition

$FTPVOIP\text{-}is\text{-}other :: id \Rightarrow (adr_{ip}, ftpvoip) \text{ packet} \Rightarrow bool$ **where**
 $FTPVOIP\text{-}is\text{-}other = (\lambda i p. (id\ p = i \wedge content\ p = other))$

fun $FTPVOIP\text{-}are\text{-}other$ **where**

$FTPVOIP\text{-}are\text{-}other\ i\ (x\#xs) = (FTPVOIP\text{-}is\text{-}other\ i\ x \wedge FTPVOIP\text{-}are\text{-}other\ i\ xs)$
 $| FTPVOIP\text{-}are\text{-}other\ i\ [] = True$

fun $last\text{-}opened\text{-}port$ **where**

$last\text{-}opened\text{-}port\ i\ ((j,s,d,ftp\text{-}port\text{-}request\ p)\#xs) = (if\ i=j\ then\ p\ else\ last\text{-}opened\text{-}port\ i\ xs)$
 $| last\text{-}opened\text{-}port\ i\ (x\#xs) = last\text{-}opened\text{-}port\ i\ xs$
 $| last\text{-}opened\text{-}port\ x\ [] = undefined$

fun $FTPVOIP\text{-}FTP\text{-}STA ::$

$((adr_{ip}, ftpvoip) \text{ history}, adr_{ip}, ftpvoip) \text{ FWStateTransition}$

where

$FTPVOIP\text{-}FTP\text{-}STA\ ((i,s,d,ftp\text{-}port\text{-}request\ pr), (InL, policy)) =$
 $(if\ not\text{-}before\ (FTPVOIP\text{-}is\text{-}close\ i)\ (FTPVOIP\text{-}is\text{-}init\ i)\ InL \wedge$
 $dest\text{-}port\ (i,s,d,ftp\text{-}port\text{-}request\ pr) = (21::port)\ then$
 $Some\ (((i,s,d,ftp\text{-}port\text{-}request\ pr)\#InL, policy) ++$

(allow-from-to-port pr (subnet-of d) (subnet-of s)))
 else Some (((i,s,d,ftp-port-request pr)#InL,policy)))

|FTPVOIP-FTP-STA ((i,s,d,ftp-close), (InL,policy)) =
 (if (∃ p. FTPVOIP-port-open InL i p) ∧ dest-port (i,s,d,ftp-close) = (21::port)
 then Some ((i,s,d,ftp-close)#InL, policy ++
 deny-from-to-port (last-opened-port i InL) (subnet-of d) (subnet-of s))
 else Some (((i,s,d,ftp-close)#InL, policy)))

|FTPVOIP-FTP-STA (p, s) = Some (p#(fst s),snd s)

fun FTPVOIP-FTP-STD ::
 ((adr_{ip}, ftpvoip) history, adr_{ip}, ftpvoip) FWStateTransition
where
 FTPVOIP-FTP-STD (p,s) = Some s

definition

FTPVOIP-is-arq :: NetworkCore.id ⇒ ('a::adr, ftpvoip) packet ⇒ bool **where**
 FTPVOIP-is-arq i p = (NetworkCore.id p = i ∧ content p = ARQ)

definition

FTPVOIP-is-fin :: id ⇒ ('a::adr, ftpvoip) packet ⇒ bool **where**
 FTPVOIP-is-fin i p = (id p = i ∧ content p = Fin)

definition

FTPVOIP-is-connect :: id ⇒ port ⇒ ('a::adr, ftpvoip) packet ⇒ bool **where**
 FTPVOIP-is-connect i port p = (id p = i ∧ content p = Connect port)

definition

FTPVOIP-is-setup :: id ⇒ port ⇒ ('a::adr, ftpvoip) packet ⇒ bool **where**
 FTPVOIP-is-setup i port p = (id p = i ∧ content p = Setup port)

We need also an operator *ports-open* to get access to the two dynamic ports.

definition

FTPVOIP-ports-open :: id ⇒ port × port ⇒ (adr_{ip}, ftpvoip) history ⇒ bool **where**
 FTPVOIP-ports-open i p L = ((not-before (FTPVOIP-is-fin i) (FTPVOIP-is-setup i (fst p))
 L) ∧
 not-before (FTPVOIP-is-fin i) (FTPVOIP-is-connect i (snd p)) L)

As we do not know which entity closes the connection, we define an operator which checks if the closer is the caller.

fun

FTPVOIP-src-is-initiator :: id ⇒ adr_{ip} ⇒ (adr_{ip},ftpvoip) history ⇒ bool **where**
 FTPVOIP-src-is-initiator i a [] = False
 |FTPVOIP-src-is-initiator i a (p#S) = (((id p = i) ∧

$$\begin{aligned}
& (\exists \text{ port. content } p = \text{Setup port}) \wedge \\
& ((\text{fst } (\text{src } p) = \text{fst } a)) \vee \\
& (\text{FTPVOIP-src-is-initiator } i \ a \ S))
\end{aligned}$$

definition *FTPVOIP-subnet-of-adr* :: *int* \Rightarrow *adr_{ip}* *net* **where**
FTPVOIP-subnet-of-adr *x* = $\{\{(a,b). \ a = x\}\}$

fun *FTPVOIP-VOIP-STA* ::
 ((*adr_{ip}*, *ftpvoip*) *history*, *adr_{ip}*, *ftpvoip*) *FWStateTransition*
where

$$\begin{aligned}
& \text{FTPVOIP-VOIP-STA } ((a,c,d,ARQ), (InL, policy)) = \\
& \quad \text{Some } (((a,c,d, ARQ)\#InL, \\
& \quad (\text{allow-from-to-port } (1719::\text{port})(\text{subnet-of } d) (\text{subnet-of } c)) \oplus policy))
\end{aligned}$$

$$\begin{aligned}
& | \text{FTPVOIP-VOIP-STA } ((a,c,d,ARJ), (InL, policy)) = \\
& \quad (\text{if } (\text{not-before } (\text{FTPVOIP-is-fin } a) (\text{FTPVOIP-is-arq } a) \ InL) \\
& \quad \quad \text{then Some } (((a,c,d,ARJ)\#InL, \\
& \quad \quad \text{deny-from-to-port } (14::\text{port}) (\text{subnet-of } c) (\text{subnet-of } d) \oplus policy)) \\
& \quad \quad \text{else Some } (((a,c,d,ARJ)\#InL,policy)))
\end{aligned}$$

$$\begin{aligned}
& | \text{FTPVOIP-VOIP-STA } ((a,c,d,ACF \text{ callee}), (InL, policy)) = \\
& \quad \text{Some } (((a,c,d,ACF \text{ callee})\#InL, \\
& \quad \text{allow-from-to-port } (1720::\text{port}) (\text{subnet-of-adr } \text{callee}) (\text{subnet-of } d) \oplus \\
& \quad \text{allow-from-to-port } (1720::\text{port}) (\text{subnet-of } d) (\text{subnet-of-adr } \text{callee}) \oplus \\
& \quad \text{deny-from-to-port } (1719::\text{port}) (\text{subnet-of } d) (\text{subnet-of } c) \oplus \\
& \quad policy))
\end{aligned}$$

$$\begin{aligned}
& | \text{FTPVOIP-VOIP-STA } ((a,c,d, \text{Setup port}), (InL, policy)) = \\
& \quad \text{Some } (((a,c,d,\text{Setup port})\#InL, \\
& \quad \text{allow-from-to-port port } (\text{subnet-of } d) (\text{subnet-of } c) \oplus policy))
\end{aligned}$$

$$\begin{aligned}
& | \text{FTPVOIP-VOIP-STA } ((a,c,d, \text{ftpvoip.Connect port}), (InL, policy)) = \\
& \quad \text{Some } (((a,c,d,\text{ftpvoip.Connect port})\#InL, \\
& \quad \text{allow-from-to-port port } (\text{subnet-of } d) (\text{subnet-of } c) \oplus policy))
\end{aligned}$$

$$\begin{aligned}
& | \text{FTPVOIP-VOIP-STA } ((a,c,d,Fin), (InL,policy)) = \\
& \quad (\text{if } \exists \ p1 \ p2. \ \text{FTPVOIP-ports-open } a \ (p1,p2) \ InL \ \text{then } (\\
& \quad \quad (\text{if } \text{FTPVOIP-src-is-initiator } a \ c \ InL \\
& \quad \quad \text{then } (\text{Some } (((a,c,d,Fin)\#InL, \\
& \quad \quad (\text{deny-from-to-port } (1720::\text{int}) (\text{subnet-of } c) (\text{subnet-of } d)) \oplus \\
& \quad \quad (\text{deny-from-to-port } (\text{snd } (\text{SOME } p. \ \text{FTPVOIP-ports-open } a \ p \ InL)) \\
& \quad \quad (\text{subnet-of } c) (\text{subnet-of } d)) \oplus
\end{aligned}$$

$(deny-from-to-port (fst (SOME p. FTPVOIP-ports-open a p InL))$
 $(subnet-of d) (subnet-of c)) \oplus policy)))$

 $else (Some (((a,c,d,Fin)\#InL,$
 $(deny-from-to-port (1720::int) (subnet-of c) (subnet-of d)) \oplus$
 $(deny-from-to-port (fst (SOME p. FTPVOIP-ports-open a p InL))$
 $(subnet-of c) (subnet-of d)) \oplus$
 $(deny-from-to-port (snd (SOME p. FTPVOIP-ports-open a p InL))$
 $(subnet-of d) (subnet-of c)) \oplus policy))))))$

 $else$
 $(Some (((a,c,d,Fin)\#InL,policy))))$

$| FTPVOIP-VOIP-STA (p, (InL, policy)) =$
 $Some ((p\#InL,policy))$

fun *FTPVOIP-VOIP-STD* ::
 $((adr_{ip}, ftpvoip) history, adr_{ip}, ftpvoip) FWStateTransition$
where
 $FTPVOIP-VOIP-STD (p,s) = Some s$

definition *FTP-VOIP-STA* :: $((adr_{ip}, ftpvoip) history, adr_{ip}, ftpvoip) FWStateTransition$
where
 $FTP-VOIP-STA = (\lambda(x,x). Some x) o-m ((FTPVOIP-FTP-STA \otimes_S FTPVOIP-VOIP-STA$
 $o (\lambda (p,x). (p,x,x))))$

definition *FTP-VOIP-STD* :: $((adr_{ip}, ftpvoip) history, adr_{ip}, ftpvoip) FWStateTransition$
where
 $FTP-VOIP-STD = (\lambda(x,x). Some x) o-m ((FTPVOIP-FTP-STD \otimes_S FTPVOIP-VOIP-STD$
 $o (\lambda (p,x). (p,x,x))))$

definition *FTPVOIP-TRPolicy* **where**
 $FTPVOIP-TRPolicy = policy2MON ($
 $((FTP-VOIP-STA,FTP-VOIP-STD) \otimes_{\nabla} applyPolicy) o (\lambda (x,(y,z)). ((x,z),(x,(y,z))))))$

lemmas *FTPVOIP-ST-simps* = *Let-def in-subnet-def src-def dest-def*
subnet-of-def id-def FTPVOIP-port-open-def
FTPVOIP-is-init-def FTPVOIP-is-data-def FTPVOIP-is-port-request-def FTPVOIP-is-close-def
p-accept-def content-def PortCombinators exI
NetworkCore.id-def adr_{ip}Lemmas

datatype *ftp-states2* = *FS0* | *FS1* | *FS2* | *FS3*

datatype *voip-states2* = *V0* | *V1* | *V2* | *V3* | *V4* | *V5*

The constant *is-voip* checks if a trace corresponds to a legal VoIP protocol, given the IP-addresses of the three entities, the ID, and the two dynamic ports.

fun *FTPVOIP-is-voip* :: *voip-states2* \Rightarrow *address* \Rightarrow *address* \Rightarrow *address* \Rightarrow *id* \Rightarrow *port* \Rightarrow
port \Rightarrow (*adr_{ip}*, *ftpvoip*) *history* \Rightarrow *bool*

where

FTPVOIP-is-voip H s d g i p1 p2 [] = (*H* = *V5*)
|FTPVOIP-is-voip H s d g i p1 p2 (x#InL) =
 (((λ (*id*,*sr*,*de*,*co*).
 (((*id* = *i* \wedge
 (*H* = *V4* \wedge ((*sr* = (*s*,1719) \wedge *de* = (*g*,1719) \wedge *co* = *ARQ* \wedge
FTPVOIP-is-voip V5 s d g i p1 p2 InL))) \vee
 (*H* = *V0* \wedge *sr* = (*g*,1719) \wedge *de* = (*s*,1719) \wedge *co* = *ARJ* \wedge
FTPVOIP-is-voip V4 s d g i p1 p2 InL) \vee
 (*H* = *V3* \wedge *sr* = (*g*,1719) \wedge *de* = (*s*,1719) \wedge *co* = *ACF d* \wedge
FTPVOIP-is-voip V4 s d g i p1 p2 InL) \vee
 (*H* = *V2* \wedge *sr* = (*s*,1720) \wedge *de* = (*d*,1720) \wedge *co* = *Setup p1* \wedge
FTPVOIP-is-voip V3 s d g i p1 p2 InL) \vee
 (*H* = *V1* \wedge *sr* = (*d*,1720) \wedge *de* = (*s*,1720) \wedge *co* = *Connect p2* \wedge
FTPVOIP-is-voip V2 s d g i p1 p2 InL) \vee
 (*H* = *V1* \wedge *sr* = (*s*,*p1*) \wedge *de* = (*d*,*p2*) \wedge *co* = *Stream* \wedge
FTPVOIP-is-voip V1 s d g i p1 p2 InL) \vee
 (*H* = *V1* \wedge *sr* = (*d*,*p2*) \wedge *de* = (*s*,*p1*) \wedge *co* = *Stream* \wedge
FTPVOIP-is-voip V1 s d g i p1 p2 InL) \vee
 (*H* = *V0* \wedge *sr* = (*d*,1720) \wedge *de* = (*s*,1720) \wedge *co* = *Fin* \wedge
FTPVOIP-is-voip V1 s d g i p1 p2 InL) \vee
 (*H* = *V0* \wedge *sr* = (*s*,1720) \wedge *de* = (*d*,1720) \wedge *co* = *Fin* \wedge
FTPVOIP-is-voip V1 s d g i p1 p2 InL)))))) *x*)

Finally, *NB-voip* returns the set of protocol traces which correspond to a correct protocol run given the three addresses, the ID, and the two dynamic ports.

definition

FTPVOIP-NB-voip :: *address* \Rightarrow *address* \Rightarrow *address* \Rightarrow *id* \Rightarrow *port* \Rightarrow *port* \Rightarrow
 (*adr_{ip}*, *ftpvoip*) *history* *set* **where**
FTPVOIP-NB-voip s d g i p1 p2 = {*x*. (*FTPVOIP-is-voip V0 s d g i p1 p2 x*)}

$$\begin{aligned}
& \text{FTPVOIP-is-ftp} :: \text{ftp-states2} \Rightarrow \text{adr}_{ip} \Rightarrow \text{adr}_{ip} \Rightarrow \text{id} \Rightarrow \text{port} \Rightarrow \\
& (\text{adr}_{ip}, \text{ftpvoip}) \text{ history} \Rightarrow \text{bool} \\
\text{where} \\
& \text{FTPVOIP-is-ftp } H \text{ c s i p } [] = (H = \text{FS3}) \\
& |\text{FTPVOIP-is-ftp } H \text{ c s i p } (x \# \text{InL}) = (\text{snd } s = 21 \wedge ((\lambda (id, sr, de, co). (((id = i \wedge (\\
& (H = \text{FS2} \wedge sr = c \wedge de = s \wedge co = \text{ftp-init} \wedge \text{FTPVOIP-is-ftp } \text{FS3 } c \text{ s i p } \text{InL}) \vee \\
& (H = \text{FS1} \wedge sr = c \wedge de = s \wedge co = \text{ftp-port-request } p \wedge \text{FTPVOIP-is-ftp } \text{FS2 } c \text{ s i p } \text{InL}) \\
& \vee \\
& (H = \text{FS1} \wedge sr = s \wedge de = (\text{fst } c, p) \wedge co = \text{ftp-data} \wedge \text{FTPVOIP-is-ftp } \text{FS1 } c \text{ s i p } \text{InL}) \vee \\
& (H = \text{FS0} \wedge sr = c \wedge de = s \wedge co = \text{ftp-close} \wedge \text{FTPVOIP-is-ftp } \text{FS1 } c \text{ s i p } \text{InL})))))) x))
\end{aligned}$$
$$\begin{aligned} &FTPVOIP-NB-ftp :: adr_{ip} \text{ src} \Rightarrow adr_{ip} \text{ dest} \Rightarrow id \Rightarrow port \Rightarrow \\ &\quad (adr_{ip}, ftpvoip) \text{ history set } \mathbf{where} \\ &FTPVOIP-NB-ftp \text{ s d i p} = \{x. (FTPVOIP-is-ftp \text{ FS0 s d i p } x)\} \end{aligned}$$
$$\begin{aligned} \text{ftp-voip-interleaved} :: & \text{adr}_{ip} \text{ src} \Rightarrow \text{adr}_{ip} \text{ dest} \Rightarrow \text{id} \Rightarrow \text{port} \Rightarrow \\ & \text{address} \Rightarrow \text{address} \Rightarrow \text{address} \Rightarrow \text{id} \Rightarrow \text{port} \Rightarrow \text{port} \Rightarrow \\ & (\text{adr}_{ip}, \text{ftpvoip}) \text{ history set} \end{aligned}$$
$$\begin{aligned} &ftp\text{-}voip\text{-}interleaved\ s1\ d1\ i1\ p1\ vs\ vd\ vg\ vi\ vp1\ vp2 = \\ &\{x. (FTPVOIP\text{-}is\text{-}ftp\ FS0\ s1\ d1\ i1\ p1\ (packet\text{-}with\text{-}id\ x\ i1)) \wedge \\ &\quad (FTPVOIP\text{-}is\text{-}voip\ V0\ vs\ vd\ vg\ vi\ vp1\ vp2\ (packet\text{-}with\text{-}id\ x\ vi))\} \end{aligned}$$

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A. Appendix

theory *NormalisationGenericProofs*
imports *FWNormalisationCore*
begin

This theory contains the generic proofs of the normalisation procedure, i.e. those which are independent from the concrete semantical interpretation function.

lemma *domNMT*: $\text{dom } X \neq \{\} \implies X \neq \emptyset$
 $\langle \text{proof} \rangle$

lemma *denyNMT*: $\text{deny-all} \neq \emptyset$
 $\langle \text{proof} \rangle$

lemma *wellformed-policy1-charn*[*rule-format*] : $\text{wellformed-policy1 } p \longrightarrow \text{DenyAll} \in \text{set } p \longrightarrow (\exists p'. p = \text{DenyAll} \# p' \wedge \text{DenyAll} \notin \text{set } p')$
 $\langle \text{proof} \rangle$

lemma *singleCombinatorsConc*: $\text{singleCombinators } (x \# xs) \implies \text{singleCombinators } xs$
 $\langle \text{proof} \rangle$

lemma *aux0-0*: $\text{singleCombinators } x \implies \neg (\exists a b. (a \oplus b) \in \text{set } x)$
 $\langle \text{proof} \rangle$

lemma *aux0-4*: $(a \in \text{set } x \vee a \in \text{set } y) = (a \in \text{set } (x @ y))$
 $\langle \text{proof} \rangle$

lemma *aux0-1*: $\llbracket \text{singleCombinators } xs; \text{singleCombinators } [x] \rrbracket \implies \text{singleCombinators } (x \# xs)$
 $\langle \text{proof} \rangle$

lemma *aux0-6*: $\llbracket \text{singleCombinators } xs; \neg (\exists a b. x = a \oplus b) \rrbracket \implies \text{singleCombinators } (x \# xs)$
 $\langle \text{proof} \rangle$

lemma *aux0-5*: $\neg (\exists a b. (a \oplus b) \in \text{set } x) \implies \text{singleCombinators } x$
 $\langle \text{proof} \rangle$

lemma *ANDConc*[*rule-format*]: $\text{allNetsDistinct } (a \# p) \longrightarrow \text{allNetsDistinct } (p)$
 $\langle \text{proof} \rangle$

lemma *aux6*: *twoNetsDistinct a1 a2 a b* \implies
 $\text{dom } (\text{deny-all-from-to } a1 \ a2) \cap \text{dom } (\text{deny-all-from-to } a \ b) = \{\}$
 $\langle \text{proof} \rangle$

lemma *aux5*[*rule-format*]: $(\text{DenyAllFromTo } a \ b) \in \text{set } p \longrightarrow a \in \text{set } (\text{net-list } p)$
 $\langle \text{proof} \rangle$

lemma *aux5a*[*rule-format*]: $(\text{DenyAllFromTo } b \ a) \in \text{set } p \longrightarrow a \in \text{set } (\text{net-list } p)$
 $\langle \text{proof} \rangle$

lemma *aux5c*[*rule-format*]:
 $(\text{AllowPortFromTo } a \ b \ po) \in \text{set } p \longrightarrow a \in \text{set } (\text{net-list } p)$
 $\langle \text{proof} \rangle$

lemma *aux5d*[*rule-format*]:
 $(\text{AllowPortFromTo } b \ a \ po) \in \text{set } p \longrightarrow a \in \text{set } (\text{net-list } p)$
 $\langle \text{proof} \rangle$

lemma *aux10*[*rule-format*]: $a \in \text{set } (\text{net-list } p) \longrightarrow a \in \text{set } (\text{net-list-aux } p)$
 $\langle \text{proof} \rangle$

lemma *srcInNetListaux*[*simp*]: $\llbracket x \in \text{set } p; \text{singleCombinators}[x]; x \neq \text{DenyAll} \rrbracket \implies$
 $\text{srcNet } x \in \text{set } (\text{net-list-aux } p)$
 $\langle \text{proof} \rangle$

lemma *destInNetListaux*[*simp*]: $\llbracket x \in \text{set } p; \text{singleCombinators}[x]; x \neq \text{DenyAll} \rrbracket \implies$
 $\text{destNet } x \in \text{set } (\text{net-list-aux } p)$
 $\langle \text{proof} \rangle$

lemma *tND1*: $\llbracket \text{allNetsDistinct } p; x \in \text{set } p; y \in \text{set } p; a = \text{srcNet } x;$
 $b = \text{destNet } x; c = \text{srcNet } y; d = \text{destNet } y; a \neq c;$
 $\text{singleCombinators}[x]; x \neq \text{DenyAll}; \text{singleCombinators}[y];$
 $y \neq \text{DenyAll} \rrbracket \implies \text{twoNetsDistinct } a \ b \ c \ d$
 $\langle \text{proof} \rangle$

lemma *tND2*: $\llbracket \text{allNetsDistinct } p; x \in \text{set } p; y \in \text{set } p; a = \text{srcNet } x;$
 $b = \text{destNet } x; c = \text{srcNet } y; d = \text{destNet } y; b \neq d;$
 $\text{singleCombinators}[x]; x \neq \text{DenyAll}; \text{singleCombinators}[y];$
 $y \neq \text{DenyAll} \rrbracket \implies \text{twoNetsDistinct } a \ b \ c \ d$
 $\langle \text{proof} \rangle$

lemma *tND*: $\llbracket \text{allNetsDistinct } p; x \in \text{set } p; y \in \text{set } p; a = \text{srcNet } x;$
 $b = \text{destNet } x; c = \text{srcNet } y; d = \text{destNet } y; a \neq c \vee b \neq d;$
 $\text{singleCombinators}[x]; x \neq \text{DenyAll}; \text{singleCombinators}[y]; y \neq \text{DenyAll} \rrbracket$
 $\implies \text{twoNetsDistinct } a \ b \ c \ d$
 $\langle \text{proof} \rangle$

lemma *aux7*: $\llbracket \text{DenyAllFromTo } a \ b \in \text{set } p; \text{allNetsDistinct } ((\text{DenyAllFromTo } c \ d) \# p);$
 $a \neq c \vee b \neq d \rrbracket \implies \text{twoNetsDistinct } a \ b \ c \ d$
 $\langle \text{proof} \rangle$

lemma *aux7a*: $\llbracket \text{DenyAllFromTo } a \ b \in \text{set } p;$
 $\text{allNetsDistinct } ((\text{AllowPortFromTo } c \ d \ po) \# p); a \neq c \vee b \neq d \rrbracket \implies$
 $\text{twoNetsDistinct } a \ b \ c \ d$
 $\langle \text{proof} \rangle$

lemma *nDComm*: **assumes** *ab*: *netsDistinct a b* **shows** *ba*: *netsDistinct b a*
 $\langle \text{proof} \rangle$

lemma *tNDComm*:
assumes *abcd*: *twoNetsDistinct a b c d* **shows** *twoNetsDistinct c d a b*
 $\langle \text{proof} \rangle$

lemma *aux[rule-format]*: $a \in \text{set } (\text{removeShadowRules2 } p) \longrightarrow a \in \text{set } p$
 $\langle \text{proof} \rangle$

lemma *aux12*: $\llbracket a \in x; b \notin x \rrbracket \implies a \neq b$
 $\langle \text{proof} \rangle$

lemma *ND0aux1[rule-format]*: $\text{DenyAllFromTo } x \ y \in \text{set } b \implies$
 $x \in \text{set } (\text{net-list-aux } b)$
 $\langle \text{proof} \rangle$

lemma *ND0aux2[rule-format]*: $\text{DenyAllFromTo } x \ y \in \text{set } b \implies$
 $y \in \text{set } (\text{net-list-aux } b)$
 $\langle \text{proof} \rangle$

lemma *ND0aux3[rule-format]*: $\text{AllowPortFromTo } x \ y \ p \in \text{set } b \implies$
 $x \in \text{set } (\text{net-list-aux } b)$
 $\langle \text{proof} \rangle$

lemma *ND0aux4[rule-format]*: $\text{AllowPortFromTo } x \ y \ p \in \text{set } b \implies$
 $y \in \text{set } (\text{net-list-aux } b)$
 $\langle \text{proof} \rangle$

lemma *aNDSubsetaux*[*rule-format*]: *singleCombinators a* \longrightarrow *set a* \subseteq *set b* \longrightarrow
 $\text{set } (\text{net-list-aux } a) \subseteq \text{set } (\text{net-list-aux } b)$
 <proof>

lemma *aNDSetsEqaux*[*rule-format*]: *singleCombinators a* \longrightarrow *singleCombinators b* \longrightarrow
 $\text{set } a = \text{set } b \longrightarrow \text{set } (\text{net-list-aux } a) = \text{set } (\text{net-list-aux } b)$
 <proof>

lemma *aNDSubset*: $\llbracket \text{singleCombinators } a; \text{set } a \subseteq \text{set } b; \text{allNetsDistinct } b \rrbracket \Longrightarrow$
 $\text{allNetsDistinct } a$
 <proof>

lemma *aNDSetsEq*: $\llbracket \text{singleCombinators } a; \text{singleCombinators } b; \text{set } a = \text{set } b; \text{allNetsDistinct } b \rrbracket \Longrightarrow$
 $\text{allNetsDistinct } a$
 <proof>

lemma *SCConca*: $\llbracket \text{singleCombinators } p; \text{singleCombinators } [a] \rrbracket \Longrightarrow$
 $\text{singleCombinators } (a \# p)$
 <proof>

lemma *aux3*[*simp*]: $\llbracket \text{singleCombinators } p; \text{singleCombinators } [a]; \text{allNetsDistinct } (a \# p) \rrbracket \Longrightarrow$
 $\text{allNetsDistinct } (a \# a \# p)$
 <proof>

lemma *wp1-aux1a*[*rule-format*]: $xs \neq [] \longrightarrow \text{wellformed-policy1-strong } (xs @ [x]) \longrightarrow$
 $\text{wellformed-policy1-strong } xs$
 <proof>

lemma *wp1alternative-RS1*[*rule-format*]: $\text{DenyAll} \in \text{set } p \longrightarrow$
 $\text{wellformed-policy1-strong } (\text{removeShadowRules1 } p)$
 <proof>

lemma *wellformed-eq*: $\text{DenyAll} \in \text{set } p \longrightarrow$
 $((\text{wellformed-policy1 } p) = (\text{wellformed-policy1-strong } p))$
 <proof>

lemma *set-insort*: $\text{set}(\text{insort } x \text{ } xs \text{ } l) = \text{insert } x (\text{set } xs)$
 <proof>

lemma *set-sort*[*simp*]: $\text{set}(\text{sort } xs \text{ } l) = \text{set } xs$
 <proof>

lemma *set-sortQ*: $\text{set}(\text{qsort } xs \text{ } l) = \text{set } xs$
 <proof>

lemma *aux79*[rule-format]: $y \in \text{set } (\text{insort } x \ a \ l) \longrightarrow y \neq x \longrightarrow y \in \text{set } a$
 <proof>

lemma *aux80*: $\llbracket y \notin \text{set } p; y \neq x \rrbracket \implies y \notin \text{set } (\text{insort } x \ (\text{sort } p \ l) \ l)$
 <proof>

lemma *WP1Conca*: $\text{DenyAll} \notin \text{set } p \implies \text{wellformed-policy1 } (a \# p)$
 <proof>

lemma *saux*[simp]: $(\text{insort } \text{DenyAll } p \ l) = \text{DenyAll} \# p$
 <proof>

lemma *saux3*[rule-format]: $\text{DenyAllFromTo } a \ b \in \text{set list} \longrightarrow$
 $\text{DenyAllFromTo } c \ d \notin \text{set list} \longrightarrow (a \neq c) \vee (b \neq d)$
 <proof>

lemma *waux2*[rule-format]: $(\text{DenyAll} \notin \text{set } xs) \longrightarrow \text{wellformed-policy1 } xs$
 <proof>

lemma *waux3*[rule-format]: $\llbracket x \neq a; x \notin \text{set } p \rrbracket \implies x \notin \text{set } (\text{insort } a \ p \ l)$
 <proof>

lemma *wellformed1-sorted-aux*[rule-format]: $\text{wellformed-policy1 } (x \# p) \implies$
 $\text{wellformed-policy1 } (\text{insort } x \ p \ l)$
 <proof>

lemma *wellformed1-sorted-auxQ*[rule-format]: $\text{wellformed-policy1 } (p) \implies$
 $\text{wellformed-policy1 } (q\text{sort } p \ l)$
 <proof>

lemma *SR1Subset*: $\text{set } (\text{removeShadowRules1 } p) \subseteq \text{set } p$
 <proof>

lemma *SCSubset*[rule-format]: $\text{singleCombinators } b \longrightarrow \text{set } a \subseteq \text{set } b \longrightarrow$
 $\text{singleCombinators } a$
 <proof>

lemma *setInsert*[simp]: $\text{set list} \subseteq \text{insert } a \ (\text{set list})$
 <proof>

lemma *SC-RS1*[rule-format,simp]: $\text{singleCombinators } p \longrightarrow \text{allNetsDistinct } p \longrightarrow$
 $\text{singleCombinators } (\text{removeShadowRules1 } p)$
 <proof>

lemma *RS2Set*[rule-format]: $\text{set } (\text{removeShadowRules2 } p) \subseteq \text{set } p$

$\langle \text{proof} \rangle$

lemma *WP1*: $a \notin \text{set } \text{list} \implies a \notin \text{set } (\text{removeShadowRules2 } \text{list})$

$\langle \text{proof} \rangle$

lemma *denyAllDom[simp]*: $x \in \text{dom } (\text{deny-all})$

$\langle \text{proof} \rangle$

lemma *lCdom2*: $(\text{list2FWpolicy } (a @ (b @ c))) = (\text{list2FWpolicy } ((a @ b) @ c))$

$\langle \text{proof} \rangle$

lemma *SCConcEnd*: $\text{singleCombinators } (xs @ [xa]) \implies \text{singleCombinators } xs$

$\langle \text{proof} \rangle$

lemma *list2FWpolicyconc[rule-format]*: $a \neq [] \longrightarrow$

$(\text{list2FWpolicy } (xa \# a)) = (xa) \oplus (\text{list2FWpolicy } a)$

$\langle \text{proof} \rangle$

lemma *wp1n-tl [rule-format]*: $\text{wellformed-policy1-strong } p \longrightarrow$

$p = (\text{DenyAll} \# (\text{tl } p))$

$\langle \text{proof} \rangle$

lemma *foo2*:

$\llbracket a \notin \text{set } ps; a \notin \text{set } ss; \text{set } p = \text{set } s; p = (a \# (ps)); s = (a \# ss) \rrbracket \implies$

$\text{set } (ps) = \text{set } (ss)$

$\langle \text{proof} \rangle$

lemma *SCnotConc[rule-format,simp]*: $a \oplus b \in \text{set } p \longrightarrow \text{singleCombinators } p \longrightarrow \text{False}$

$\langle \text{proof} \rangle$

lemma *auxx8*: $\text{removeShadowRules1-alternative-rev } [x] = [x]$

$\langle \text{proof} \rangle$

lemma *RS1End[rule-format]*: $x \neq \text{DenyAll} \longrightarrow \text{removeShadowRules1 } (xs @ [x]) =$

$(\text{removeShadowRules1 } xs) @ [x]$

$\langle \text{proof} \rangle$

lemma *aux114*: $x \neq \text{DenyAll} \implies \text{removeShadowRules1-alternative-rev } (x \# xs) =$

$x \# (\text{removeShadowRules1-alternative-rev } xs)$

$\langle \text{proof} \rangle$

lemma *aux115[rule-format]*: $x \neq \text{DenyAll} \implies \text{removeShadowRules1-alternative } (xs @ [x])$

$= (\text{removeShadowRules1-alternative } xs) @ [x]$

$\langle \text{proof} \rangle$

lemma *RS1-DA[simp]*: $\text{removeShadowRules1 } (xs @ [\text{DenyAll}]) = [\text{DenyAll}]$

<proof>

lemma *rSR1-eq*: *removeShadowRules1-alternative* = *removeShadowRules1*

<proof>

lemma *domInterMT*[*rule-format*]: $\llbracket \text{dom } a \cap \text{dom } b = \{\}; x \in \text{dom } a \rrbracket \implies x \notin \text{dom } b$

<proof>

lemma *domComm*: $\text{dom } a \cap \text{dom } b = \text{dom } b \cap \text{dom } a$

<proof>

lemma *r-not-DA-in-tl*[*rule-format*]: $\text{wellformed-policy1-strong } p \longrightarrow a \in \text{set } p \longrightarrow$
 $a \neq \text{DenyAll} \longrightarrow a \in \text{set } (\text{tl } p)$

<proof>

lemma *wp1-aux1aa*[*rule-format*]: $\text{wellformed-policy1-strong } p \longrightarrow \text{DenyAll} \in \text{set } p$

<proof>

lemma *mauxa*: $(\exists r. a \ b = \lfloor r \rfloor) = (a \ b \neq \perp)$

<proof>

lemma *l2p-aux*[*rule-format*]: $\text{list} \neq [] \longrightarrow$
 $\text{list2FWpolicy } (a \ \# \ \text{list}) = a \oplus (\text{list2FWpolicy } \text{list})$

<proof>

lemma *l2p-aux2*[*rule-format*]: $\text{list} = [] \implies \text{list2FWpolicy } (a \ \# \ \text{list}) = a$

<proof>

lemma *aux7aa*: $\llbracket \text{AllowPortFromTo } a \ b \ po \in \text{set } p;$
 $\text{allNetsDistinct } ((\text{AllowPortFromTo } c \ d \ po) \ \# \ p); a \neq c \vee b \neq d \rrbracket \implies$
 $\text{twoNetsDistinct } a \ b \ c \ d$

<proof>

lemma *ANDConcEnd*: $\llbracket \text{allNetsDistinct } (xs \ @ \ [xa]); \text{singleCombinators } xs \rrbracket \implies$
 $\text{allNetsDistinct } xs$

<proof>

lemma *WP1ConcEnd*[*rule-format*]:
 $\text{wellformed-policy1 } (xs @ [xa]) \longrightarrow \text{wellformed-policy1 } xs$

<proof>

lemma *NDComm*: $\text{netsDistinct } a \ b = \text{netsDistinct } b \ a$

<proof>

lemma *wellformed1-sorted*[*simp*]:
assumes *wp1*: *wellformed-policy1* *p*
shows *wellformed-policy1* (*sort* *p* *l*)

$\langle proof \rangle$

lemma *wellformed1-sortedQ[simp]*:
 assumes *wp1: wellformed-policy1 p*
 shows *wellformed-policy1 (qsort p l)*
 $\langle proof \rangle$

lemma *SC1[simp]*: *singleCombinators p \implies singleCombinators (removeShadowRules1 p)*
 $\langle proof \rangle$

lemma *SC2[simp]*: *singleCombinators p \implies singleCombinators (removeShadowRules2 p)*
 $\langle proof \rangle$

lemma *SC3[simp]*: *singleCombinators p \implies singleCombinators (sort p l)*
 $\langle proof \rangle$

lemma *SC3Q[simp]*: *singleCombinators p \implies singleCombinators (qsort p l)*
 $\langle proof \rangle$

lemma *aND-RS1[simp]*: $\llbracket \text{singleCombinators } p; \text{allNetsDistinct } p \rrbracket \implies$
 allNetsDistinct (removeShadowRules1 p)
 $\langle proof \rangle$

lemma *aND-RS2[simp]*: $\llbracket \text{singleCombinators } p; \text{allNetsDistinct } p \rrbracket \implies$
 allNetsDistinct (removeShadowRules2 p)
 $\langle proof \rangle$

lemma *aND-sort[simp]*: $\llbracket \text{singleCombinators } p; \text{allNetsDistinct } p \rrbracket \implies$
 allNetsDistinct (sort p l)
 $\langle proof \rangle$

lemma *aND-sortQ[simp]*: $\llbracket \text{singleCombinators } p; \text{allNetsDistinct } p \rrbracket \implies$
 allNetsDistinct (qsort p l)
 $\langle proof \rangle$

lemma *inRS2[rule-format,simp]*: $x \notin \text{set } p \longrightarrow x \notin \text{set (removeShadowRules2 } p)$
 $\langle proof \rangle$

lemma *distinct-RS2*[*rule-format,simp*]: $\text{distinct } p \longrightarrow$
 $\text{distinct } (\text{removeShadowRules2 } p)$
 ⟨*proof*⟩

lemma *setPaireq*: $\{x, y\} = \{a, b\} \implies x = a \wedge y = b \vee x = b \wedge y = a$
 ⟨*proof*⟩

lemma *position-positive*[*rule-format*]: $a \in \text{set } l \longrightarrow \text{position } a \ l > 0$
 ⟨*proof*⟩

lemma *pos-noteq*[*rule-format*]:
 $a \in \text{set } l \longrightarrow b \in \text{set } l \longrightarrow c \in \text{set } l \longrightarrow a \neq b \longrightarrow$
 $(\text{position } a \ l) <= (\text{position } b \ l) \longrightarrow$
 $(\text{position } b \ l) <= (\text{position } c \ l) \longrightarrow$
 $a \neq c$
 ⟨*proof*⟩

lemma *setPair-noteq*: $\{a, b\} \neq \{c, d\} \implies \neg ((a = c) \wedge (b = d))$
 ⟨*proof*⟩

lemma *setPair-noteq-allow*: $\{a, b\} \neq \{c, d\} \implies \neg ((a = c) \wedge (b = d) \wedge P)$
 ⟨*proof*⟩

lemma *order-trans*:
 $\llbracket \text{in-list } x \ l; \text{in-list } y \ l; \text{in-list } z \ l; \text{singleCombinators } [x];$
 $\text{singleCombinators } [y]; \text{singleCombinators } [z]; \text{smaller } x \ y \ l; \text{smaller } y \ z \ l \rrbracket \implies$
 $\text{smaller } x \ z \ l$
 ⟨*proof*⟩

lemma *sortedConcStart*[*rule-format*]:
 $\text{sorted } (a \ \# \ aa \ \# \ p) \ l \longrightarrow \text{in-list } a \ l \longrightarrow \text{in-list } aa \ l \longrightarrow \text{all-in-list } p \ l \longrightarrow$
 $\text{singleCombinators } [a] \longrightarrow \text{singleCombinators } [aa] \longrightarrow \text{singleCombinators } p \longrightarrow$
 $\text{sorted } (a \ \# \ p) \ l$
 ⟨*proof*⟩

lemma *singleCombinatorsStart*[*simp*]: $\text{singleCombinators } (x \ \# \ xs) \implies$
 $\text{singleCombinators } [x]$
 ⟨*proof*⟩

lemma *sorted-is-smaller*[*rule-format*]:
 $\text{sorted } (a \ \# \ p) \ l \longrightarrow \text{in-list } a \ l \longrightarrow \text{in-list } b \ l \longrightarrow \text{all-in-list } p \ l \longrightarrow$
 $\text{singleCombinators } [a] \longrightarrow \text{singleCombinators } p \longrightarrow b \in \text{set } p \longrightarrow \text{smaller } a \ b \ l$
 ⟨*proof*⟩

lemma *sortedConcEnd*[*rule-format*]: $\text{sorted } (a \ \# \ p) \ l \longrightarrow \text{in-list } a \ l \longrightarrow$
 $\text{all-in-list } p \ l \longrightarrow \text{singleCombinators } [a] \longrightarrow$
 $\text{singleCombinators } p \longrightarrow \text{sorted } p \ l$
 ⟨*proof*⟩

lemma *in-set-in-list*[rule-format]: $a \in \text{set } p \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{in-list } a \ l$
 <proof>

lemma *sorted-Consb*[rule-format]:
 $\text{all-in-list } (x\#xs) \ l \longrightarrow \text{singleCombinators } (x\#xs) \longrightarrow$
 $(\text{sorted } xs \ l \ \& \ (\text{ALL } y:\text{set } xs. \text{smaller } x \ y \ l)) \longrightarrow (\text{sorted } (x\#xs) \ l)$
 <proof>

lemma *sorted-Cons*: $\llbracket \text{all-in-list } (x\#xs) \ l; \text{singleCombinators } (x\#xs) \rrbracket \Longrightarrow$
 $(\text{sorted } xs \ l \ \& \ (\text{ALL } y:\text{set } xs. \text{smaller } x \ y \ l)) = (\text{sorted } (x\#xs) \ l)$
 <proof>

lemma *smaller-antisym*: $\llbracket \neg \text{smaller } a \ b \ l; \text{in-list } a \ l; \text{in-list } b \ l;$
 $\text{singleCombinators } [a]; \text{singleCombinators } [b] \rrbracket \Longrightarrow$
 $\text{smaller } b \ a \ l$
 <proof>

lemma *set-insort-insert*: $\text{set } (\text{insort } x \ xs \ l) \subseteq \text{insert } x \ (\text{set } xs)$
 <proof>

lemma *all-in-listSubset*[rule-format]: $\text{all-in-list } b \ l \longrightarrow \text{singleCombinators } a \longrightarrow$
 $\text{set } a \subseteq \text{set } b \longrightarrow \text{all-in-list } a \ l$
 <proof>

lemma *singleCombinators-insort*: $\llbracket \text{singleCombinators } [x]; \text{singleCombinators } xs \rrbracket \Longrightarrow$
 $\text{singleCombinators } (\text{insort } x \ xs \ l)$
 <proof>

lemma *all-in-list-insort*: $\llbracket \text{all-in-list } xs \ l; \text{singleCombinators } (x\#xs);$
 $\text{in-list } x \ l \rrbracket \Longrightarrow \text{all-in-list } (\text{insort } x \ xs \ l) \ l$
 <proof>

lemma *sorted-ConsA*: $\llbracket \text{all-in-list } (x\#xs) \ l; \text{singleCombinators } (x\#xs) \rrbracket \Longrightarrow$
 $(\text{sorted } (x\#xs) \ l) = (\text{sorted } xs \ l \ \& \ (\text{ALL } y:\text{set } xs. \text{smaller } x \ y \ l))$
 <proof>

lemma *is-in-insort*: $y \in \text{set } xs \Longrightarrow y \in \text{set } (\text{insort } x \ xs \ l)$
 <proof>

lemma *sorted-insorta*[rule-format]:
 $\text{sorted } (\text{insort } x \ xs \ l) \ l \longrightarrow \text{all-in-list } (x\#xs) \ l \longrightarrow \text{distinct } (x\#xs) \longrightarrow$
 $\text{singleCombinators } [x] \longrightarrow \text{singleCombinators } xs \longrightarrow \text{sorted } xs \ l$
 <proof>

lemma *sorted-insortb*[rule-format]:
 $\text{sorted } xs \ l \longrightarrow \text{all-in-list } (x\#xs) \ l \longrightarrow \text{distinct } (x\#xs) \longrightarrow$
 $\text{singleCombinators } [x] \longrightarrow \text{singleCombinators } xs \longrightarrow \text{sorted } (\text{insort } x \ xs \ l) \ l$
 <proof>

lemma *sorted-insort*: $\llbracket \text{all-in-list } (x\#xs) \ l; \text{distinct}(x\#xs); \text{singleCombinators } [x];$
 $\text{singleCombinators } xs \rrbracket \implies$
 $\text{sorted } (\text{insort } x \ xs \ l) \ l = \text{sorted } xs \ l$

<proof>

lemma *distinct-insort*: $\text{distinct } (\text{insort } x \ xs \ l) = (x \notin \text{set } xs \wedge \text{distinct } xs)$

<proof>

lemma *distinct-sort[simp]*: $\text{distinct } (\text{sort } xs \ l) = \text{distinct } xs$

<proof>

lemma *sort-is-sorted[rule-format]*: $\text{all-in-list } p \ l \longrightarrow \text{distinct } p \longrightarrow$
 $\text{singleCombinators } p \longrightarrow \text{sorted } (\text{sort } p \ l) \ l$

<proof>

lemma *smaller-sym[rule-format]*: $\text{all-in-list } [a] \ l \longrightarrow \text{smaller } a \ a \ l$

<proof>

lemma *SC-sublist[rule-format]*: $\text{singleCombinators } xs \implies$
 $\text{singleCombinators } (qsort \ [y \leftarrow xs. \ P \ y] \ l)$

<proof>

lemma *all-in-list-sublist[rule-format]*: $\text{singleCombinators } xs \longrightarrow \text{all-in-list } xs \ l \longrightarrow$
 $\text{all-in-list } (qsort \ [y \leftarrow xs. \ P \ y] \ l) \ l$

<proof>

lemma *SC-sublist2[rule-format]*: $\text{singleCombinators } xs \longrightarrow$
 $\text{singleCombinators } ([y \leftarrow xs. \ P \ y])$

<proof>

lemma *all-in-list-sublist2[rule-format]*: $\text{singleCombinators } xs \longrightarrow \text{all-in-list } xs \ l \longrightarrow$
 $\text{all-in-list } ([y \leftarrow xs. \ P \ y]) \ l$

<proof>

lemma *all-in-listAppend[rule-format]*:
 $\text{all-in-list } (xs) \ l \longrightarrow \text{all-in-list } (ys) \ l \longrightarrow \text{all-in-list } (xs \ @ \ ys) \ l$

<proof>

lemma *distinct-sortQ[rule-format]*: $\text{singleCombinators } xs \longrightarrow$

all-in-list xs l \longrightarrow *distinct xs* \longrightarrow *distinct (qsort xs l)*
 <proof>

lemma *singleCombinatorsAppend*[rule-format]:
singleCombinators xs \longrightarrow *singleCombinators ys* \longrightarrow *singleCombinators (xs @ ys)*
 <proof>

lemma *sorted-append1*[rule-format]:
all-in-list xs l \longrightarrow *singleCombinators xs* \longrightarrow
all-in-list ys l \longrightarrow
singleCombinators ys \longrightarrow
 (*sorted (xs@ys) l* \longrightarrow (*sorted xs l* & *sorted ys l* & ($\forall x \in \text{set } xs. \forall y \in \text{set } ys. \text{smaller } x \ y \ l$)))
 <proof>

lemma *sorted-append2*[rule-format]:
all-in-list xs l \longrightarrow *singleCombinators xs* \longrightarrow
all-in-list ys l \longrightarrow
singleCombinators ys \longrightarrow
 (*sorted xs l* & *sorted ys l* & ($\forall x \in \text{set } xs. \forall y \in \text{set } ys. \text{smaller } x \ y \ l$)) \longrightarrow (*sorted (xs@ys) l*)
 <proof>

lemma *sorted-append*[rule-format]:
all-in-list xs l \longrightarrow *singleCombinators xs* \longrightarrow
all-in-list ys l \longrightarrow
singleCombinators ys \longrightarrow (*sorted (xs@ys) l*) =
 (*sorted xs l* & *sorted ys l* & ($\forall x \in \text{set } xs. \forall y \in \text{set } ys. \text{smaller } x \ y \ l$))
 <proof>

lemma *sort-is-sortedQ*[rule-format]: *all-in-list p l* \longrightarrow
singleCombinators p \longrightarrow *sorted (qsort p l) l*
 <proof>

lemma *inSet-not-MT*: $a \in \text{set } p \implies p \neq []$
 <proof>

lemma *RS1n-assoc*: $x \neq \text{DenyAll} \implies \text{removeShadowRules1-alternative } xs \ @ \ [x] =$
removeShadowRules1-alternative (xs @ [x])
 <proof>

lemma *RS1n-nMT*[rule-format,simp]: $p \neq [] \longrightarrow \text{removeShadowRules1-alternative } p \neq []$
 <proof>

lemma *RS1N-DA*[simp]: *removeShadowRules1-alternative (a@[DenyAll])* = *[DenyAll]*
 <proof>

lemma *WP1n-DA-notinSet*[rule-format]: *wellformed-policy1-strong* $p \longrightarrow$
 $DenyAll \notin set\ (tl\ p)$

<proof>

lemma *mt-sym*: $dom\ a \cap dom\ b = \{\} \implies dom\ b \cap dom\ a = \{\}$

<proof>

lemma *DAnotTL*[rule-format]:

$xs \neq [] \longrightarrow wellformed-policy1\ (xs\ @\ [DenyAll]) \longrightarrow False$

<proof>

lemma *AND-tl*[rule-format]: *allNetsDistinct* (p) $\longrightarrow allNetsDistinct\ (tl\ p)$

<proof>

lemma *distinct-tl*[rule-format]: *distinct* $p \longrightarrow distinct\ (tl\ p)$

<proof>

lemma *SC-tl*[rule-format]: *singleCombinators* (p) $\longrightarrow singleCombinators\ (tl\ p)$

<proof>

lemma *Conc-not-MT*: $p = x\#\!xs \implies p \neq []$

<proof>

lemma *wp1-tl*[rule-format]: $p \neq [] \wedge wellformed-policy1\ p \longrightarrow$

$wellformed-policy1\ (tl\ p)$

<proof>

lemma *wp1-eq*[rule-format]: *wellformed-policy1-strong* $p \implies wellformed-policy1\ p$

<proof>

lemma *wellformed1-alternative-sorted*: *wellformed-policy1-strong* $p \implies$

$wellformed-policy1-strong\ (sort\ p\ l)$

<proof>

lemma *wp1n-RS2*[rule-format]: *wellformed-policy1-strong* $p \longrightarrow$

$wellformed-policy1-strong\ (removeShadowRules2\ p)$

<proof>

lemma *RS2-NMT*[rule-format]: $p \neq [] \longrightarrow removeShadowRules2\ p \neq []$

<proof>

lemma *wp1-alternative-not-mt[simp]: wellformed-policy1-strong $p \implies p \neq []$*
 <proof>

lemma *AIL1[rule-format,simp]: all-in-list $p\ l \longrightarrow$
 all-in-list (removeShadowRules1 p) l*
 <proof>

lemma *wp1ID: wellformed-policy1-strong (insertDeny (removeShadowRules1 p))*
 <proof>

lemma *dRD[simp]: distinct (remdups p)*
 <proof>

lemma *AILrd[rule-format,simp]: all-in-list $p\ l \longrightarrow$ all-in-list (remdups p) l*
 <proof>

lemma *AILiD[rule-format,simp]: all-in-list $p\ l \longrightarrow$ all-in-list (insertDeny p) l*
 <proof>

lemma *SCrd[rule-format,simp]: singleCombinators $p \longrightarrow$ singleCombinators (remdups p)*
 <proof>

lemma *SCRiD[rule-format,simp]: singleCombinators $p \longrightarrow$
 singleCombinators (insertDeny p)*
 <proof>

lemma *WP1rd[rule-format,simp]: wellformed-policy1-strong $p \longrightarrow$
 wellformed-policy1-strong (remdups p)*
 <proof>

lemma *ANDrd[rule-format,simp]: singleCombinators $p \longrightarrow$ allNetsDistinct $p \longrightarrow$
 allNetsDistinct (remdups p)*
 <proof>

lemma *ANDiD[rule-format,simp]: allNetsDistinct $p \longrightarrow$
 allNetsDistinct (insertDeny p)*
 <proof>

lemma *mr-iD[rule-format]: wellformed-policy1-strong $p \longrightarrow$
 matching-rule $x\ p =$ matching-rule x (insertDeny p)*
 <proof>

lemma *WP1iD[rule-format,simp]: wellformed-policy1-strong $p \longrightarrow$*

wellformed-policy1-strong (*insertDeny* *p*)

<proof>

lemma *DAiniD*: *DenyAll* \in *set* (*insertDeny* *p*)

<proof>

lemma *p2lNmt*: *policy2list* *p* \neq []

<proof>

lemma *AIL2*[*rule-format,simp*]: *all-in-list* *p l* \longrightarrow
all-in-list (*removeShadowRules2* *p*) *l*

<proof>

lemma *SCConc*: $\llbracket \text{singleCombinators } x; \text{singleCombinators } y \rrbracket \Longrightarrow$
singleCombinators (*x@y*)

<proof>

lemma *SCp2l*: *singleCombinators* (*policy2list* *p*)

<proof>

lemma *subnetAux*: *Dd* \cap *A* \neq {} \Longrightarrow *A* \subseteq *B* \Longrightarrow *Dd* \cap *B* \neq {}

<proof>

lemma *soadisj*: $\llbracket x \in \text{subnetsOfAdr } a; y \in \text{subnetsOfAdr } a \rrbracket \Longrightarrow \neg \text{netsDistinct } x \ y$

<proof>

lemma *not-member*: $\neg \text{member } a \ (x \oplus y) \Longrightarrow \neg \text{member } a \ x$

<proof>

lemma *soadisj2*: $(\forall \ a \ x \ y. \ x \in \text{subnetsOfAdr } a \wedge y \in \text{subnetsOfAdr } a \longrightarrow$
 $\neg \text{netsDistinct } x \ y)$

<proof>

lemma *ndFalse1*: $\llbracket (\forall \ a \ b \ c \ d. \ (a,b) \in A \wedge (c,d) \in B \longrightarrow \text{netsDistinct } a \ c);$
 $\exists (a, b) \in A. \ a \in \text{subnetsOfAdr } D;$
 $\exists (a, b) \in B. \ a \in \text{subnetsOfAdr } D \rrbracket$
 $\Longrightarrow \text{False}$

<proof>

lemma *ndFalse2*: $\llbracket (\forall \ a \ b \ c \ d. \ (a,b) \in A \wedge (c,d) \in B \longrightarrow \text{netsDistinct } b \ d);$
 $\exists (a, b) \in A. \ b \in \text{subnetsOfAdr } D;$
 $\exists (a, b) \in B. \ b \in \text{subnetsOfAdr } D \rrbracket$
 $\Longrightarrow \text{False}$

<proof>

lemma *tndFalse*: $\llbracket (\forall \ a \ b \ c \ d. \ (a,b) \in A \wedge (c,d) \in B \longrightarrow \text{twoNetsDistinct } a \ b \ c \ d);$

$$\begin{aligned} & \exists (a, b) \in A. a \in \text{subnetsOfAdr } (D::('a::\text{adr})) \wedge b \in \text{subnetsOfAdr } (F::'a); \\ & \exists (a, b) \in B. a \in \text{subnetsOfAdr } D \wedge b \in \text{subnetsOfAdr } F \\ & \implies \text{False} \\ & \langle \text{proof} \rangle
\end{aligned}$$

lemma *sepnMT*[rule-format]: $p \neq [] \longrightarrow (\text{separate } p) \neq []$
 $\langle \text{proof} \rangle$

lemma *sepDA*[rule-format]: $\text{DenyAll} \notin \text{set } p \longrightarrow \text{DenyAll} \notin \text{set } (\text{separate } p)$
 $\langle \text{proof} \rangle$

lemma *setnMT*: $\text{set } a = \text{set } b \implies a \neq [] \implies b \neq []$
 $\langle \text{proof} \rangle$

lemma *sortnMT*: $p \neq [] \implies \text{sort } p \text{ l} \neq []$
 $\langle \text{proof} \rangle$

lemma *idNMT*[rule-format]: $p \neq [] \longrightarrow \text{insertDenies } p \neq []$
 $\langle \text{proof} \rangle$

lemma *OTNoTN*[rule-format]: $\text{OnlyTwoNets } p \longrightarrow x \neq \text{DenyAll} \longrightarrow x \in \text{set } p \longrightarrow$
 $\text{onlyTwoNets } x$
 $\langle \text{proof} \rangle$

lemma *first-isIn*[rule-format]:
 $\neg \text{member DenyAll } x \longrightarrow (\text{first-srcNet } x, \text{first-destNet } x) \in \text{sdnets } x$
 $\langle \text{proof} \rangle$

lemma *sdnets2*: $\llbracket \exists a \ b. \text{sdnets } x = \{(a, b), (b, a)\}; \neg \text{member DenyAll } x \rrbracket \implies$
 $\text{sdnets } x = \{(\text{first-srcNet } x, \text{first-destNet } x),$
 $(\text{first-destNet } x, \text{first-srcNet } x)\}$
 $\langle \text{proof} \rangle$

lemma *alternativelistconc1*[rule-format]: $a \in \text{set } (\text{net-list-aux } [x]) \longrightarrow$
 $a \in \text{set } (\text{net-list-aux } [x, y])$
 $\langle \text{proof} \rangle$

lemma *alternativelistconc2*[rule-format]: $a \in \text{set } (\text{net-list-aux } [x]) \longrightarrow$
 $a \in \text{set } (\text{net-list-aux } [y, x])$
 $\langle \text{proof} \rangle$

lemma *noDA*[rule-format]: $\text{noDenyAll } xs \longrightarrow s \in \text{set } xs \longrightarrow \neg \text{member DenyAll } s$
 $\langle \text{proof} \rangle$

lemma *isInAlternativeList*:
 $(aa \in \text{set } (\text{net-list-aux } [a]) \vee aa \in \text{set } (\text{net-list-aux } p))$
 $\implies aa \in \text{set } (\text{net-list-aux } (a \# p))$
 $\langle \text{proof} \rangle$

lemma *netlistaux*: $x \in \text{set } (\text{net-list-aux } (a \# p)) \implies$
 $x \in \text{set } (\text{net-list-aux } ([a])) \vee x \in \text{set } (\text{net-list-aux } (p))$
 <proof>

lemma *firstInNet*[rule-format]: $\neg \text{member DenyAll } a \longrightarrow$
 $\text{first-destNet } a \in \text{set } (\text{net-list-aux } (a \# p))$
 <proof>

lemma *firstInNeta*[rule-format]: $\neg \text{member DenyAll } a \longrightarrow$
 $\text{first-srcNet } a \in \text{set } (\text{net-list-aux } (a \# p))$
 <proof>

lemma *disjComm*: $\text{disjSD-2 } a \ b \implies \text{disjSD-2 } b \ a$
 <proof>

lemma *disjSD2aux*: $\llbracket \text{disjSD-2 } a \ b; \neg \text{member DenyAll } a; \neg \text{member DenyAll } b \rrbracket \implies$
 $\text{disjSD-2 } (\text{DenyAllFromTo } (\text{first-srcNet } a) (\text{first-destNet } a) \oplus$
 $\text{DenyAllFromTo } (\text{first-destNet } a) (\text{first-srcNet } a) \oplus a) \ b$
 <proof>

lemma *noDA1eq*[rule-format]: $\text{noDenyAll } p \longrightarrow \text{noDenyAll1 } p$
 <proof>

lemma *noDA1C*[rule-format]: $\text{noDenyAll1 } (a \# p) \longrightarrow \text{noDenyAll1 } p$
 <proof>

lemma *disjSD-2IDa*: $\llbracket \text{disjSD-2 } x \ y; \neg \text{member DenyAll } x; \neg \text{member DenyAll } y;$
 $a = (\text{first-srcNet } x); b = (\text{first-destNet } x) \rrbracket \implies$
 $\text{disjSD-2 } ((\text{DenyAllFromTo } a \ b) \oplus (\text{DenyAllFromTo } b \ a) \oplus x) \ y$
 <proof>

lemma *noDAID*[rule-format]: $\text{noDenyAll } p \longrightarrow \text{noDenyAll } (\text{insertDenies } p)$
 <proof>

lemma *isInIDo*[rule-format]: $\text{noDenyAll } p \longrightarrow s \in \text{set } (\text{insertDenies } p) \longrightarrow$
 $(\exists! a. s = (\text{DenyAllFromTo } (\text{first-srcNet } a) (\text{first-destNet } a)) \oplus$
 $(\text{DenyAllFromTo } (\text{first-destNet } a) (\text{first-srcNet } a)) \oplus a \wedge a \in \text{set } p)$
 <proof>

lemma *id-aux1*[rule-format]: $\text{DenyAllFromTo } (\text{first-srcNet } s) (\text{first-destNet } s) \oplus$
 $\text{DenyAllFromTo } (\text{first-destNet } s) (\text{first-srcNet } s) \oplus s \in \text{set } (\text{insertDenies } p)$
 $\longrightarrow s \in \text{set } p$
 <proof>

lemma *id-aux2*:
 $\llbracket \text{noDenyAll } p; (\forall s. s \in \text{set } p \longrightarrow \text{disjSD-2 } a \ s); \neg \text{member DenyAll } a;$
 $((\text{DenyAllFromTo } (\text{first-srcNet } s) (\text{first-destNet } s)) \oplus (\text{DenyAllFromTo } (\text{first-destNet } s) (\text{first-srcNet } s))) \rrbracket$

$(first_destNet\ s)\ (first_srcNet\ s) \oplus s \in set\ (insertDenies\ p) \implies$
 $disjSD-2\ a\ ((DenyAllFromTo\ (first_srcNet\ s)\ (first_destNet\ s)) \oplus$
 $(DenyAllFromTo\ (first_destNet\ s)\ (first_srcNet\ s)) \oplus s)$
 $\langle proof \rangle$

lemma *id-aux4*[rule-format]: $\llbracket noDenyAll\ p; (\forall s. s \in set\ p \longrightarrow$
 $disjSD-2\ a\ s); s \in set\ (insertDenies\ p); \neg member\ DenyAll\ a \rrbracket \implies disjSD-2\ a\ s$
 $\langle proof \rangle$

lemma *sepNetsID*[rule-format]: $noDenyAll1\ p \longrightarrow separated\ p \longrightarrow$
 $separated\ (insertDenies\ p)$
 $\langle proof \rangle$

lemma *aNDDA*[rule-format]: $allNetsDistinct\ p \longrightarrow allNetsDistinct(DenyAll\ \#p)$
 $\langle proof \rangle$

lemma *OTNConc*[rule-format]: $OnlyTwoNets\ (y\ \# z) \longrightarrow OnlyTwoNets\ z$
 $\langle proof \rangle$

lemma *first-bothNetsd*: $\neg member\ DenyAll\ x \implies$
 $first_bothNet\ x = \{first_srcNet\ x, first_destNet\ x\}$
 $\langle proof \rangle$

lemma *bNaux*:
 $\llbracket \neg member\ DenyAll\ x; \neg member\ DenyAll\ y; first_bothNet\ x = first_bothNet\ y \rrbracket$
 $\implies \{first_srcNet\ x, first_destNet\ x\} = \{first_srcNet\ y, first_destNet\ y\}$
 $\langle proof \rangle$

lemma *setPair*: $\{a, b\} = \{a, d\} \implies b = d$
 $\langle proof \rangle$

lemma *setPair1*: $\{a, b\} = \{d, a\} \implies b = d$
 $\langle proof \rangle$

lemma *setPair4*: $\{a, b\} = \{c, d\} \implies a \neq c \implies a = d$
 $\langle proof \rangle$

lemma *otnaux1*: $\{x, y, x, y\} = \{x, y\}$
 $\langle proof \rangle$

lemma *OTNIDaux4*: $\{x, y, x\} = \{y, x\}$
 $\langle proof \rangle$

lemma *setPair5*: $\{a, b\} = \{c, d\} \implies a \neq c \implies a = d$
 $\langle proof \rangle$

lemma *otnaux*:

$\llbracket \text{first-bothNet } x = \text{first-bothNet } y; \neg \text{member DenyAll } x; \neg \text{member DenyAll } y; \\ \text{onlyTwoNets } y; \text{onlyTwoNets } x \rrbracket \implies \\ \text{onlyTwoNets } (x \oplus y)$
 <proof>

lemma OTNSepaux: $\llbracket \text{onlyTwoNets } (a \oplus y) \wedge \text{OnlyTwoNets } z \longrightarrow \\ \text{OnlyTwoNets } (\text{separate } (a \oplus y \# z)); \\ \neg \text{member DenyAll } a; \\ \neg \text{member DenyAll } y; \text{noDenyAll } z; \\ \text{onlyTwoNets } a; \text{OnlyTwoNets } (y \# z); \text{first-bothNet } (a) = \text{first-bothNet } y \rrbracket \\ \implies \text{OnlyTwoNets } (\text{separate } (a \oplus y \# z))$
 <proof>

lemma OTNSEp[rule-format]: $\text{noDenyAll1 } p \longrightarrow \text{OnlyTwoNets } p \longrightarrow \\ \text{OnlyTwoNets } (\text{separate } p)$
 <proof>

lemma nda[rule-format]: $\text{singleCombinators } (a \# p) \longrightarrow \text{noDenyAll } p \longrightarrow \\ \text{noDenyAll1 } (a \# p)$
 <proof>

lemma nDAcharn[rule-format]: $\text{noDenyAll } p = (\forall r \in \text{set } p. \neg \text{member DenyAll } r)$
 <proof>

lemma nDAeqSet: $\text{set } p = \text{set } s \implies \text{noDenyAll } p = \text{noDenyAll } s$
 <proof>

lemma nDASCaux[rule-format]: $\text{DenyAll} \notin \text{set } p \longrightarrow \text{singleCombinators } p \longrightarrow \\ r \in \text{set } p \longrightarrow \neg \text{member DenyAll } r$
 <proof>

lemma nDASC[rule-format]: $\text{wellformed-policy1 } p \longrightarrow \text{singleCombinators } p \longrightarrow \\ \text{noDenyAll1 } p$
 <proof>

lemma noDAAll[rule-format]: $\text{noDenyAll } p = (\neg \text{memberP DenyAll } p)$
 <proof>

lemma memberPsep[symmetric]: $\text{memberP } x \ p = \text{memberP } x \ (\text{separate } p)$
 <proof>

lemma noDAsep[rule-format]: $\text{noDenyAll } p \implies \text{noDenyAll } (\text{separate } p)$
 <proof>

lemma noDA1sep[rule-format]: $\text{noDenyAll1 } p \longrightarrow \text{noDenyAll1 } (\text{separate } p)$
 <proof>

lemma isInAlternativeLista: $(aa \in \text{set } (\text{net-list-aux } [a])) \implies \\ aa \in \text{set } (\text{net-list-aux } (a \# p))$

<proof>

lemma *isInAlternativeListb*: $(aa \in \text{set } (\text{net-list-aux } p)) \implies$
 $aa \in \text{set } (\text{net-list-aux } (a \# p))$

<proof>

lemma *ANDSepaux*: $\text{allNetsDistinct } (x \# y \# z) \implies \text{allNetsDistinct } (x \oplus y \# z)$

<proof>

lemma *netlistalternativeSeparateaux*:

$\text{net-list-aux } [y] @ \text{net-list-aux } z = \text{net-list-aux } (y \# z)$

<proof>

lemma *netlistalternativeSeparate*: $\text{net-list-aux } p = \text{net-list-aux } (\text{separate } p)$

<proof>

lemma *ANDSepaux2*: $\llbracket \text{allNetsDistinct } (x \# y \# z);$
 $\text{allNetsDistinct } (\text{separate } (y \# z)) \rrbracket$
 $\implies \text{allNetsDistinct } (x \# \text{separate } (y \# z))$

<proof>

lemma *ANDSep[rule-format]*: $\text{allNetsDistinct } p \longrightarrow \text{allNetsDistinct } (\text{separate } p)$

<proof>

lemma *wp1-alternativesep[rule-format]*: $\text{wellformed-policy1-strong } p \longrightarrow$
 $\text{wellformed-policy1-strong } (\text{separate } p)$

<proof>

lemma *noDAsort[rule-format]*: $\text{noDenyAll1 } p \longrightarrow \text{noDenyAll1 } (\text{sort } p \ l)$

<proof>

lemma *OTNSC[rule-format]*: $\text{singleCombinators } p \longrightarrow \text{OnlyTwoNets } p$

<proof>

lemma *fMTaux*: $\neg \text{member DenyAll } x \implies \text{first-bothNet } x \neq \{\}$

<proof>

lemma *fl2[rule-format]*: $\text{firstList } (\text{separate } p) = \text{firstList } p$

<proof>

lemma *fl3[rule-format]*: $\text{NetsCollected } p \longrightarrow (\text{first-bothNet } x \neq \text{firstList } p \longrightarrow$
 $(\forall a \in \text{set } p. \text{first-bothNet } x \neq \text{first-bothNet } a) \longrightarrow \text{NetsCollected } (x \# p))$

<proof>

lemma *sortedConc[rule-format]*: $\text{sorted } (a \# p) \ l \longrightarrow \text{sorted } p \ l$

<proof>

lemma *smalleraux2*:

$\{a,b\} \in \text{set } l \implies \{c,d\} \in \text{set } l \implies \{a,b\} \neq \{c,d\} \implies$
 $\text{smaller } (\text{DenyAllFromTo } a \ b) \ (\text{DenyAllFromTo } c \ d) \ l \implies$
 $\neg \text{smaller } (\text{DenyAllFromTo } c \ d) \ (\text{DenyAllFromTo } a \ b) \ l$
 $\langle \text{proof} \rangle$

lemma *smalleraux2a*:

$\{a,b\} \in \text{set } l \implies \{c,d\} \in \text{set } l \implies \{a,b\} \neq \{c,d\} \implies$
 $\text{smaller } (\text{DenyAllFromTo } a \ b) \ (\text{AllowPortFromTo } c \ d \ p) \ l \implies$
 $\neg \text{smaller } (\text{AllowPortFromTo } c \ d \ p) \ (\text{DenyAllFromTo } a \ b) \ l$
 $\langle \text{proof} \rangle$

lemma *smalleraux2b*:

$\{a,b\} \in \text{set } l \implies \{c,d\} \in \text{set } l \implies \{a,b\} \neq \{c,d\} \implies y = \text{DenyAllFromTo } a \ b \implies$
 $\text{smaller } (\text{AllowPortFromTo } c \ d \ p) \ y \ l \implies$
 $\neg \text{smaller } y \ (\text{AllowPortFromTo } c \ d \ p) \ l$
 $\langle \text{proof} \rangle$

lemma *smalleraux2c*:

$\{a,b\} \in \text{set } l \implies \{c,d\} \in \text{set } l \implies \{a,b\} \neq \{c,d\} \implies y = \text{AllowPortFromTo } a \ b \ q \implies$
 $\text{smaller } (\text{AllowPortFromTo } c \ d \ p) \ y \ l \implies \neg \text{smaller } y \ (\text{AllowPortFromTo } c \ d \ p) \ l$
 $\langle \text{proof} \rangle$

lemma *smalleraux3*:

assumes $x \in \text{set } l$
assumes $y \in \text{set } l$
assumes $x \neq y$
assumes $x = \text{bothNet } a$
assumes $y = \text{bothNet } b$
assumes $\text{smaller } a \ b \ l$
assumes $\text{singleCombinators } [a]$
assumes $\text{singleCombinators } [b]$
shows $\neg \text{smaller } b \ a \ l$
 $\langle \text{proof} \rangle$

lemma *smalleraux3a*:

$a \neq \text{DenyAll} \implies b \neq \text{DenyAll} \implies \text{in-list } b \ l \implies \text{in-list } a \ l \implies$
 $\text{bothNet } a \neq \text{bothNet } b \implies \text{smaller } a \ b \ l \implies \text{singleCombinators } [a] \implies$
 $\text{singleCombinators } [b] \implies \neg \text{smaller } b \ a \ l$
 $\langle \text{proof} \rangle$

lemma *posaux[rule-format]*: $\text{position } a \ l < \text{position } b \ l \longrightarrow a \neq b$
 $\langle \text{proof} \rangle$

lemma *posaux6[rule-format]*: $a \in \text{set } l \longrightarrow b \in \text{set } l \longrightarrow a \neq b \longrightarrow$
 $\text{position } a \ l \neq \text{position } b \ l$
 $\langle \text{proof} \rangle$

lemma *notSmallerTransaux*[rule-format]:

$\llbracket x \neq \text{DenyAll}; r \neq \text{DenyAll}; \text{singleCombinators } [x]; \text{singleCombinators } [y];$
 $\text{singleCombinators } [r]; \neg \text{smaller } y \ x \ l; \text{smaller } x \ y \ l; \text{smaller } x \ r \ l;$
 $\text{smaller } y \ r \ l; \text{in-list } x \ l; \text{in-list } y \ l; \text{in-list } r \ l \rrbracket \implies$
 $\neg \text{smaller } r \ x \ l$

<proof>

lemma *notSmallerTrans*[rule-format]:

$x \neq \text{DenyAll} \longrightarrow r \neq \text{DenyAll} \longrightarrow \text{singleCombinators } (x \# y \# z) \longrightarrow$
 $\neg \text{smaller } y \ x \ l \longrightarrow \text{sorted } (x \# y \# z) \ l \longrightarrow r \in \text{set } z \longrightarrow$
 $\text{all-in-list } (x \# y \# z) \ l \longrightarrow \neg \text{smaller } r \ x \ l$

<proof>

lemma *NCSaux1*[rule-format]:

$\text{noDenyAll } p \longrightarrow \{x, y\} \in \text{set } l \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{singleCombinators } p \longrightarrow$
 $\text{sorted } (\text{DenyAllFromTo } x \ y \ \# \ p) \ l \longrightarrow \{x, y\} \neq \text{firstList } p \longrightarrow$
 $\text{DenyAllFromTo } u \ v \in \text{set } p \longrightarrow \{x, y\} \neq \{u, v\}$

<proof>

lemma *posaux3*[rule-format]:

$a \in \text{set } l \longrightarrow b \in \text{set } l \longrightarrow a \neq b \longrightarrow \text{position } a \ l \neq \text{position } b \ l$

<proof>

lemma *posaux4*[rule-format]: $\text{singleCombinators } [a] \longrightarrow a \neq \text{DenyAll} \longrightarrow$

$b \neq \text{DenyAll} \longrightarrow \text{in-list } a \ l \longrightarrow \text{in-list } b \ l \longrightarrow$
 $\text{smaller } a \ b \ l \longrightarrow x = (\text{bothNet } a) \longrightarrow$

$y = (\text{bothNet } b) \longrightarrow \text{position } x \ l \leq \text{position } y \ l$

<proof>

lemma *NCSaux2*[rule-format]:

$\text{noDenyAll } p \longrightarrow \{a, b\} \in \text{set } l \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{singleCombinators } p \longrightarrow$
 $\text{sorted } (\text{DenyAllFromTo } a \ b \ \# \ p) \ l \longrightarrow \{a, b\} \neq \text{firstList } p \longrightarrow$
 $\text{AllowPortFromTo } u \ v \ w \in \text{set } p \longrightarrow \{a, b\} \neq \{u, v\}$

<proof>

lemma *NCSaux3*[rule-format]:

$\text{noDenyAll } p \longrightarrow \{a, b\} \in \text{set } l \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{singleCombinators } p \longrightarrow$
 $\text{sorted } (\text{AllowPortFromTo } a \ b \ w \ \# \ p) \ l \longrightarrow \{a, b\} \neq \text{firstList } p \longrightarrow$
 $\text{DenyAllFromTo } u \ v \in \text{set } p \longrightarrow \{a, b\} \neq \{u, v\}$

<proof>

lemma *NCSaux4*[rule-format]:

$\text{noDenyAll } p \longrightarrow \{a, b\} \in \text{set } l \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{singleCombinators } p \longrightarrow$
 $\text{sorted } (\text{AllowPortFromTo } a \ b \ c \ \# \ p) \ l \longrightarrow \{a, b\} \neq \text{firstList } p \longrightarrow$
 $\text{AllowPortFromTo } u \ v \ w \in \text{set } p \longrightarrow \{a, b\} \neq \{u, v\}$

<proof>

lemma *NetsCollectedSorted*[rule-format]:

noDenyAll1 p \longrightarrow *all-in-list p l* \longrightarrow *singleCombinators p* \longrightarrow *sorted p l* \longrightarrow
NetsCollected p

\langle *proof* \rangle

lemma *NetsCollectedSort*: *distinct p* \implies *noDenyAll1 p* \implies *all-in-list p l* \implies
singleCombinators p \implies *NetsCollected (sort p l)*

\langle *proof* \rangle

lemma *fBNsep*[rule-format]: $(\forall a \in \text{set } z. \{b, c\} \neq \text{first-bothNet } a) \longrightarrow$
 $(\forall a \in \text{set } (\text{separate } z). \{b, c\} \neq \text{first-bothNet } a)$

\langle *proof* \rangle

lemma *fBNsep1*[rule-format]: $(\forall a \in \text{set } z. \text{first-bothNet } x \neq \text{first-bothNet } a) \longrightarrow$
 $(\forall a \in \text{set } (\text{separate } z). \text{first-bothNet } x \neq \text{first-bothNet } a)$

\langle *proof* \rangle

lemma *NetsCollectedSepauxa*:

$\llbracket \{b, c\} \neq \text{firstList } z; \text{noDenyAll1 } z;$
 $(\forall a \in \text{set } z. \{b, c\} \neq \text{first-bothNet } a); \text{NetsCollected } (z);$
 $\text{NetsCollected } (\text{separate } (z)); \{b, c\} \neq \text{firstList } (\text{separate } (z));$
 $a \in \text{set } (\text{separate } (z)) \rrbracket \implies$
 $\{b, c\} \neq \text{first-bothNet } a$

\langle *proof* \rangle

lemma *NetsCollectedSepaux*:

$\llbracket \text{first-bothNet } (x :: ('a, 'b) \text{Combinators}) \neq \text{first-bothNet } y; \neg \text{member } \text{DenyAll } y \wedge$
 $\text{noDenyAll } z;$
 $(\forall a \in \text{set } z. \text{first-bothNet } x \neq \text{first-bothNet } a) \wedge \text{NetsCollected } (y \# z);$
 $\text{NetsCollected } (\text{separate } (y \# z)); \text{first-bothNet } x \neq \text{firstList } (\text{separate } (y \# z));$
 $a \in \text{set } (\text{separate } (y \# z)) \rrbracket \implies$
 $\text{first-bothNet } (x :: ('a, 'b) \text{Combinators}) \neq \text{first-bothNet } (a :: ('a, 'b) \text{Combinators})$

\langle *proof* \rangle

lemma *NetsCollectedSep*[rule-format]: *noDenyAll1 p* \longrightarrow *NetsCollected p* \longrightarrow
NetsCollected (separate p)

\langle *proof* \rangle

lemma *OTNaux*:

$onlyTwoNets\ a \implies \neg member\ DenyAll\ a \implies (x,y) \in sdnets\ a \implies$
 $(x = first-srcNet\ a \wedge y = first-destNet\ a) \vee$
 $(x = first-destNet\ a \wedge y = first-srcNet\ a)$
 <proof>

lemma *sdnets-charn*: $onlyTwoNets\ a \implies \neg member\ DenyAll\ a \implies$
 $sdnets\ a = \{(first-srcNet\ a, first-destNet\ a)\} \vee$
 $sdnets\ a = \{(first-srcNet\ a, first-destNet\ a), (first-destNet\ a, first-srcNet\ a)\}$
 <proof>

lemma *first-bothNet-charn*[*rule-format*]: $\neg member\ DenyAll\ a \longrightarrow$
 $first-bothNet\ a = \{first-srcNet\ a, first-destNet\ a\}$
 <proof>

lemma *sdnets-noteq*:
 $\llbracket onlyTwoNets\ a; onlyTwoNets\ aa; first-bothNet\ a \neq first-bothNet\ aa;$
 $\neg member\ DenyAll\ a; \neg member\ DenyAll\ aa \rrbracket$
 $\implies sdnets\ a \neq sdnets\ aa$
 <proof>

lemma *fbn-noteq*:
 $\llbracket onlyTwoNets\ a; onlyTwoNets\ aa; first-bothNet\ a \neq first-bothNet\ aa;$
 $\neg member\ DenyAll\ a; \neg member\ DenyAll\ aa; allNetsDistinct\ [a, aa] \rrbracket \implies$
 $first-srcNet\ a \neq first-srcNet\ aa \vee first-srcNet\ a \neq first-destNet\ aa \vee$
 $first-destNet\ a \neq first-srcNet\ aa \vee first-destNet\ a \neq first-destNet\ aa$
 <proof>

lemma *NCisSD2aux*:
 $\llbracket onlyTwoNets\ a; onlyTwoNets\ aa; first-bothNet\ a \neq first-bothNet\ aa;$
 $\neg member\ DenyAll\ a; \neg member\ DenyAll\ aa; allNetsDistinct\ [a, aa] \rrbracket \implies$
 $disjSD-2\ a\ aa$
 <proof>

lemma *ANDaux3*[*rule-format*]: $y \in set\ xs \longrightarrow a \in set\ (net-list-aux\ [y]) \longrightarrow$
 $a \in set\ (net-list-aux\ xs)$
 <proof>

lemma *ANDaux2*: $allNetsDistinct\ (x \# xs) \implies y \in set\ xs$
 $\implies allNetsDistinct\ [x,y]$
 <proof>

lemma *NCisSD2*[*rule-format*]:
 $\llbracket \neg member\ DenyAll\ a; OnlyTwoNets\ (a \# p); NetsCollected2\ (a \# p);$
 $NetsCollected\ (a \# p); noDenyAll\ (p); allNetsDistinct\ (a \# p); s \in set\ p \rrbracket \implies$
 $disjSD-2\ a\ s$
 <proof>

lemma *separatedNC*[*rule-format*]:

$OnlyTwoNets\ p \longrightarrow NetsCollected2\ p \longrightarrow NetsCollected\ p \longrightarrow noDenyAll1\ p \longrightarrow$
 $allNetsDistinct\ p \longrightarrow separated\ p$
 $\langle proof \rangle$

lemma $NC2Sep[rule-format]: noDenyAll1\ p \longrightarrow NetsCollected2\ (separate\ p)$
 $\langle proof \rangle$

lemma $separatedSep[rule-format]:$
 $OnlyTwoNets\ p \longrightarrow NetsCollected2\ p \longrightarrow NetsCollected\ p \longrightarrow noDenyAll1\ p \longrightarrow$
 $allNetsDistinct\ p \longrightarrow separated\ (separate\ p)$
 $\langle proof \rangle$

lemma $rADnMT[rule-format]: p \neq [] \longrightarrow removeAllDuplicates\ p \neq []$
 $\langle proof \rangle$

lemma $remDupsNMT[rule-format]: p \neq [] \longrightarrow$
 $remdups\ p \neq []$
 $\langle proof \rangle$

lemma $sets-distinct1: (n::int) \neq m \implies \{(a,b). a = n\} \neq \{(a,b). a = m\}$
 $\langle proof \rangle$

lemma $sets-distinct2: (m::int) \neq n \implies \{(a,b). a = n\} \neq \{(a,b). a = m\}$
 $\langle proof \rangle$

lemma $sets-distinct5: (n::int) < m \implies \{(a,b). a = n\} \neq \{(a,b). a = m\}$
 $\langle proof \rangle$

lemma $sets-distinct6: (m::int) < n \implies \{(a,b). a = n\} \neq \{(a,b). a = m\}$
 $\langle proof \rangle$

end

theory *NormalisationIntegerPortProof*
imports *NormalisationGenericProofs*
begin

Normalisation proofs which are specific to the IntegerPort address representation.

lemma $ConcAssoc: C((A \oplus B) \oplus D) = C(A \oplus (B \oplus D))$
 $\langle proof \rangle$

lemma *aux26[simp]: twoNetsDistinct a b c d \implies
 $\text{dom } (C \text{ (AllowPortFromTo a b p)}) \cap \text{dom } (C \text{ (DenyAllFromTo c d)}) = \{\}$*
 <proof>

lemma *wp2-aux[rule-format]: wellformed-policy2 (xs @ [x]) \longrightarrow
 wellformed-policy2 xs*
 <proof>

lemma *Cdom2: $x \in \text{dom}(C \text{ b}) \implies C \text{ (a } \oplus \text{ b)} \text{ x} = (C \text{ b)} \text{ x}$*
 <proof>

lemma *wp2Conc[rule-format]: wellformed-policy2 (x#xs) \implies wellformed-policy2 xs*
 <proof>

lemma *DAimpliesMR-E[rule-format]: DenyAll \in set p \longrightarrow
 $(\exists \text{ r. applied-rule-rev } C \text{ x p} = \text{Some r})$*
 <proof>

lemma *DAimplieMR[rule-format]: DenyAll \in set p \implies applied-rule-rev C x p \neq None*
 <proof>

lemma *MRList1[rule-format]: $x \in \text{dom } (C \text{ a}) \implies \text{applied-rule-rev } C \text{ x (b@[a])} = \text{Some a}$*
 <proof>

lemma *MRList2: $x \in \text{dom } (C \text{ a}) \implies \text{applied-rule-rev } C \text{ x (c@b@[a])} = \text{Some a}$*
 <proof>

lemma *MRList3: $x \notin \text{dom } (C \text{ xa}) \implies$
 $\text{applied-rule-rev } C \text{ x (a @ b \# xs @ [xa])} = \text{applied-rule-rev } C \text{ x (a @ b \# xs)}$*
 <proof>

lemma *CConcEnd[rule-format]: C a x = Some y \longrightarrow
 $C \text{ (list2FWpolicy (xs @ [a])) x} = \text{Some y}$*
 (is ?P xs)
 <proof>

lemma *CConcStartaux: $\llbracket C \text{ a x} = \text{None} \rrbracket \implies (C \text{ aa ++ C a)} \text{ x} = C \text{ aa x}$*
 <proof>

lemma *CConcStart[rule-format]: xs $\neq [] \longrightarrow C \text{ a x} = \text{None} \longrightarrow$
 $C \text{ (list2FWpolicy (xs @ [a])) x} = C \text{ (list2FWpolicy xs) x}$*
 <proof>

lemma *mrNnt[simp]: applied-rule-rev C x p = Some a $\implies p \neq []$*
 <proof>

lemma *mr-is-C*[*rule-format*]: *applied-rule-rev C x p = Some a* \longrightarrow
 $C (list2FWpolicy (p)) x = C a x$
 <proof>

lemma *CConcStart2*: $\llbracket p \neq []; x \notin dom (C a) \rrbracket \implies$
 $C (list2FWpolicy (p@[a])) x = C (list2FWpolicy p)x$
 <proof>

lemma *CConcEnd1*: $\llbracket q@p \neq []; x \notin dom (C a) \rrbracket \implies$
 $C (list2FWpolicy (q@p@[a])) x = C (list2FWpolicy (q@p))x$
 <proof>

lemma *CConcEnd2*[*rule-format*]: $x \in dom (C a) \longrightarrow$
 $C (list2FWpolicy (xs @ [a])) x = C a x$
 (is ?P xs)
 <proof>

lemma *bar3*: $x \in dom (C (list2FWpolicy (xs @ [xa]))) \implies$
 $x \in dom (C (list2FWpolicy xs)) \vee x \in dom (C xa)$
 <proof>

lemma *CeqEnd*[*rule-format,simp*]: $a \neq [] \longrightarrow x \in dom (C (list2FWpolicy a)) \longrightarrow$
 $C (list2FWpolicy (b@a)) x = (C (list2FWpolicy a)) x$
 <proof>

lemma *CConcStartA*[*rule-format,simp*]: $x \in dom (C a) \longrightarrow$
 $x \in dom (C (list2FWpolicy (a \# b)))$
 (is ?P b)
 <proof>

lemma *domConc*: $\llbracket x \in dom (C (list2FWpolicy b)); b \neq [] \rrbracket \implies$
 $x \in dom (C (list2FWpolicy (a@b)))$
 <proof>

lemma *CeqStart*[*rule-format,simp*]:
 $x \notin dom (C (list2FWpolicy a)) \longrightarrow a \neq [] \longrightarrow b \neq [] \longrightarrow$
 $C (list2FWpolicy (b@a)) x = (C (list2FWpolicy b)) x$
 <proof>

lemma *C-eq-if-mr-eq2*: $\llbracket applied-rule-rev C x a = Some r; applied-rule-rev C x b = Some r; a \neq []; b \neq [] \rrbracket \implies$
 $(C (list2FWpolicy a)) x = (C (list2FWpolicy b)) x$
 <proof>

lemma *nMRtoNone*[*rule-format*]: $p \neq [] \longrightarrow applied-rule-rev C x p = None \longrightarrow$

$C \text{ (list2FWpolicy } p) \ x = \text{None}$
 $\langle \text{proof} \rangle$

lemma *C-eq-if-mr-eq*:
 $\llbracket \text{applied-rule-rev } C \ x \ b = \text{applied-rule-rev } C \ x \ a; a \neq []; b \neq [] \rrbracket \implies$
 $(C \text{ (list2FWpolicy } a)) \ x = (C \text{ (list2FWpolicy } b)) \ x$
 $\langle \text{proof} \rangle$

lemma *notmatching-notdom*: $\text{applied-rule-rev } C \ x \ (p@[a]) \neq \text{Some } a \implies x \notin \text{dom } (C \ a)$
 $\langle \text{proof} \rangle$

lemma *foo3a[rule-format]*: $\text{applied-rule-rev } C \ x \ (a@[b]@c) = \text{Some } b \longrightarrow r \in \text{set } c \longrightarrow$
 $b \notin \text{set } c \longrightarrow x \notin \text{dom } (C \ r)$
 $\langle \text{proof} \rangle$

lemma *foo3D*: $\llbracket \text{wellformed-policy1 } p; p = (\text{DenyAll}\#ps);$
 $\text{applied-rule-rev } C \ x \ p = \text{Some } \text{DenyAll}; r \in \text{set } ps \rrbracket \implies x \notin \text{dom } (C \ r)$
 $\langle \text{proof} \rangle$

lemma *foo4[rule-format]*: $\text{set } p = \text{set } s \wedge (\forall \ r. r \in \text{set } p \longrightarrow x \notin \text{dom } (C \ r)) \longrightarrow$
 $(\forall \ r. r \in \text{set } s \longrightarrow x \notin \text{dom } (C \ r))$
 $\langle \text{proof} \rangle$

lemma *foo5b[rule-format]*: $x \in \text{dom } (C \ b) \longrightarrow (\forall \ r. r \in \text{set } c \longrightarrow x \notin \text{dom } (C \ r)) \longrightarrow$
 $\text{applied-rule-rev } C \ x \ (b\#c) = \text{Some } b$
 $\langle \text{proof} \rangle$

lemma *mr-first*: $\llbracket x \in \text{dom } (C \ b); (\forall \ r. r \in \text{set } c \longrightarrow x \notin \text{dom } (C \ r)); s = b\#c \rrbracket \implies$
 $\text{applied-rule-rev } C \ x \ s = \text{Some } b$
 $\langle \text{proof} \rangle$

lemma *mr-charn[rule-format]*: $a \in \text{set } p \longrightarrow (x \in \text{dom } (C \ a)) \longrightarrow$
 $(\forall \ r. r \in \text{set } p \wedge x \in \text{dom } (C \ r) \longrightarrow r = a) \longrightarrow$
 $\text{applied-rule-rev } C \ x \ p = \text{Some } a$
 $\langle \text{proof} \rangle$

lemma *foo8*: $\llbracket (\forall \ r. r \in \text{set } p \wedge x \in \text{dom } (C \ r) \longrightarrow r = a); \text{set } p = \text{set } s \rrbracket \implies$
 $(\forall \ r. r \in \text{set } s \wedge x \in \text{dom } (C \ r) \longrightarrow r = a)$
 $\langle \text{proof} \rangle$

lemma *mrConcEnd[rule-format]*: $\text{applied-rule-rev } C \ x \ (b \# p) = \text{Some } a \longrightarrow a \neq b \longrightarrow$
 $\text{applied-rule-rev } C \ x \ p = \text{Some } a$
 $\langle \text{proof} \rangle$

lemma *wp3tl[rule-format]*: $\text{wellformed-policy3 } p \longrightarrow \text{wellformed-policy3 } (tl \ p)$
 $\langle \text{proof} \rangle$

lemma *wp3Conc*[rule-format]: *wellformed-policy3* (*a* # *p*) \longrightarrow *wellformed-policy3* *p*
 <proof>

lemma *foo98*[rule-format]: *applied-rule-rev* *C* *x* (*aa* # *p*) = *Some* *a* \longrightarrow *x* \in *dom* (*C* *r*) \longrightarrow
 r \in *set* *p*
 \longrightarrow *a* \in *set* *p*
 <proof>

lemma *mrMTNone*[simp]: *applied-rule-rev* *C* *x* [] = *None*
 <proof>

lemma *DAAux*[simp]: *x* \in *dom* (*C* *DenyAll*)
 <proof>

lemma *mrSet*[rule-format]: *applied-rule-rev* *C* *x* *p* = *Some* *r* \longrightarrow *r* \in *set* *p*
 <proof>

lemma *mr-not-Conc*: *singleCombinators* *p* \implies *applied-rule-rev* *C* *x* *p* \neq *Some* (*a* \oplus *b*)
 <proof>

lemma *foo25*[rule-format]: *wellformed-policy3* (*p*@[*x*]) \longrightarrow *wellformed-policy3* *p*
 <proof>

lemma *mr-in-dom*[rule-format]: *applied-rule-rev* *C* *x* *p* = *Some* *a* \longrightarrow *x* \in *dom* (*C* *a*)
 <proof>

lemma *wp3EndMT*[rule-format]: *wellformed-policy3* (*p*@[*xs*]) \longrightarrow
 $\text{AllowPortFromTo } a \ b \ po \in \text{set } p \longrightarrow$
 $\text{dom } (C \ (\text{AllowPortFromTo } a \ b \ po)) \cap \text{dom } (C \ xs) = \{\}$
 <proof>

lemma *foo29*: $\llbracket \text{dom } (C \ a) \neq \{\}; \text{dom } (C \ a) \cap \text{dom } (C \ b) = \{\} \rrbracket \implies a \neq b$
 <proof>

lemma *foo28*: $\llbracket \text{AllowPortFromTo } a \ b \ po \in \text{set } p;$
 $\text{dom } (C \ (\text{AllowPortFromTo } a \ b \ po)) \neq \{\}; (\text{wellformed-policy3 } (p@[x])) \rrbracket$
 $\implies x \neq \text{AllowPortFromTo } a \ b \ po$
 <proof>

lemma *foo28a*[rule-format]: *x* \in *dom* (*C* *a*) \implies *dom* (*C* *a*) \neq {}
 <proof>

lemma *allow-deny-dom*[simp]: $\text{dom } (C \ (\text{AllowPortFromTo } a \ b \ po)) \subseteq$
 $\text{dom } (C \ (\text{DenyAllFromTo } a \ b))$

<proof>

lemma *DenyAllowDisj*: $\text{dom } (C \text{ (AllowPortFromTo } a \ b \ p)) \neq \{\} \implies$
 $\text{dom } (C \text{ (DenyAllFromTo } a \ b)) \cap \text{dom } (C \text{ (AllowPortFromTo } a \ b \ p)) \neq \{\}$

<proof>

lemma *foo31*: $\llbracket (\forall \ r. \ r \in \text{set } p \wedge x \in \text{dom } (C \ r) \longrightarrow$
 $(r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll});$
 $\text{set } p = \text{set } s \rrbracket \implies$
 $(\forall \ r. \ r \in \text{set } s \wedge x \in \text{dom } (C \ r) \longrightarrow$
 $(r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll}))$

<proof>

lemma *wp1-axa*: *wellformed-policy1-strong* $p \implies (\exists \ r. \text{applied-rule-rev } C \ x \ p = \text{Some } r)$

<proof>

lemma *deny-dom[simp]*: $\text{twoNetsDistinct } a \ b \ c \ d \implies \text{dom } (C \text{ (DenyAllFromTo } a \ b)) \cap$
 $\text{dom } (C \text{ (DenyAllFromTo } c \ d)) = \{\}$

<proof>

lemma *domTrans*: $\llbracket \text{dom } a \subseteq \text{dom } b; \text{dom}(b) \cap \text{dom}(c) = \{\} \rrbracket \implies \text{dom}(a) \cap \text{dom}(c) = \{\}$

<proof>

lemma *DomInterAllowsMT*: $\llbracket \text{twoNetsDistinct } a \ b \ c \ d \rrbracket \implies$
 $\text{dom } (C \text{ (AllowPortFromTo } a \ b \ p)) \cap \text{dom } (C \text{ (AllowPortFromTo } c \ d \ po)) = \{\}$

<proof>

lemma *DomInterAllowsMT-Ports*: $\llbracket p \neq po \rrbracket \implies$
 $\text{dom } (C \text{ (AllowPortFromTo } a \ b \ p)) \cap \text{dom } (C \text{ (AllowPortFromTo } c \ d \ po)) = \{\}$

<proof>

lemma *wellformed-policy3-charn[rule-format]*:
 $\text{singleCombinators } p \longrightarrow \text{distinct } p \longrightarrow \text{allNetsDistinct } p \longrightarrow$
 $\text{wellformed-policy1 } p \longrightarrow \text{wellformed-policy2 } p \longrightarrow \text{wellformed-policy3 } p$

<proof>

lemma *DistinctNetsDenyAllow*:
 $\llbracket \text{DenyAllFromTo } b \ c \in \text{set } p; \text{AllowPortFromTo } a \ d \ po \in \text{set } p; \text{allNetsDistinct } p;$
 $\text{dom } (C \text{ (DenyAllFromTo } b \ c)) \cap \text{dom } (C \text{ (AllowPortFromTo } a \ d \ po)) \neq \{\} \rrbracket$
 $\implies b = a \wedge c = d$

<proof>

lemma *DistinctNetsAllowAllow*:

$\llbracket \text{AllowPortFromTo } b \ c \ po \in \text{set } p; \text{AllowPortFromTo } a \ d \ po \in \text{set } p;$
 $\text{allNetsDistinct } p; \text{dom } (C \ (\text{AllowPortFromTo } b \ c \ po)) \cap$
 $\text{dom } (C \ (\text{AllowPortFromTo } a \ d \ po)) \neq \{\}$ \rrbracket
 $\implies b = a \wedge c = d \wedge po = po$

$\langle \text{proof} \rangle$

lemma *WP2RS2[simp]*:

$\llbracket \text{singleCombinators } p;$
 $\text{distinct } p;$
 $\text{allNetsDistinct } p \rrbracket \implies$
 $\text{wellformed-policy2 } (\text{removeShadowRules2 } p)$

$\langle \text{proof} \rangle$

lemma *AD-aux*: $\llbracket \text{AllowPortFromTo } a \ b \ po \in \text{set } p; \text{DenyAllFromTo } c \ d \in \text{set } p;$

$\text{allNetsDistinct } p; \text{singleCombinators } p;$

$a \neq c \vee b \neq d \rrbracket$

$\implies \text{dom } (C \ (\text{AllowPortFromTo } a \ b \ po)) \cap \text{dom } (C \ (\text{DenyAllFromTo } c \ d)) = \{\}$

$\langle \text{proof} \rangle$

lemma *sorted-WP2[rule-format]*: $\text{sorted } p \ l \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{distinct } p \longrightarrow$

$\text{allNetsDistinct } p \longrightarrow \text{singleCombinators } p \longrightarrow \text{wellformed-policy2 } p$

$\langle \text{proof} \rangle$

lemma *wellformed2-sorted[simp]*: $\llbracket \text{all-in-list } p \ l; \text{distinct } p; \text{allNetsDistinct } p;$

$\text{singleCombinators } p \rrbracket \implies \text{wellformed-policy2 } (\text{sort } p \ l)$

$\langle \text{proof} \rangle$

lemma *wellformed2-sortedQ[simp]*: $\llbracket \text{all-in-list } p \ l; \text{distinct } p; \text{allNetsDistinct } p;$

$\text{singleCombinators } p \rrbracket \implies \text{wellformed-policy2 } (\text{qsort } p \ l)$

$\langle \text{proof} \rangle$

lemma *C-DenyAll[simp]*: $C \ (\text{list2FWpolicy } (xs \ @ \ [\text{DenyAll}])) \ x = \text{Some } (\text{deny } ())$

$\langle \text{proof} \rangle$

lemma *C-eq-RS1n*:

$C(\text{list2FWpolicy } (\text{removeShadowRules1-alternative } p)) = C(\text{list2FWpolicy } p)$
 <proof>

lemma *C-eq-RS1[simp]*: $p \neq [] \implies$
 $C(\text{list2FWpolicy } (\text{removeShadowRules1 } p)) = C(\text{list2FWpolicy } p)$
 <proof>

lemma *EX-MR-aux[rule-format]*: $\text{applied-rule-rev } C \ x \ (\text{DenyAll } \# \ p) \neq \text{Some } \text{DenyAll} \longrightarrow$
 $(\exists y. \text{applied-rule-rev } C \ x \ p = \text{Some } y)$
 <proof>

lemma *EX-MR* : $\llbracket \text{applied-rule-rev } C \ x \ p \neq (\text{Some } \text{DenyAll}); p = \text{DenyAll} \# ps \rrbracket \implies$
 $(\text{applied-rule-rev } C \ x \ p = \text{applied-rule-rev } C \ x \ ps)$
 <proof>

lemma *mr-not-DA*:
 $\llbracket \text{wellformed-policy1-strong } s; \text{applied-rule-rev } C \ x \ p = \text{Some } (\text{DenyAllFromTo } a \ ab);$
 $\text{set } p = \text{set } s \rrbracket \implies \text{applied-rule-rev } C \ x \ s \neq \text{Some } \text{DenyAll}$
 <proof>

lemma *domsMT-notND-DD*:
 $\llbracket \text{dom } (C \ (\text{DenyAllFromTo } a \ b)) \cap \text{dom } (C \ (\text{DenyAllFromTo } c \ d)) \neq \{\} \rrbracket \implies$
 $\neg \text{netsDistinct } a \ c$
 <proof>

lemma *domsMT-notND-DD2*:
 $\llbracket \text{dom } (C \ (\text{DenyAllFromTo } a \ b)) \cap \text{dom } (C \ (\text{DenyAllFromTo } c \ d)) \neq \{\} \rrbracket \implies$
 $\neg \text{netsDistinct } b \ d$
 <proof>

lemma *domsMT-notND-DD3*:
 $\llbracket x \in \text{dom } (C \ (\text{DenyAllFromTo } a \ b)); x \in \text{dom } (C \ (\text{DenyAllFromTo } c \ d)) \rrbracket \implies$
 $\neg \text{netsDistinct } a \ c$
 <proof>

lemma *domsMT-notND-DD4*:
 $\llbracket x \in \text{dom } (C \ (\text{DenyAllFromTo } a \ b)); x \in \text{dom } (C \ (\text{DenyAllFromTo } c \ d)) \rrbracket \implies$
 $\neg \text{netsDistinct } b \ d$
 <proof>

lemma *NetsEq-if-sameP-DD*:
 $\llbracket \text{allNetsDistinct } p; u \in \text{set } p; v \in \text{set } p; u = (\text{DenyAllFromTo } a \ b);$
 $v = (\text{DenyAllFromTo } c \ d); x \in \text{dom } (C \ (u)); x \in \text{dom } (C \ (v)) \rrbracket \implies$
 $a = c \wedge b = d$

<proof>

lemma *rule-charn1*:
 assumes *aND*: *allNetsDistinct* *p*
 and *mr-is-allow*: *applied-rule-rev* *C* *x* *p* = *Some* (*AllowPortFromTo* *a* *b* *po*)
 and *SC*: *singleCombinators* *p*
 and *inp*: *r* ∈ *set* *p*
 and *inDom*: *x* ∈ *dom* (*C* *r*)
 shows (*r* = *AllowPortFromTo* *a* *b* *po* ∨ *r* = *DenyAllFromTo* *a* *b* ∨ *r* = *DenyAll*)
<proof>

lemma *noneMTsubset[rule-format]*: *noneMT* *C* *a* \longrightarrow *set* *b* \subseteq *set* *a* \longrightarrow *noneMT* *C* *b*
<proof>

lemma *nMTSort*: *noneMT* *C* *p* \implies *noneMT* *C* (*sort* *p* *l*)
<proof>

lemma *nMTSortQ*: *noneMT* *C* *p* \implies *noneMT* *C* (*qsort* *p* *l*)
<proof>

lemma *wp3char[rule-format]*: *noneMT* *C* *xs* \wedge *C* (*AllowPortFromTo* *a* *b* *po*) = *empty* \wedge
 wellformed-policy3 (*xs* @ [*DenyAllFromTo* *a* *b*]) \longrightarrow
 AllowPortFromTo *a* *b* *po* \notin *set* *xs*
<proof>

lemma *wp3charn[rule-format]*:
 assumes *domAllow*: *dom* (*C* (*AllowPortFromTo* *a* *b* *po*)) \neq {}
 and *wp3*: *wellformed-policy3* (*xs* @ [*DenyAllFromTo* *a* *b*])
 shows *allowNotInList*: *AllowPortFromTo* *a* *b* *po* \notin *set* *xs*
<proof>

lemma *rule-charn2*:
 assumes *aND*: *allNetsDistinct* *p*
 and *wp1*: *wellformed-policy1* *p*
 and *SC*: *singleCombinators* *p*
 and *wp3*: *wellformed-policy3* *p*
 and *allow-in-list*: *AllowPortFromTo* *c* *d* *po* ∈ *set* *p*
 and *x-in-dom-allow*: *x* ∈ *dom* (*C* (*AllowPortFromTo* *c* *d* *po*))
 shows *applied-rule-rev* *C* *x* *p* = *Some* (*AllowPortFromTo* *c* *d* *po*)

$\langle \text{proof} \rangle$

lemma *rule-charn3*:

$\llbracket \text{wellformed-policy1 } p; \text{allNetsDistinct } p; \text{singleCombinators } p;$
 $\text{wellformed-policy3 } p; \text{applied-rule-rev } C \ x \ p = \text{Some } (\text{DenyAllFromTo } c \ d);$
 $\text{AllowPortFromTo } a \ b \ po \in \text{set } p \rrbracket \implies x \notin \text{dom } (C \ (\text{AllowPortFromTo } a \ b \ po))$

$\langle \text{proof} \rangle$

lemma *rule-charn4*:

assumes *wp1*: *wellformed-policy1* *p*
and *aND*: *allNetsDistinct* *p*
and *SC*: *singleCombinators* *p*
and *wp3*: *wellformed-policy3* *p*
and *DA*: *DenyAll* \notin *set* *p*
and *mr*: *applied-rule-rev* *C* *x* *p* = *Some* (*DenyAllFromTo* *a* *b*)
and *rinp*: *r* \in *set* *p*
and *xindom*: *x* \in *dom* (*C* *r*)
shows *r* = *DenyAllFromTo* *a* *b*
 $\langle \text{proof} \rangle$

lemma *foo31a*: $\llbracket (\forall \ r. \ r \in \text{set } p \wedge x \in \text{dom } (C \ r) \longrightarrow$
 $(r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll}));$
 $\text{set } p = \text{set } s; r \in \text{set } s; x \in \text{dom } (C \ r) \rrbracket \implies$
 $(r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll})$
 $\langle \text{proof} \rangle$

lemma *aux4*[*rule-format*]:

applied-rule-rev *C* *x* (*a*#*p*) = *Some* *a* $\longrightarrow a \notin \text{set } (p) \longrightarrow \text{applied-rule-rev } C \ x \ p = \text{None}$
 $\langle \text{proof} \rangle$

lemma *mrDA-tl*:

assumes *mr-DA*: *applied-rule-rev* *C* *x* *p* = *Some* *DenyAll*
and *wp1n*: *wellformed-policy1-strong* *p*
shows *applied-rule-rev* *C* *x* (*tl* *p*) = *None*
 $\langle \text{proof} \rangle$

lemma *rule-charnDAFT*:

$\llbracket \text{wellformed-policy1-strong } p; \text{allNetsDistinct } p; \text{singleCombinators } p;$
 $\text{wellformed-policy3 } p; \text{applied-rule-rev } C \ x \ p = \text{Some } (\text{DenyAllFromTo } a \ b);$
 $r \in \text{set } (\text{tl } p); x \in \text{dom } (C \ r) \rrbracket \implies r = \text{DenyAllFromTo } a \ b$
 $\langle \text{proof} \rangle$

lemma *mrDenyAll-is-unique*:

$\llbracket \text{wellformed-policy1-strong } p; \text{applied-rule-rev } C \ x \ p = \text{Some } \text{DenyAll};$
 $r \in \text{set } (\text{tl } p) \rrbracket \implies x \notin \text{dom } (C \ r)$
 $\langle \text{proof} \rangle$

theorem *C-eq-Sets-mr*:

assumes *sets-eq*: $\text{set } p = \text{set } s$
and *SC*: *singleCombinators* *p*
and *wp1-p*: *wellformed-policy1-strong* *p*
and *wp1-s*: *wellformed-policy1-strong* *s*
and *wp3-p*: *wellformed-policy3* *p*
and *wp3-s*: *wellformed-policy3* *s*
and *aND*: *allNetsDistinct* *p*

shows *applied-rule-rev* *C* *x* *p* = *applied-rule-rev* *C* *x* *s*

<proof>

lemma *C-eq-Sets*:

$\llbracket \text{singleCombinators } p; \text{wellformed-policy1-strong } p; \text{wellformed-policy1-strong } s;$
 $\text{wellformed-policy3 } p; \text{wellformed-policy3 } s; \text{allNetsDistinct } p; \text{set } p = \text{set } s \rrbracket \implies$
 $C (\text{list2FWpolicy } p) x = C (\text{list2FWpolicy } s) x$
<proof>

lemma *C-eq-sorted*: $\llbracket \text{distinct } p; \text{all-in-list } p l; \text{singleCombinators } p;$

$\text{wellformed-policy1-strong } p; \text{wellformed-policy3 } p; \text{allNetsDistinct } p \rrbracket \implies$
 $C (\text{list2FWpolicy } (\text{sort } p l)) = C (\text{list2FWpolicy } p)$

<proof>

lemma *C-eq-sortedQ*: $\llbracket \text{distinct } p; \text{all-in-list } p l; \text{singleCombinators } p;$

$\text{wellformed-policy1-strong } p; \text{wellformed-policy3 } p; \text{allNetsDistinct } p \rrbracket \implies$
 $C (\text{list2FWpolicy } (\text{qsort } p l)) = C (\text{list2FWpolicy } p)$

<proof>

lemma *C-eq-RS2-mr*: *applied-rule-rev* *C* *x* (*removeShadowRules2* *p*) = *applied-rule-rev* *C* *x* *p*

<proof>

lemma *C-eq-None[rule-format]*: $p \neq [] \implies \text{applied-rule-rev } C x p = \text{None} \longrightarrow$

$C (\text{list2FWpolicy } p) x = \text{None}$

<proof>

lemma *C-eq-None2*:

$\llbracket a \neq []; b \neq []; \text{applied-rule-rev } C x a = \text{None}; \text{applied-rule-rev } C x b = \text{None} \rrbracket \implies$
 $(C (\text{list2FWpolicy } a)) x = (C (\text{list2FWpolicy } b)) x$

<proof>

lemma *C-eq-RS2: wellformed-policy1-strong* $p \implies$
 $C \text{ (list2FWpolicy (removeShadowRules2 } p)) = C \text{ (list2FWpolicy } p)$
 $\langle \text{proof} \rangle$

lemma *noneMTRS2: noneMT* $C \text{ } p \implies \text{noneMT } C \text{ (removeShadowRules2 } p)$
 $\langle \text{proof} \rangle$

lemma *CconcNone: dom (C a) = {} ; p ≠ []* \implies
 $C \text{ (list2FWpolicy (a \# p)) } x = C \text{ (list2FWpolicy } p) \text{ } x$
 $\langle \text{proof} \rangle$

lemma *noneMTrd[rule-format]: noneMT* $C \text{ } p \longrightarrow \text{noneMT } C \text{ (remdups } p)$
 $\langle \text{proof} \rangle$

lemma *DARS3[rule-format]: DenyAll* $\notin \text{set } p \longrightarrow \text{DenyAll} \notin \text{set (removeShadowRules3 } C \text{ } p)$
 $\langle \text{proof} \rangle$

lemma *DAnMT: dom (C DenyAll) ≠ {}*
 $\langle \text{proof} \rangle$

lemma *DAnMT2: C DenyAll ≠ empty*
 $\langle \text{proof} \rangle$

lemma *wp1n-RS3[rule-format,simp]: wellformed-policy1-strong* $p \longrightarrow$
 $\text{wellformed-policy1-strong (removeShadowRules3 } C \text{ } p)$
 $\langle \text{proof} \rangle$

lemma *AILRS3[rule-format,simp]: all-in-list* $p \text{ } l \longrightarrow$
 $\text{all-in-list (removeShadowRules3 } C \text{ } p) \text{ } l$
 $\langle \text{proof} \rangle$

lemma *SCRS3[rule-format,simp]: singleCombinators* $p \longrightarrow$
 $\text{singleCombinators(removeShadowRules3 } C \text{ } p)$
 $\langle \text{proof} \rangle$

lemma *RS3subset: set (removeShadowRules3 C p) ⊆ set p*
 $\langle \text{proof} \rangle$

lemma *ANDRS3[simp]: [singleCombinators p; allNetsDistinct p]* \implies
 $\text{allNetsDistinct (removeShadowRules3 } C \text{ } p)$

$\langle \text{proof} \rangle$

lemma *nlpaux*: $x \notin \text{dom } (C \ b) \implies C \ (a \oplus b) \ x = C \ a \ x$

$\langle \text{proof} \rangle$

lemma *notindom*[*rule-format*]: $a \in \text{set } p \longrightarrow x \notin \text{dom } (C \ (\text{list2FWpolicy } p)) \longrightarrow x \notin \text{dom } (C \ a)$

$\langle \text{proof} \rangle$

lemma *C-eq-rd*[*rule-format*]: $p \neq [] \implies C \ (\text{list2FWpolicy } (\text{remdups } p)) = C \ (\text{list2FWpolicy } p)$

$\langle \text{proof} \rangle$

lemma *nMT-domMT*: $\llbracket \neg \text{notMTpolicy } C \ p; p \neq [] \rrbracket \implies r \notin \text{dom } (C \ (\text{list2FWpolicy } p))$

$\langle \text{proof} \rangle$

lemma *C-eq-RS3-aux*[*rule-format*]: $\text{notMTpolicy } C \ p \implies C \ (\text{list2FWpolicy } p) \ x = C \ (\text{list2FWpolicy } (\text{removeShadowRules3 } C \ p)) \ x$

$\langle \text{proof} \rangle$

lemma *C-eq-id*: $\text{wellformed-policy1-strong } p \implies C \ (\text{list2FWpolicy } (\text{insertDeny } p)) = C \ (\text{list2FWpolicy } p)$

$\langle \text{proof} \rangle$

lemma *C-eq-RS3*: $\text{notMTpolicy } C \ p \implies C \ (\text{list2FWpolicy } (\text{removeShadowRules3 } C \ p)) = C \ (\text{list2FWpolicy } p)$

$\langle \text{proof} \rangle$

lemma *NMPrd*[*rule-format*]: $\text{notMTpolicy } C \ p \longrightarrow \text{notMTpolicy } C \ (\text{remdups } p)$

$\langle \text{proof} \rangle$

lemma *NMPDA*[*rule-format*]: $\text{DenyAll} \in \text{set } p \longrightarrow \text{notMTpolicy } C \ p$

$\langle \text{proof} \rangle$

lemma *NMPiD*[*rule-format*]: $\text{notMTpolicy } C \ (\text{insertDeny } p)$

$\langle \text{proof} \rangle$

lemma *list2FWpolicy2list*[*rule-format*]: $C \ (\text{list2FWpolicy}(\text{policy2list } p)) = (C \ p)$

$\langle \text{proof} \rangle$

lemmas $C\text{-eq-Lemmas} = \text{noneMTRS2 noneMTrd SCp2l}$
 $\text{wp1n-RS2 wp1ID NMPiD wp1alternative-RS1}$
 $\text{p2lNmt list2FWpolicy2list wellformed-policy3-charn waux2 wp1-eq}$

lemmas $C\text{-eq-subst-Lemmas} = C\text{-eq-sorted } C\text{-eq-sortedQ } C\text{-eq-RS2 } C\text{-eq-rd } C\text{-eq-RS3 } C\text{-eq-id}$

lemma $C\text{-eq-All-untilSorted}$:

$\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{ all-in-list } (\text{policy2list } p) \text{ l};$
 $\text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C(\text{list2FWpolicy } (\text{sort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } C$
 $(\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \text{ l})) = C \text{ p}$
 $\langle \text{proof} \rangle$

lemma $C\text{-eq-All-untilSortedQ}$:

$\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{ all-in-list } (\text{policy2list } p) \text{ l};$
 $\text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C(\text{list2FWpolicy } (\text{qsort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } C$
 $(\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \text{ l})) = C \text{ p}$
 $\langle \text{proof} \rangle$

lemma $C\text{-eq-All-untilSorted-withSimps}$:

$\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{ all-in-list } (\text{policy2list } p) \text{ l};$
 $\text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C(\text{list2FWpolicy } (\text{sort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } C (\text{insertDeny}$
 $(\text{removeShadowRules1 } (\text{policy2list } p)))))) \text{ l})) = C \text{ p}$
 $\langle \text{proof} \rangle$

lemma $C\text{-eq-All-untilSorted-withSimpsQ}$:

$\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{ all-in-list } (\text{policy2list } p) \text{ l};$
 $\text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C(\text{list2FWpolicy } (\text{qsort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } C (\text{insertDeny}$
 $(\text{removeShadowRules1 } (\text{policy2list } p)))))) \text{ l})) = C \text{ p}$
 $\langle \text{proof} \rangle$

lemma $\text{InDomConc}[\text{rule-format}]: p \neq [] \longrightarrow x \in \text{dom } (C (\text{list2FWpolicy } (p))) \longrightarrow$
 $x \in \text{dom } (C (\text{list2FWpolicy } (a\#p)))$

$\langle \text{proof} \rangle$

lemma *not-in-member*[rule-format]: $\text{member } a \ b \longrightarrow x \notin \text{dom } (C \ b) \longrightarrow x \notin \text{dom } (C \ a)$
 <proof>

lemma *src-in-sdnets*[rule-format]: $\neg \text{member } \text{DenyAll } x \longrightarrow p \in \text{dom } (C \ x) \longrightarrow$
 $\text{subnetsOfAdr } (\text{src } p) \cap (\text{fst-set } (\text{sdnets } x)) \neq \{\}$
 <proof>

lemma *dest-in-sdnets*[rule-format]: $\neg \text{member } \text{DenyAll } x \longrightarrow p \in \text{dom } (C \ x) \longrightarrow$
 $\text{subnetsOfAdr } (\text{dest } p) \cap (\text{snd-set } (\text{sdnets } x)) \neq \{\}$
 <proof>

lemma *sdnets-in-subnets*[rule-format]: $p \in \text{dom } (C \ x) \longrightarrow \neg \text{member } \text{DenyAll } x \longrightarrow$
 $(\exists (a,b) \in \text{sdnets } x. a \in \text{subnetsOfAdr } (\text{src } p) \wedge b \in \text{subnetsOfAdr } (\text{dest } p))$
 <proof>

lemma *disjSD-no-p-in-both*[rule-format]:
 $\llbracket \text{disjSD-2 } x \ y; \neg \text{member } \text{DenyAll } x; \neg \text{member } \text{DenyAll } y;$
 $p \in \text{dom } (C \ x); p \in \text{dom } (C \ y) \rrbracket \Longrightarrow \text{False}$
 <proof>

lemma *list2FWpolicy-eq*: $zs \neq [] \Longrightarrow$
 $C \ (\text{list2FWpolicy } (x \oplus y \ \# \ z)) \ p = C \ (x \oplus \text{list2FWpolicy } (y \ \# \ z)) \ p$
 <proof>

lemma *dom-sep*[rule-format]: $x \in \text{dom } (C \ (\text{list2FWpolicy } p)) \longrightarrow$
 $x \in \text{dom } (C \ (\text{list2FWpolicy } (\text{separate } p)))$
 <proof>

lemma *domdConcStart*[rule-format]: $x \in \text{dom } (C \ (\text{list2FWpolicy } (a \ \# \ b))) \longrightarrow$
 $x \notin \text{dom } (C \ (\text{list2FWpolicy } b))$
 $\longrightarrow x \in \text{dom } (C \ (a))$
 <proof>

lemma *sep-dom2-aux*: $\llbracket x \in \text{dom } (C \ (\text{list2FWpolicy } (a \oplus y \ \# \ z))) \rrbracket$
 $\Longrightarrow x \in \text{dom } (C \ (a \oplus \text{list2FWpolicy } (y \ \# \ z)))$
 <proof>

lemma *sep-dom2-aux2*:
 $\llbracket (x \in \text{dom } (C \ (\text{list2FWpolicy } (\text{separate } (y \ \# \ z)))) \longrightarrow$
 $x \in \text{dom } (C \ (\text{list2FWpolicy } (y \ \# \ z)));$
 $x \in \text{dom } (C \ (\text{list2FWpolicy } (a \ \# \ \text{separate } (y \ \# \ z)))) \rrbracket$

$\implies x \in \text{dom } (C \text{ (list2FWpolicy } (a \oplus y \# z)))$
 $\langle \text{proof} \rangle$

lemma *sep-dom2*[rule-format]:
 $x \in \text{dom } (C \text{ (list2FWpolicy } (\text{separate } p))) \longrightarrow x \in \text{dom } (C \text{ (list2FWpolicy } (p)))$
 $\langle \text{proof} \rangle$

lemma *sepDom*: $\text{dom } (C \text{ (list2FWpolicy } p)) = \text{dom } (C \text{ (list2FWpolicy } (\text{separate } p)))$
 $\langle \text{proof} \rangle$

lemma *C-eq-s-ext*[rule-format]: $p \neq [] \longrightarrow$
 $C \text{ (list2FWpolicy } (\text{separate } p)) \ a = C \text{ (list2FWpolicy } p) \ a$
 $\langle \text{proof} \rangle$

lemma *C-eq-s*: $p \neq [] \implies C \text{ (list2FWpolicy } (\text{separate } p)) = C \text{ (list2FWpolicy } p)$
 $\langle \text{proof} \rangle$

lemma *sortnMTQ*: $p \neq [] \implies \text{qsort } p \ l \neq []$
 $\langle \text{proof} \rangle$

lemmas *C-eq-Lemmas-sep* =
C-eq-Lemmas sortnMT sortnMTQ RS2-NMT NMPrd NMPRS3

lemma *C-eq-until-separated*:
 $\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{all-in-list } (\text{policy2list } p) \ l; \text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C \text{ (list2FWpolicy } (\text{separate } (\text{sort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } C$
 $(\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \ l))) = C \ p$
 $\langle \text{proof} \rangle$

lemma *C-eq-until-separatedQ*:
 $\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{all-in-list } (\text{policy2list } p) \ l; \text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C \text{ (list2FWpolicy } (\text{separate } (\text{qsort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } C$
 $(\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \ l))) = C \ p$
 $\langle \text{proof} \rangle$

lemma *domID*[rule-format]: $p \neq [] \wedge x \in \text{dom}(C(\text{list2FWpolicy } p)) \longrightarrow$
 $x \in \text{dom } (C(\text{list2FWpolicy}(\text{insertDenies } p)))$
 $\langle \text{proof} \rangle$

lemma *DA-is-deny*:

$x \in \text{dom } (C (\text{DenyAllFromTo } a \ b \oplus \text{DenyAllFromTo } b \ a \oplus \text{DenyAllFromTo } a \ b)) \implies$
 $C (\text{DenyAllFromTo } a \ b \oplus \text{DenyAllFromTo } b \ a \oplus \text{DenyAllFromTo } a \ b) \ x = \text{Some } (\text{deny } ())$
 $\langle \text{proof} \rangle$

lemma *iDdomAux*[rule-format]:
 $p \neq [] \longrightarrow x \notin \text{dom } (C (\text{list2FWpolicy } p)) \longrightarrow$
 $x \in \text{dom } (C (\text{list2FWpolicy } (\text{insertDenies } p))) \longrightarrow$
 $C (\text{list2FWpolicy } (\text{insertDenies } p)) \ x = \text{Some } (\text{deny } ())$
 $\langle \text{proof} \rangle$

lemma *iD-isD*[rule-format]: $p \neq [] \longrightarrow x \notin \text{dom } (C (\text{list2FWpolicy } p))$
 $\longrightarrow C (\text{DenyAll} \oplus \text{list2FWpolicy } (\text{insertDenies } p)) \ x = C \ \text{DenyAll} \ x$
 $\langle \text{proof} \rangle$

lemma *inDomConc*: $\llbracket x \notin \text{dom } (C \ a); x \notin \text{dom } (C (\text{list2FWpolicy } p)) \rrbracket \implies$
 $x \notin \text{dom } (C (\text{list2FWpolicy } (a \# p)))$
 $\langle \text{proof} \rangle$

lemma *domsdisj*[rule-format]: $p \neq [] \longrightarrow (\forall \ x \ s. \ s \in \text{set } p \wedge x \in \text{dom } (C \ A) \longrightarrow$
 $x \notin \text{dom } (C \ s)) \longrightarrow y \in \text{dom } (C \ A) \longrightarrow$
 $y \notin \text{dom } (C (\text{list2FWpolicy } p))$
 $\langle \text{proof} \rangle$

lemma *isSepaux*:
 $\llbracket p \neq []; \text{noDenyAll } (a \# p); \text{separated } (a \ \# \ p);$
 $x \in \text{dom } (C (\text{DenyAllFromTo } (\text{first-srcNet } a) (\text{first-destNet } a) \oplus$
 $\text{DenyAllFromTo } (\text{first-destNet } a) (\text{first-srcNet } a) \oplus a)) \rrbracket \implies$
 $x \notin \text{dom } (C (\text{list2FWpolicy } p))$
 $\langle \text{proof} \rangle$

lemma *noneMTsep*[rule-format]: $\text{noneMT } C \ p \longrightarrow \text{noneMT } C \ (\text{separate } p)$
 $\langle \text{proof} \rangle$

lemma *dom-id*:
 $\llbracket \text{noDenyAll } (a \# p); \text{separated } (a \# p); p \neq []; x \notin \text{dom } (C (\text{list2FWpolicy } p));$
 $x \in \text{dom } (C \ (a)) \rrbracket$
 $\implies x \notin \text{dom } (C (\text{list2FWpolicy } (\text{insertDenies } p)))$
 $\langle \text{proof} \rangle$

lemma *C-eq-iD-aux2*[rule-format]:
 $\text{noDenyAll1 } p \longrightarrow$
 $\text{separated } p \longrightarrow$
 $p \neq [] \longrightarrow$
 $x \in \text{dom } (C (\text{list2FWpolicy } p)) \longrightarrow$

$C(\text{list2FWpolicy } (\text{insertDenies } p)) \ x = C(\text{list2FWpolicy } p) \ x$
 <proof>

lemma *C-eq-iD*: $\llbracket \text{separated } p; \text{noDenyAll1 } p; \text{wellformed-policy1-strong } p \rrbracket \implies$
 $C(\text{list2FWpolicy } (\text{insertDenies } p)) = C(\text{list2FWpolicy } p)$
 <proof>

lemma *noDAsortQ[rule-format]*: $\text{noDenyAll1 } p \longrightarrow \text{noDenyAll1 } (\text{qsort } p \ l)$
 <proof>

lemma *NetsCollectedSortQ*: $\text{distinct } p \implies \text{noDenyAll1 } p \implies \text{all-in-list } p \ l \implies$
 $\text{singleCombinators } p \implies \text{NetsCollected } (\text{qsort } p \ l)$
 <proof>

lemmas *CLemmas* = *nMTSort nMTSortQ noneMTRS2 noneMTrd*
noDAsort noDAsortQ nDASC wp1-eq wp1IID
SCp2l ANDSep wp1n-RS2
OTNSEp OTNSC noDA1sep wp1-alternativesep wellformed-eq
wellformed1-alternative-sorted

lemmas *C-eqLemmas-id* = *CLemmas NC2Sep NetsCollectedSep*
NetsCollectedSort NetsCollectedSortQ separatedNC

lemma *C-eq-Until-InsertDenies*: $\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{all-in-list}$
 $(\text{policy2list } p) \ l; \text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C(\text{list2FWpolicy } ((\text{insertDenies } (\text{separate } (\text{sort } (\text{removeShadowRules2 } (\text{remdups}$
 $(\text{removeShadowRules3 } C(\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \ l)))))) =$
 $C \ p$
 <proof>

lemma *C-eq-Until-InsertDeniesQ*: $\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{all-in-list}$
 $(\text{policy2list } p) \ l; \text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $C(\text{list2FWpolicy } ((\text{insertDenies } (\text{separate } (\text{qsort } (\text{removeShadowRules2 } (\text{remdups}$
 $(\text{removeShadowRules3 } C(\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \ l)))))) =$
 $C \ p$
 <proof>

lemma *C-eq-RD-aux[rule-format]*: $C(p) \ x = C(\text{removeDuplicates } p) \ x$
 <proof>

lemma *C-eq-RAD-aux*[rule-format]:

$p \neq [] \longrightarrow C \text{ (list2FWpolicy } p) \ x = C \text{ (list2FWpolicy (removeAllDuplicates } p)) \ x$
 <proof>

lemma *C-eq-RAD*:

$p \neq [] \implies C \text{ (list2FWpolicy } p) = C \text{ (list2FWpolicy (removeAllDuplicates } p))$
 <proof>

lemma *C-eq-compile*:

$\llbracket \text{DenyAll} \in \text{set (policy2list } p); \text{all-in-list (policy2list } p) \ l; \text{allNetsDistinct (policy2list } p) \rrbracket \implies$
 $C \text{ (list2FWpolicy (removeAllDuplicates (insertDenies (separate (sort$
 $\text{(removeShadowRules2 (remdups (removeShadowRules3 } C \text{ (insertDeny$
 $\text{(removeShadowRules1 (policy2list } p)))))) \ l)))) = C \ p$
 <proof>

lemma *C-eq-compileQ*:

$\llbracket \text{DenyAll} \in \text{set (policy2list } p); \text{all-in-list (policy2list } p) \ l; \text{allNetsDistinct (policy2list } p) \rrbracket \implies$
 $C \text{ (list2FWpolicy (removeAllDuplicates (insertDenies (separate (qsort$
 $\text{(removeShadowRules2 (remdups (removeShadowRules3 } C \text{ (insertDeny$
 $\text{(removeShadowRules1 (policy2list } p)))))) \ l)))) = C \ p$
 <proof>

lemma *C-eq-normalize*:

$\llbracket \text{DenyAll} \in \text{set (policy2list } p); \text{allNetsDistinct (policy2list } p);$
 $\text{all-in-list (policy2list } p) \text{ (Nets-List } p) \rrbracket \implies$
 $C \text{ (list2FWpolicy (normalize } p)) = C \ p$
 <proof>

lemma *C-eq-normalizeQ*:

$\llbracket \text{DenyAll} \in \text{set (policy2list } p); \text{allNetsDistinct (policy2list } p);$
 $\text{all-in-list (policy2list } p) \text{ (Nets-List } p) \rrbracket \implies$
 $C \text{ (list2FWpolicy (normalizeQ } p)) = C \ p$
 <proof>

lemma *domSubset3*: $\text{dom } (C \text{ (DenyAll } \oplus x)) = \text{dom } (C \text{ (DenyAll)})$
 ⟨proof⟩

lemma *domSubset4*: $\text{dom } (C \text{ (DenyAllFromTo } x \ y \oplus \text{DenyAllFromTo } y \ x \oplus \text{AllowPortFromTo } x \ y \text{ dn})) =$
 $\text{dom } (C \text{ (DenyAllFromTo } x \ y \oplus \text{DenyAllFromTo } y \ x))$
 ⟨proof⟩

lemma *domSubset5*:
 $\text{dom } (C \text{ (DenyAllFromTo } x \ y \oplus \text{DenyAllFromTo } y \ x \oplus \text{AllowPortFromTo } y \ x \text{ dn})) =$
 $\text{dom } (C \text{ (DenyAllFromTo } x \ y \oplus \text{DenyAllFromTo } y \ x))$
 ⟨proof⟩

lemma *domSubset1*: $\text{dom } (C \text{ (DenyAllFromTo one two } \oplus \text{DenyAllFromTo two one } \oplus \text{AllowPortFromTo one two dn } \oplus x)) =$
 $\text{dom } (C \text{ (DenyAllFromTo one two } \oplus \text{DenyAllFromTo two one } \oplus x))$
 ⟨proof⟩

lemma *domSubset2*:
 $\text{dom } (C \text{ (DenyAllFromTo one two } \oplus \text{DenyAllFromTo two one } \oplus \text{AllowPortFromTo two one dn } \oplus x)) =$
 $\text{dom } (C \text{ (DenyAllFromTo one two } \oplus \text{DenyAllFromTo two one } \oplus x))$
 ⟨proof⟩

lemma *ConcAssoc2*: $C \text{ (} X \oplus Y \oplus ((A \oplus B) \oplus D)) = C \text{ (} X \oplus Y \oplus A \oplus B \oplus D)$
 ⟨proof⟩

lemma *ConcAssoc3*: $C \text{ (} X \oplus ((Y \oplus A) \oplus D)) = C \text{ (} X \oplus Y \oplus A \oplus D)$
 ⟨proof⟩

lemma *RS3-NMT[rule-format]*: $\text{DenyAll} \in \text{set } p \longrightarrow$
 $\text{removeShadowRules3 } C \ p \neq []$
 ⟨proof⟩

lemma *norm-notMT*: $\text{DenyAll} \in \text{set } (\text{policy2list } p) \implies \text{normalize } p \neq []$
 ⟨proof⟩

lemma *norm-notMTQ*: $\text{DenyAll} \in \text{set } (\text{policy2list } p) \implies \text{normalizeQ } p \neq []$
 ⟨proof⟩

lemma *domDA*: $\text{dom } (C \text{ (DenyAll } \oplus A)) = \text{dom } (C \text{ (DenyAll)})$
 <proof>

lemmas *domain-reasoning* = *domDA ConcAssoc2 domSubset1 domSubset2*
domSubset3 domSubset4 domSubset5 domSubsetDistr1
domSubsetDistr2 domSubsetDistrA domSubsetDistrD coerc-assoc ConcAssoc
ConcAssoc3

The following lemmas help with the normalisation

lemma *list2policyR-Start*[rule-format]: $p \in \text{dom } (C \text{ } a) \longrightarrow$
 $C \text{ (list2policyR } (a \# \text{ list})) \text{ } p = C \text{ } a \text{ } p$
 <proof>

lemma *list2policyR-End*: $p \notin \text{dom } (C \text{ } a) \implies$
 $C \text{ (list2policyR } (a \# \text{ list})) \text{ } p = (C \text{ } a \oplus \text{ list2policy } (\text{map } C \text{ list})) \text{ } p$
 <proof>

lemma *l2polR-eq-el*[rule-format]: $N \neq [] \longrightarrow$
 $C \text{ (list2policyR } N) \text{ } p = \text{ (list2policy } (\text{map } C \text{ } N)) \text{ } p$
 <proof>

lemma *l2polR-eq*: $N \neq [] \implies$
 $C \text{ (list2policyR } N) = \text{ (list2policy } (\text{map } C \text{ } N))$
 <proof>

lemma *list2FWpolicys-eq-el*[rule-format]:
 $\text{Filter} \neq [] \longrightarrow C \text{ (list2policyR } \text{Filter}) \text{ } p = C \text{ (list2FWpolicy } (\text{rev } \text{Filter})) \text{ } p$
 <proof>

lemma *list2FWpolicys-eq*:
 $\text{Filter} \neq [] \implies$
 $C \text{ (list2policyR } \text{Filter}) = C \text{ (list2FWpolicy } (\text{rev } \text{Filter}))$
 <proof>

lemma *list2FWpolicys-eq-sym*:
 $\text{Filter} \neq [] \implies$
 $C \text{ (list2policyR } (\text{rev } \text{Filter})) = C \text{ (list2FWpolicy } \text{Filter})$
 <proof>

lemma *p-eq[rule-format]*: $p \neq [] \longrightarrow$
 $list2policy \ (map \ C \ (rev \ p)) = C \ (list2FWpolicy \ p)$
 $\langle proof \rangle$

lemma *p-eq2[rule-format]*: $normalize \ x \neq [] \longrightarrow$
 $C \ (list2FWpolicy \ (normalize \ x)) = C \ x \longrightarrow$
 $list2policy \ (map \ C \ (rev \ (normalize \ x))) = C \ x$
 $\langle proof \rangle$

lemma *p-eq2Q[rule-format]*: $normalizeQ \ x \neq [] \longrightarrow$
 $C \ (list2FWpolicy \ (normalizeQ \ x)) = C \ x \longrightarrow$
 $list2policy \ (map \ C \ (rev \ (normalizeQ \ x))) = C \ x$
 $\langle proof \rangle$

lemma *list2listNMT[rule-format]*: $x \neq [] \longrightarrow map \ sem \ x \neq []$
 $\langle proof \rangle$

lemma *Norm-Distr2*:
 $r \ o\text{-}f \ ((P \otimes_2 \ (list2policy \ Q)) \ o \ d) =$
 $(list2policy \ ((P \otimes_L \ Q) \ (op \otimes_2 \ r \ d)))$
 $\langle proof \rangle$

lemma *NATDistr*: $\llbracket N \neq []; F = C \ (list2policyR \ N) \rrbracket \Longrightarrow$
 $((\lambda \ (x,y). \ x) \ o\text{-}f \ ((NAT \otimes_2 \ F) \ o \ (\lambda \ x. \ (x,x)))) =$
 $(list2policy \ (\ ((NAT \otimes_L \ (map \ C \ N)) \ (op \otimes_2 \$
 $(\lambda \ (x,y). \ x) \ (\lambda \ x. \ (x,x))))))$
 $\langle proof \rangle$

lemma *C-eq-normalize-manual*:
 $\llbracket DenyAll \in set \ (policy2list \ p);$
 $allNetsDistinct \ (policy2list \ p);$
 $all\text{-}in\text{-}list \ (policy2list \ p) \ ll \rrbracket \Longrightarrow$
 $C \ (list2FWpolicy \ (normalize\text{-}manual\text{-}order \ p \ l)) = C \ p$
 $\langle proof \rangle$

lemma *p-eq2-manualQ[rule-format]*: $normalize\text{-}manual\text{-}orderQ \ x \ l \neq [] \longrightarrow$
 $C \ (list2FWpolicy \ (normalize\text{-}manual\text{-}orderQ \ x \ l)) = C \ x \longrightarrow$

list2policy (map *C* (rev (normalize-manual-orderQ *x l*))) = *C x*
 <proof>

lemma *norm-notMT-manualQ*: *DenyAll* ∈ set (*policy2list p*) ⇒ *normalize-manual-orderQ p*
l ≠ []
 <proof>

lemma *C-eq-normalize-manualQ*:
 [[*DenyAll* ∈ set (*policy2list p*);
allNetsDistinct (*policy2list p*);
all-in-list (*policy2list p*) *l*]] ⇒
C (*list2FWpolicy* (*normalize-manual-orderQ p l*)) = *C p*
 <proof>

lemma *p-eq2-manual[rule-format]*: *normalize-manual-order x l* ≠ [] →
C (*list2FWpolicy* (*normalize-manual-order x l*)) = *C x* →
list2policy (map *C* (rev (*normalize-manual-order x l*))) = *C x*
 <proof>

lemma *norm-notMT-manual*: *DenyAll* ∈ set (*policy2list p*) ⇒ *normalize-manual-order p l* ≠
 []
 <proof>

As an example, how this theorems can be used for a concrete normalisation instantiation.

lemma *normalizeNAT*: [[*DenyAll* ∈ set (*policy2list Filter*);
allNetsDistinct (*policy2list Filter*);
all-in-list (*policy2list Filter*) (*Nets-List Filter*)] ⇒
 ((λ (*x,y*). *x*) o-f (((*NAT* ⊗₂ *C Filter*) o (λ *x*. (*x,x*)))) =
list2policy ((*NAT* ⊗_L (map *C* (rev (*normalize Filter*)))) (op ⊗₂) (λ (*x,y*). *x*) (λ *x*. (*x,x*)))
 <proof>

lemma *domSimpl[simp]*: dom (*C* (*A* ⊕ *DenyAll*)) = dom (*C* (*DenyAll*))
 <proof>

The followin theorems can be applied when prepending the usual normalisation with an additional step and using another semantical interpretation function. This is a general recipe which can be applied whenever one nees to combine several normalisation strategies.

lemma *CRotate-eq-rotateC*: *CRotate p* = *C* (*rotatePolicy p*)

$\langle \text{proof} \rangle$

lemma *DAinRotate*: $\llbracket \text{DenyAll} \in \text{set} (\text{policy2list } p) \rrbracket$
 $\implies \text{DenyAll} \in \text{set} (\text{policy2list} (\text{rotatePolicy } p))$
 $\langle \text{proof} \rangle$

lemma *DAUniv*: $\text{dom} (C\text{Rotate} (P \oplus \text{DenyAll})) = \text{UNIV}$
 $\langle \text{proof} \rangle$

lemma *p-eq2R[rule-format]*: $\text{normalize} (\text{rotatePolicy } x) \neq [] \longrightarrow$
 $C (\text{list2FWpolicy} (\text{normalize} (\text{rotatePolicy } x))) = C\text{Rotate } x \longrightarrow$
 $\text{list2policy} (\text{map } C (\text{rev} (\text{normalize} (\text{rotatePolicy } x)))) = C\text{Rotate } x$
 $\langle \text{proof} \rangle$

lemma *C-eq-normalizeRotate*:
 $\llbracket \text{DenyAll} \in \text{set} (\text{policy2list } p) \rrbracket$;
 $\text{allNetsDistinct} (\text{policy2list} (\text{rotatePolicy } p));$
 $\text{all-in-list} (\text{policy2list} (\text{rotatePolicy } p)) (\text{Nets-List} (\text{rotatePolicy } p)) \rrbracket \implies$
 $C (\text{list2FWpolicy} (\text{removeAllDuplicates} (\text{insertDenies} (\text{separate} (\text{sort}$
 $\quad (\text{removeShadowRules2} (\text{remdups} (\text{removeShadowRules3 } C (\text{insertDeny}$
 $\quad (\text{removeShadowRules1} (\text{policy2list} (\text{rotatePolicy } p)))))))$
 $(\text{Nets-List} (\text{rotatePolicy } p)))))) = C\text{Rotate } p$
 $\langle \text{proof} \rangle$

lemma *C-eq-normalizeRotate2*:
 $\llbracket \text{DenyAll} \in \text{set} (\text{policy2list } p) \rrbracket$;
 $\text{allNetsDistinct} (\text{policy2list} (\text{rotatePolicy } p));$
 $\text{all-in-list} (\text{policy2list} (\text{rotatePolicy } p)) (\text{Nets-List} (\text{rotatePolicy } p)) \rrbracket \implies$
 $C (\text{list2FWpolicy} (\text{normalize} (\text{rotatePolicy } p))) = C\text{Rotate } p$
 $\langle \text{proof} \rangle$

end

theory *NormalisationIPPProofs*
imports *NormalisationIntegerPortProof*
begin

Normalisation proofs which are specific to the IntegerProtocol address representation.

lemma *ConcAssoc*: $Cp((A \oplus B) \oplus D) = Cp(A \oplus (B \oplus D))$
 <proof>

lemma *aux26[simp]*: $twoNetsDistinct\ a\ b\ c\ d \implies$
 $dom\ (Cp\ (AllowPortFromTo\ a\ b\ p)) \cap dom\ (Cp\ (DenyAllFromTo\ c\ d)) = \{\}$
 <proof>

lemma *wp2-aux[rule-format]*: $wellformed-policy2Pr\ (xs\ @\ [x]) \longrightarrow$
 $wellformed-policy2Pr\ xs$
 <proof>

lemma *Cdom2*: $x \in dom(Cp\ b) \implies Cp\ (a \oplus b)\ x = (Cp\ b)\ x$
 <proof>

lemma *wp2Conc[rule-format]*: $wellformed-policy2Pr\ (x\ \#xs) \implies wellformed-policy2Pr\ xs$
 <proof>

lemma *DAimpliesMR-E[rule-format]*: $DenyAll \in set\ p \longrightarrow$
 $(\exists\ r.\ applied-rule-rev\ Cp\ x\ p = Some\ r)$
 <proof>

lemma *DAimplieMR[rule-format]*: $DenyAll \in set\ p \implies applied-rule-rev\ Cp\ x\ p \neq None$
 <proof>

lemma *MRList1[rule-format]*: $x \in dom\ (Cp\ a) \implies applied-rule-rev\ Cp\ x\ (b@[a]) = Some\ a$
 <proof>

lemma *MRList2*: $x \in dom\ (Cp\ a) \implies applied-rule-rev\ Cp\ x\ (c@b@[a]) = Some\ a$
 <proof>

lemma *MRList3*: $x \notin dom\ (Cp\ xa) \implies$
 $applied-rule-rev\ Cp\ x\ (a\ @\ b\ \#xs\ @\ [xa]) = applied-rule-rev\ Cp\ x\ (a\ @\ b\ \#xs)$
 <proof>

lemma *CConcEnd[rule-format]*: $Cp\ a\ x = Some\ y \longrightarrow$
 $Cp\ (list2FWpolicy\ (xs\ @\ [a]))\ x = Some\ y$
 (is ?P xs)
 <proof>

lemma *CConcStartaux*: $\llbracket Cp\ a\ x = None \rrbracket \implies (Cp\ aa\ ++\ Cp\ a)\ x = Cp\ aa\ x$
 <proof>

lemma *CConcStart[rule-format]*: $xs \neq [] \longrightarrow Cp\ a\ x = None \longrightarrow$
 $Cp\ (list2FWpolicy\ (xs\ @\ [a]))\ x = Cp\ (list2FWpolicy\ xs)\ x$
 <proof>

lemma *mrNnt[simp]: applied-rule-rev Cp x p = Some a \implies p \neq []*
<proof>

lemma *mr-is-C[rule-format]: applied-rule-rev Cp x p = Some a \longrightarrow*
Cp (list2FWpolicy (p)) x = Cp a x
<proof>

lemma *CConcStart2: [p \neq []; x \notin dom (Cp a)] \implies*
Cp (list2FWpolicy (p@[a])) x = Cp (list2FWpolicy p)x
<proof>

lemma *CConcEnd1: [q@[p] \neq []; x \notin dom (Cp a)] \implies*
Cp (list2FWpolicy (q@[p@[a]])) x = Cp (list2FWpolicy (q@[p]))x
<proof>

lemma *CConcEnd2[rule-format]: x \in dom (Cp a) \longrightarrow*
Cp (list2FWpolicy (xs @ [a])) x = Cp a x
(is ?P xs)
<proof>

lemma *bar3: x \in dom (Cp (list2FWpolicy (xs @ [xa]))) \implies*
x \in dom (Cp (list2FWpolicy xs)) \vee x \in dom (Cp xa)
<proof>

lemma *CeqEnd[rule-format,simp]: a \neq [] \longrightarrow x \in dom (Cp (list2FWpolicy a)) \longrightarrow*
Cp (list2FWpolicy (b@a)) x = (Cp (list2FWpolicy a)) x
<proof>

lemma *CConcStartA[rule-format,simp]: x \in dom (Cp a) \longrightarrow*
x \in dom (Cp (list2FWpolicy (a # b)))
(is ?P b)
<proof>

lemma *domConc: [x \in dom (Cp (list2FWpolicy b)); b \neq []] \implies*
x \in dom (Cp (list2FWpolicy (a@b)))
<proof>

lemma *CeqStart[rule-format,simp]:*
x \notin dom (Cp (list2FWpolicy a)) \longrightarrow a \neq [] \longrightarrow b \neq [] \longrightarrow
Cp (list2FWpolicy (b@a)) x = (Cp (list2FWpolicy b)) x
<proof>

lemma *C-eq-if-mr-eq2: [applied-rule-rev Cp x a = Some r; applied-rule-rev Cp x b = Some r;*
a \neq []; b \neq []] \implies

$(Cp \text{ (list2FWpolicy } a)) \ x = (Cp \text{ (list2FWpolicy } b)) \ x$
 $\langle \text{proof} \rangle$

lemma *nMRtoNone*[rule-format]: $p \neq [] \longrightarrow \text{applied-rule-rev } Cp \ x \ p = \text{None} \longrightarrow$
 $Cp \text{ (list2FWpolicy } p) \ x = \text{None}$
 $\langle \text{proof} \rangle$

lemma *C-eq-if-mr-eq*:
 $\llbracket \text{applied-rule-rev } Cp \ x \ b = \text{applied-rule-rev } Cp \ x \ a; a \neq []; b \neq [] \rrbracket \implies$
 $(Cp \text{ (list2FWpolicy } a)) \ x = (Cp \text{ (list2FWpolicy } b)) \ x$
 $\langle \text{proof} \rangle$

lemma *notmatching-notdom*: $\text{applied-rule-rev } Cp \ x \ (p@[a]) \neq \text{Some } a \implies x \notin \text{dom } (Cp \ a)$
 $\langle \text{proof} \rangle$

lemma *foo3a*[rule-format]: $\text{applied-rule-rev } Cp \ x \ (a@[b]@c) = \text{Some } b \longrightarrow r \in \text{set } c \longrightarrow$
 $b \notin \text{set } c \longrightarrow x \notin \text{dom } (Cp \ r)$
 $\langle \text{proof} \rangle$

lemma *foo3D*: $\llbracket \text{wellformed-policy1 } p; p = (\text{DenyAll} \# ps);$
 $\text{applied-rule-rev } Cp \ x \ p = \text{Some } \text{DenyAll}; r \in \text{set } ps \rrbracket \implies x \notin \text{dom } (Cp \ r)$
 $\langle \text{proof} \rangle$

lemma *foo4*[rule-format]: $\text{set } p = \text{set } s \wedge (\forall \ r. r \in \text{set } p \longrightarrow x \notin \text{dom } (Cp \ r)) \longrightarrow$
 $(\forall \ r. r \in \text{set } s \longrightarrow x \notin \text{dom } (Cp \ r))$
 $\langle \text{proof} \rangle$

lemma *foo5b*[rule-format]: $x \in \text{dom } (Cp \ b) \longrightarrow (\forall \ r. r \in \text{set } c \longrightarrow x \notin \text{dom } (Cp \ r)) \longrightarrow$
 $\text{applied-rule-rev } Cp \ x \ (b\#c) = \text{Some } b$
 $\langle \text{proof} \rangle$

lemma *mr-first*: $\llbracket x \in \text{dom } (Cp \ b); (\forall \ r. r \in \text{set } c \longrightarrow x \notin \text{dom } (Cp \ r)); s = b\#c \rrbracket \implies$
 $\text{applied-rule-rev } Cp \ x \ s = \text{Some } b$
 $\langle \text{proof} \rangle$

lemma *mr-cha*[rule-format]: $a \in \text{set } p \longrightarrow (x \in \text{dom } (Cp \ a)) \longrightarrow$
 $(\forall \ r. r \in \text{set } p \wedge x \in \text{dom } (Cp \ r) \longrightarrow r = a) \longrightarrow$
 $\text{applied-rule-rev } Cp \ x \ p = \text{Some } a$
 $\langle \text{proof} \rangle$

lemma *foo8*: $\llbracket (\forall \ r. r \in \text{set } p \wedge x \in \text{dom } (Cp \ r) \longrightarrow r = a); \text{set } p = \text{set } s \rrbracket \implies$
 $(\forall \ r. r \in \text{set } s \wedge x \in \text{dom } (Cp \ r) \longrightarrow r = a)$
 $\langle \text{proof} \rangle$

lemma *mrConcEnd*[rule-format]: $\text{applied-rule-rev } Cp \ x \ (b \# p) = \text{Some } a \longrightarrow a \neq b \longrightarrow$
 $\text{applied-rule-rev } Cp \ x \ p = \text{Some } a$
 $\langle \text{proof} \rangle$

lemma *wp3tl[rule-format]: wellformed-policy3Pr p \longrightarrow wellformed-policy3Pr (tl p)*
<proof>

lemma *wp3Conc[rule-format]: wellformed-policy3Pr (a#p) \longrightarrow wellformed-policy3Pr p*
<proof>

lemma *foo98[rule-format]: applied-rule-rev Cp x (aa # p) = Some a \longrightarrow x \in dom (Cp r) \longrightarrow*
r \in set p
 \longrightarrow a \in set p
<proof>

lemma *mrMTNone[simp]: applied-rule-rev Cp x [] = None*
<proof>

lemma *DAAux[simp]: x \in dom (Cp DenyAll)*
<proof>

lemma *mrSet[rule-format]: applied-rule-rev Cp x p = Some r \longrightarrow r \in set p*
<proof>

lemma *mr-not-Conc: singleCombinators p \implies applied-rule-rev Cp x p \neq Some (a \oplus b)*
<proof>

lemma *foo25[rule-format]: wellformed-policy3Pr (p@[x]) \longrightarrow wellformed-policy3Pr p*
<proof>

lemma *mr-in-dom[rule-format]: applied-rule-rev Cp x p = Some a \longrightarrow x \in dom (Cp a)*
<proof>

lemma *wp3EndMT[rule-format]: wellformed-policy3Pr (p@[xs]) \longrightarrow*
AllowPortFromTo a b po \in set p \longrightarrow
dom (Cp (AllowPortFromTo a b po)) \cap dom (Cp xs) = {}
<proof>

lemma *foo29: \llbracket dom (Cp a) \neq {}; dom (Cp a) \cap dom (Cp b) = {} $\rrbracket \implies a \neq b$*
<proof>

lemma *foo28: \llbracket AllowPortFromTo a b po \in set p;*
dom (Cp (AllowPortFromTo a b po)) \neq {}; (wellformed-policy3Pr (p@[x])) \rrbracket
 $\implies x \neq$ AllowPortFromTo a b po
<proof>

lemma *foo28a[rule-format]: x \in dom (Cp a) \implies dom (Cp a) \neq {}*

$\langle \text{proof} \rangle$

lemma *allow-deny-dom[simp]*: $\text{dom } (Cp \text{ (AllowPortFromTo } a \ b \ po)) \subseteq \text{dom } (Cp \text{ (DenyAllFromTo } a \ b))$

$\langle \text{proof} \rangle$

lemma *DenyAllowDisj*: $\text{dom } (Cp \text{ (AllowPortFromTo } a \ b \ p)) \neq \{\} \implies \text{dom } (Cp \text{ (DenyAllFromTo } a \ b)) \cap \text{dom } (Cp \text{ (AllowPortFromTo } a \ b \ p)) \neq \{\}$

$\langle \text{proof} \rangle$

lemma *foo31*: $\llbracket (\forall \ r. \ r \in \text{set } p \wedge x \in \text{dom } (Cp \ r) \longrightarrow (r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll})) ; \text{set } p = \text{set } s \rrbracket \implies (\forall \ r. \ r \in \text{set } s \wedge x \in \text{dom } (Cp \ r) \longrightarrow (r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll}))$

$\langle \text{proof} \rangle$

lemma *wp1-auxa*: *wellformed-policy1-strong* $p \implies (\exists \ r. \text{applied-rule-rev } Cp \ x \ p = \text{Some } r)$

$\langle \text{proof} \rangle$

lemma *deny-dom[simp]*: $\text{twoNetsDistinct } a \ b \ c \ d \implies \text{dom } (Cp \text{ (DenyAllFromTo } a \ b)) \cap \text{dom } (Cp \text{ (DenyAllFromTo } c \ d)) = \{\}$

$\langle \text{proof} \rangle$

lemma *domTrans*: $\llbracket \text{dom } a \subseteq \text{dom } b ; \text{dom}(b) \cap \text{dom}(c) = \{\} \rrbracket \implies \text{dom}(a) \cap \text{dom}(c) = \{\}$

$\langle \text{proof} \rangle$

lemma *DomInterAllowsMT*: $\llbracket \text{twoNetsDistinct } a \ b \ c \ d \rrbracket \implies \text{dom } (Cp \text{ (AllowPortFromTo } a \ b \ p)) \cap \text{dom } (Cp \text{ (AllowPortFromTo } c \ d \ po)) = \{\}$

$\langle \text{proof} \rangle$

lemma *DomInterAllowsMT-Ports*: $\llbracket p \neq po \rrbracket \implies \text{dom } (Cp \text{ (AllowPortFromTo } a \ b \ p)) \cap \text{dom } (Cp \text{ (AllowPortFromTo } c \ d \ po)) = \{\}$

$\langle \text{proof} \rangle$

lemma *wellformed-policy3-charn[rule-format]*:
 $\text{singleCombinators } p \longrightarrow \text{distinct } p \longrightarrow \text{allNetsDistinct } p \longrightarrow \text{wellformed-policy1 } p \longrightarrow \text{wellformed-policy2Pr } p \longrightarrow \text{wellformed-policy3Pr } p$

$\langle \text{proof} \rangle$

lemma *DistinctNetsDenyAllow*:

$\llbracket \text{DenyAllFromTo } b \ c \in \text{set } p; \text{AllowPortFromTo } a \ d \ po \in \text{set } p; \text{allNetsDistinct } p; \\ \text{dom } (Cp \ (\text{DenyAllFromTo } b \ c)) \cap \text{dom } (Cp \ (\text{AllowPortFromTo } a \ d \ po)) \neq \{\}\rrbracket \\ \implies b = a \wedge c = d$
 <proof>

lemma *DistinctNetsAllowAllow*:

$\llbracket \text{AllowPortFromTo } b \ c \ po \in \text{set } p; \text{AllowPortFromTo } a \ d \ po \in \text{set } p; \\ \text{allNetsDistinct } p; \text{dom } (Cp \ (\text{AllowPortFromTo } b \ c \ po)) \cap \\ \text{dom } (Cp \ (\text{AllowPortFromTo } a \ d \ po)) \neq \{\}\rrbracket \\ \implies b = a \wedge c = d \wedge po = po$
 <proof>

lemma *WP2RS2[simp]*:

$\llbracket \text{singleCombinators } p; \\ \text{distinct } p; \\ \text{allNetsDistinct } p \rrbracket \implies \\ \text{wellformed-policy2Pr } (\text{removeShadowRules2 } p)$
 <proof>

lemma *AD-aux*: $\llbracket \text{AllowPortFromTo } a \ b \ po \in \text{set } p; \text{DenyAllFromTo } c \ d \in \text{set } p; \\ \text{allNetsDistinct } p; \text{singleCombinators } p; \\ a \neq c \vee b \neq d \rrbracket \\ \implies \text{dom } (Cp \ (\text{AllowPortFromTo } a \ b \ po)) \cap \text{dom } (Cp \ (\text{DenyAllFromTo } c \ d)) = \{\}$
 <proof>

lemma *sorted-WP2[rule-format]*: $\text{sorted } p \ l \longrightarrow \text{all-in-list } p \ l \longrightarrow \text{distinct } p \longrightarrow \\ \text{allNetsDistinct } p \longrightarrow \text{singleCombinators } p \longrightarrow \text{wellformed-policy2Pr } p$
 <proof>

lemma *wellformed2-sorted[simp]*: $\llbracket \text{all-in-list } p \ l; \text{distinct } p; \text{allNetsDistinct } p; \\ \text{singleCombinators } p \rrbracket \implies \text{wellformed-policy2Pr } (\text{sort } p \ l)$
 <proof>

lemma *wellformed2-sortedQ[simp]*: $\llbracket \text{all-in-list } p \ l; \text{distinct } p; \text{allNetsDistinct } p; \\ \text{singleCombinators } p \rrbracket \implies \text{wellformed-policy2Pr } (\text{qsort } p \ l)$
 <proof>

lemma *C-DenyAll[simp]*: $Cp \ (\text{list2FWpolicy } (xs \ @ \ [\text{DenyAll}])) \ x = \text{Some } (\text{deny } ())$

$\langle proof \rangle$

lemma *C-eq-RS1n*:

$$Cp(list2FWpolicy (removeShadowRules1-alternative p)) = Cp(list2FWpolicy p)$$

$\langle proof \rangle$

lemma *C-eq-RS1[simp]*: $p \neq [] \implies$

$$Cp(list2FWpolicy (removeShadowRules1 p)) = Cp(list2FWpolicy p)$$

$\langle proof \rangle$

lemma *EX-MR-aux[rule-format]*: $applied-rule-rev Cp x (DenyAll \# p) \neq Some DenyAll \longrightarrow$
 $(\exists y. applied-rule-rev Cp x p = Some y)$

$\langle proof \rangle$

lemma *EX-MR* : $\llbracket applied-rule-rev Cp x p \neq (Some DenyAll); p = DenyAll \# ps \rrbracket \implies$
 $(applied-rule-rev Cp x p = applied-rule-rev Cp x ps)$

$\langle proof \rangle$

lemma *mr-not-DA*:

$$\llbracket wellformed-policy1-strong s; applied-rule-rev Cp x p = Some (DenyAllFromTo a ab);$$
$$set p = set s \rrbracket \implies applied-rule-rev Cp x s \neq Some DenyAll$$

$\langle proof \rangle$

lemma *domsMT-notND-DD*:

$$\llbracket dom (Cp (DenyAllFromTo a b)) \cap dom (Cp (DenyAllFromTo c d)) \neq \{\} \rrbracket \implies$$
$$\neg netsDistinct a c$$

$\langle proof \rangle$

lemma *domsMT-notND-DD2*:

$$\llbracket dom (Cp (DenyAllFromTo a b)) \cap dom (Cp (DenyAllFromTo c d)) \neq \{\} \rrbracket \implies$$
$$\neg netsDistinct b d$$

$\langle proof \rangle$

lemma *domsMT-notND-DD3*:

$$\llbracket x \in dom (Cp (DenyAllFromTo a b)); x \in dom (Cp (DenyAllFromTo c d)) \rrbracket \implies$$
$$\neg netsDistinct a c$$

$\langle proof \rangle$

lemma *domsMT-notND-DD4*:

$$\llbracket x \in dom (Cp (DenyAllFromTo a b)); x \in dom (Cp (DenyAllFromTo c d)) \rrbracket \implies$$
$$\neg netsDistinct b d$$

$\langle proof \rangle$

lemma *NetsEq-if-sameP-DD*:

$\llbracket \text{allNetsDistinct } p; u \in \text{set } p; v \in \text{set } p; u = (\text{DenyAllFromTo } a \ b);$
 $v = (\text{DenyAllFromTo } c \ d); x \in \text{dom } (\text{Cp } (u)); x \in \text{dom } (\text{Cp } (v)) \rrbracket \implies$
 $a = c \wedge b = d$

$\langle \text{proof} \rangle$

lemma *rule-charn1*:

assumes *aND*: *allNetsDistinct* *p*

and *mr-is-allow*: *applied-rule-rev* *Cp x p = Some (AllowPortFromTo a b po)*

and *SC*: *singleCombinators* *p*

and *inp*: *r* \in *set p*

and *inDom*: *x* \in *dom (Cp r)*

shows $(r = \text{AllowPortFromTo } a \ b \ po \vee r = \text{DenyAllFromTo } a \ b \vee r = \text{DenyAll})$

$\langle \text{proof} \rangle$

lemma *noneMTsubset[rule-format]*: *noneMT Cp a \longrightarrow set b \subseteq set a \longrightarrow noneMT Cp b*

$\langle \text{proof} \rangle$

lemma *nMTSort*: *noneMT Cp p \implies noneMT Cp (sort p l)*

$\langle \text{proof} \rangle$

lemma *nMTSortQ*: *noneMT Cp p \implies noneMT Cp (qsort p l)*

$\langle \text{proof} \rangle$

lemma *wp3char[rule-format]*: *noneMT Cp xs \wedge Cp (AllowPortFromTo a b po) = empty \wedge*
wellformed-policy3Pr (xs @ [DenyAllFromTo a b]) \longrightarrow
AllowPortFromTo a b po \notin set xs

$\langle \text{proof} \rangle$

lemma *wp3charn[rule-format]*:

assumes *domAllow*: *dom (Cp (AllowPortFromTo a b po)) \neq {}*

and *wp3*: *wellformed-policy3Pr (xs @ [DenyAllFromTo a b])*

shows *allowNotInList*: *AllowPortFromTo a b po \notin set xs*

$\langle \text{proof} \rangle$

lemma *rule-charn2*:

assumes *aND*: *allNetsDistinct* *p*

and *wp1*: *wellformed-policy1 p*

and *SC*: *singleCombinators* *p*

and *wp3*: *wellformed-policy3Pr p*
and *allow-in-list*: *AllowPortFromTo c d po ∈ set p*
and *x-in-dom-allow*: *x ∈ dom (Cp (AllowPortFromTo c d po))*
shows *applied-rule-rev Cp x p = Some (AllowPortFromTo c d po)*
 ⟨*proof*⟩

lemma *rule-charn3*:
 [[*wellformed-policy1 p*; *allNetsDistinct p*; *singleCombinators p*;
wellformed-policy3Pr p; *applied-rule-rev Cp x p = Some (DenyAllFromTo c d)*;
AllowPortFromTo a b po ∈ set p]] $\implies x \notin \text{dom } (Cp (AllowPortFromTo a b po))$
 ⟨*proof*⟩

lemma *rule-charn4*:
assumes *wp1*: *wellformed-policy1 p*
and *aND*: *allNetsDistinct p*
and *SC*: *singleCombinators p*
and *wp3*: *wellformed-policy3Pr p*
and *DA*: *DenyAll ∉ set p*
and *mr*: *applied-rule-rev Cp x p = Some (DenyAllFromTo a b)*
and *rinp*: *r ∈ set p*
and *xindom*: *x ∈ dom (Cp r)*
shows *r = DenyAllFromTo a b*
 ⟨*proof*⟩

lemma *foo31a*: [[$(\forall r. r \in \text{set } p \wedge x \in \text{dom } (Cp r) \longrightarrow$
 $(r = \text{AllowPortFromTo } a b po \vee r = \text{DenyAllFromTo } a b \vee r = \text{DenyAll}))$;
 $\text{set } p = \text{set } s$; $r \in \text{set } s$; $x \in \text{dom } (Cp r)$]] \implies
 $(r = \text{AllowPortFromTo } a b po \vee r = \text{DenyAllFromTo } a b \vee r = \text{DenyAll})$
 ⟨*proof*⟩

lemma *aux4*[*rule-format*]:
applied-rule-rev Cp x (a#p) = Some a $\longrightarrow a \notin \text{set } (p) \longrightarrow \text{applied-rule-rev Cp x p} = \text{None}$
 ⟨*proof*⟩

lemma *mrDA-tl*:
assumes *mr-DA*: *applied-rule-rev Cp x p = Some DenyAll*
and *wp1n*: *wellformed-policy1-strong p*
shows *applied-rule-rev Cp x (tl p) = None*
 ⟨*proof*⟩

lemma *rule-charnDAFT*:
 [[*wellformed-policy1-strong p*; *allNetsDistinct p*; *singleCombinators p*;
wellformed-policy3Pr p; *applied-rule-rev Cp x p = Some (DenyAllFromTo a b)*;
 $r \in \text{set } (tl p)$; $x \in \text{dom } (Cp r)$]] $\implies r = \text{DenyAllFromTo } a b$
 ⟨*proof*⟩

lemma *mrDenyAll-is-unique*:

[[*wellformed-policy1-strong* p ; *applied-rule-rev* $Cp\ x\ p = \text{Some } \text{DenyAll}$;
 $r \in \text{set } (tl\ p)$]] $\implies x \notin \text{dom } (Cp\ r)$
 <proof>

theorem *C-eq-Sets-mr*:

assumes *sets-eq*: $\text{set } p = \text{set } s$
and *SC*: *singleCombinators* p
and *wp1-p*: *wellformed-policy1-strong* p
and *wp1-s*: *wellformed-policy1-strong* s
and *wp3-p*: *wellformed-policy3Pr* p
and *wp3-s*: *wellformed-policy3Pr* s
and *aND*: *allNetsDistinct* p

shows *applied-rule-rev* $Cp\ x\ p = \text{applied-rule-rev } Cp\ x\ s$
 <proof>

lemma *C-eq-Sets*:

[[*singleCombinators* p ; *wellformed-policy1-strong* p ; *wellformed-policy1-strong* s ;
wellformed-policy3Pr p ; *wellformed-policy3Pr* s ; *allNetsDistinct* p ; $\text{set } p = \text{set } s$]] \implies
 $Cp\ (\text{list2FWpolicy } p)\ x = Cp\ (\text{list2FWpolicy } s)\ x$
 <proof>

lemma *C-eq-sorted*: [[*distinct* p ; *all-in-list* $p\ l$; *singleCombinators* p ;

wellformed-policy1-strong p ; *wellformed-policy3Pr* p ; *allNetsDistinct* p]] \implies

$Cp\ (\text{list2FWpolicy } (\text{sort } p\ l)) = Cp\ (\text{list2FWpolicy } p)$

<proof>

lemma *C-eq-sortedQ*: [[*distinct* p ; *all-in-list* $p\ l$; *singleCombinators* p ;

wellformed-policy1-strong p ; *wellformed-policy3Pr* p ; *allNetsDistinct* p]] \implies

$Cp\ (\text{list2FWpolicy } (qsort\ p\ l)) = Cp\ (\text{list2FWpolicy } p)$

<proof>

lemma *C-eq-RS2-mr*: *applied-rule-rev* $Cp\ x\ (\text{removeShadowRules2 } p) = \text{applied-rule-rev } Cp\ x\ p$

<proof>

lemma *C-eq-None[rule-format]*: $p \neq [] \implies \text{applied-rule-rev } Cp\ x\ p = \text{None} \implies$

$Cp\ (\text{list2FWpolicy } p)\ x = \text{None}$

<proof>

lemma *C-eq-None2*:

$\llbracket a \neq []; b \neq []; \text{applied-rule-rev } Cp \ x \ a = \text{None}; \text{applied-rule-rev } Cp \ x \ b = \text{None} \rrbracket \implies$
 $(Cp \ (\text{list2FWpolicy } a)) \ x = (Cp \ (\text{list2FWpolicy } b)) \ x$
 $\langle \text{proof} \rangle$

lemma *C-eq-RS2*: *wellformed-policy1-strong* $p \implies$

$Cp \ (\text{list2FWpolicy } (\text{removeShadowRules2 } p)) = Cp \ (\text{list2FWpolicy } p)$
 $\langle \text{proof} \rangle$

lemma *noneMTRS2*: *noneMT* $Cp \ p \implies \text{noneMT } Cp \ (\text{removeShadowRules2 } p)$

$\langle \text{proof} \rangle$

lemma *CconcNone*: $\llbracket \text{dom } (Cp \ a) = \{\}; p \neq [] \rrbracket \implies$

$Cp \ (\text{list2FWpolicy } (a \ \# \ p)) \ x = Cp \ (\text{list2FWpolicy } p) \ x$
 $\langle \text{proof} \rangle$

lemma *noneMTrd[rule-format]*: *noneMT* $Cp \ p \longrightarrow \text{noneMT } Cp \ (\text{remdups } p)$

$\langle \text{proof} \rangle$

lemma *DARS3[rule-format]*: *DenyAll* $\notin \text{set } p \longrightarrow \text{DenyAll} \notin \text{set } (\text{removeShadowRules3 } Cp \ p)$

$\langle \text{proof} \rangle$

lemma *DAnMT*: *dom* $(Cp \ \text{DenyAll}) \neq \{\}$

$\langle \text{proof} \rangle$

lemma *DAnMT2*: $Cp \ \text{DenyAll} \neq \text{empty}$

$\langle \text{proof} \rangle$

lemma *wp1n-RS3[rule-format,simp]*: *wellformed-policy1-strong* $p \longrightarrow$

wellformed-policy1-strong $(\text{removeShadowRules3 } Cp \ p)$
 $\langle \text{proof} \rangle$

lemma *AILRS3[rule-format,simp]*: *all-in-list* $p \ l \longrightarrow$

all-in-list $(\text{removeShadowRules3 } Cp \ p) \ l$
 $\langle \text{proof} \rangle$

lemma *SCRS3[rule-format,simp]*: *singleCombinators* $p \longrightarrow$

singleCombinators $(\text{removeShadowRules3 } Cp \ p)$
 $\langle \text{proof} \rangle$

lemma *RS3subset*: *set* $(\text{removeShadowRules3 } Cp \ p) \subseteq \text{set } p$

$\langle \text{proof} \rangle$

lemma *ANDRS3[simp]*: $\llbracket \text{singleCombinators } p; \text{allNetsDistinct } p \rrbracket \implies$
 $\text{allNetsDistinct } (\text{removeShadowRules3 } Cp \ p)$
 $\langle \text{proof} \rangle$

lemma *nlpaux*: $x \notin \text{dom } (Cp \ b) \implies Cp \ (a \oplus b) \ x = Cp \ a \ x$
 $\langle \text{proof} \rangle$

lemma *notindom[rule-format]*: $a \in \text{set } p \longrightarrow x \notin \text{dom } (Cp \ (\text{list2FWpolicy } p)) \longrightarrow$
 $x \notin \text{dom } (Cp \ a)$
 $\langle \text{proof} \rangle$

lemma *C-eq-rd[rule-format]*: $p \neq [] \implies$
 $Cp \ (\text{list2FWpolicy } (\text{remdups } p)) = Cp \ (\text{list2FWpolicy } p)$
 $\langle \text{proof} \rangle$

lemma *nMT-domMT*: $\llbracket \neg \text{notMTpolicy } Cp \ p; p \neq [] \rrbracket \implies r \notin \text{dom } (Cp \ (\text{list2FWpolicy } p))$
 $\langle \text{proof} \rangle$

lemma *C-eq-RS3-aux[rule-format]*: $\text{notMTpolicy } Cp \ p \implies$
 $Cp \ (\text{list2FWpolicy } p) \ x = Cp \ (\text{list2FWpolicy } (\text{removeShadowRules3 } Cp \ p)) \ x$
 $\langle \text{proof} \rangle$

lemma *C-eq-id: wellformed-policy1-strong* $p \implies$
 $Cp(\text{list2FWpolicy } (\text{insertDeny } p)) = Cp \ (\text{list2FWpolicy } p)$
 $\langle \text{proof} \rangle$

lemma *C-eq-RS3*: $\text{notMTpolicy } Cp \ p \implies$
 $Cp(\text{list2FWpolicy } (\text{removeShadowRules3 } Cp \ p)) = Cp \ (\text{list2FWpolicy } p)$
 $\langle \text{proof} \rangle$

lemma *NMPrd[rule-format]*: $\text{notMTpolicy } Cp \ p \longrightarrow \text{notMTpolicy } Cp \ (\text{remdups } p)$
 $\langle \text{proof} \rangle$

lemma *NMPDA[rule-format]*: $\text{DenyAll} \in \text{set } p \longrightarrow \text{notMTpolicy } Cp \ p$
 $\langle \text{proof} \rangle$

lemma *NMPiD[rule-format]*: $\text{notMTpolicy } Cp \ (\text{insertDeny } p)$
 $\langle \text{proof} \rangle$

lemma *list2FWpolicy2list*[rule-format]: $Cp\ (list2FWpolicy(policy2list\ p)) = (Cp\ p)$
 <proof>

lemmas *C-eq-Lemmas* = noneMTRS2 noneMTrd SCp2l
 wp1n-RS2 wp1ID NMPiD wp1alternative-RS1
 p2lNmt list2FWpolicy2list wellformed-policy3-chaen waux2 wp1-eq

lemmas *C-eq-subst-Lemmas* = C-eq-sorted C-eq-sortedQ C-eq-RS2 C-eq-rd C-eq-RS3 C-eq-id

lemma *C-eq-All-untilSorted*:

$\llbracket DenyAll \in set\ (policy2list\ p); all-in-list\ (policy2list\ p)\ l;$
 $allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp(list2FWpolicy\ (sort\ (removeShadowRules2\ (remdups\ (removeShadowRules3\ Cp$
 $(insertDeny\ (removeShadowRules1\ (policy2list\ p))))))\ l)) = Cp\ p$
 <proof>

lemma *C-eq-All-untilSortedQ*:

$\llbracket DenyAll \in set\ (policy2list\ p); all-in-list\ (policy2list\ p)\ l;$
 $allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp(list2FWpolicy\ (qsort\ (removeShadowRules2\ (remdups\ (removeShadowRules3\ Cp$
 $(insertDeny\ (removeShadowRules1\ (policy2list\ p))))))\ l)) = Cp\ p$
 <proof>

lemma *C-eq-All-untilSorted-withSimps*:

$\llbracket DenyAll \in set\ (policy2list\ p); all-in-list\ (policy2list\ p)\ l;$
 $allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp(list2FWpolicy\ (sort\ (removeShadowRules2\ (remdups\ (removeShadowRules3\ Cp\ (insertDeny$
 $(removeShadowRules1\ (policy2list\ p))))))\ l)) = Cp\ p$
 <proof>

lemma *C-eq-All-untilSorted-withSimpsQ*:

$\llbracket DenyAll \in set\ (policy2list\ p); all-in-list\ (policy2list\ p)\ l;$
 $allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp(list2FWpolicy\ (qsort\ (removeShadowRules2\ (remdups\ (removeShadowRules3\ Cp\ (insertDeny$
 $(removeShadowRules1\ (policy2list\ p))))))\ l)) = Cp\ p$
 <proof>

lemma *InDomConc*[rule-format]: $p \neq [] \longrightarrow x \in \text{dom } (Cp \text{ (list2FWpolicy } (p))) \longrightarrow$
 $x \in \text{dom } (Cp \text{ (list2FWpolicy } (a\#p)))$
 ⟨proof⟩

lemma *not-in-member*[rule-format]: $\text{member } a \ b \longrightarrow x \notin \text{dom } (Cp \ b) \longrightarrow x \notin \text{dom } (Cp \ a)$
 ⟨proof⟩

lemma *src-in-sdnets*[rule-format]: $\neg \text{member } \text{DenyAll } x \longrightarrow p \in \text{dom } (Cp \ x) \longrightarrow$
 $\text{subnetsOfAdr } (\text{src } p) \cap (\text{fst-set } (\text{sdnets } x)) \neq \{\}$
 ⟨proof⟩

lemma *dest-in-sdnets*[rule-format]: $\neg \text{member } \text{DenyAll } x \longrightarrow p \in \text{dom } (Cp \ x) \longrightarrow$
 $\text{subnetsOfAdr } (\text{dest } p) \cap (\text{snd-set } (\text{sdnets } x)) \neq \{\}$
 ⟨proof⟩

lemma *sdnets-in-subnets*[rule-format]: $p \in \text{dom } (Cp \ x) \longrightarrow \neg \text{member } \text{DenyAll } x \longrightarrow$
 $(\exists (a,b) \in \text{sdnets } x. a \in \text{subnetsOfAdr } (\text{src } p) \wedge b \in \text{subnetsOfAdr } (\text{dest } p))$
 ⟨proof⟩

lemma *disjSD-no-p-in-both*[rule-format]:
 $\llbracket \text{disjSD-2 } x \ y; \neg \text{member } \text{DenyAll } x; \neg \text{member } \text{DenyAll } y;$
 $p \in \text{dom } (Cp \ x); p \in \text{dom } (Cp \ y) \rrbracket \Longrightarrow \text{False}$
 ⟨proof⟩

lemma *list2FWpolicy-eq*: $zs \neq [] \Longrightarrow$
 $Cp \text{ (list2FWpolicy } (x \oplus y \# z)) \ p = Cp \text{ (list2FWpolicy } (y \# z)) \ p$
 ⟨proof⟩

lemma *dom-sep*[rule-format]: $x \in \text{dom } (Cp \text{ (list2FWpolicy } p)) \longrightarrow$
 $x \in \text{dom } (Cp \text{ (list2FWpolicy } (\text{separate } p)))$
 ⟨proof⟩

lemma *domdConcStart*[rule-format]: $x \in \text{dom } (Cp \text{ (list2FWpolicy } (a\#b))) \longrightarrow$
 $x \notin \text{dom } (Cp \text{ (list2FWpolicy } b))$
 $\longrightarrow x \in \text{dom } (Cp \ (a))$
 ⟨proof⟩

lemma *sep-dom2-aux*: $\llbracket x \in \text{dom } (Cp \text{ (list2FWpolicy } (a \oplus y \# z))) \rrbracket$
 $\Longrightarrow x \in \text{dom } (Cp \ (a \oplus \text{list2FWpolicy } (y \# z)))$
 ⟨proof⟩

lemma *sep-dom2-aux2*:

$\llbracket (x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (\text{separate } (y \# z)))) \longrightarrow$
 $x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (y \# z)))$;
 $x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (a \# \text{separate } (y \# z)))) \rrbracket$
 $\implies x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (a \oplus y \# z)))$
 $\langle \text{proof} \rangle$

lemma *sep-dom2[rule-format]*:

$x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (\text{separate } p))) \longrightarrow x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (p)))$
 $\langle \text{proof} \rangle$

lemma *sepDom*: $\text{dom } (\text{Cp } (\text{list2FWpolicy } p)) = \text{dom } (\text{Cp } (\text{list2FWpolicy } (\text{separate } p)))$

$\langle \text{proof} \rangle$

lemma *C-eq-s-ext[rule-format]*: $p \neq [] \longrightarrow$

$\text{Cp } (\text{list2FWpolicy } (\text{separate } p)) \ a = \text{Cp } (\text{list2FWpolicy } p) \ a$

$\langle \text{proof} \rangle$

lemma *C-eq-s*: $p \neq [] \implies \text{Cp } (\text{list2FWpolicy } (\text{separate } p)) = \text{Cp } (\text{list2FWpolicy } p)$

$\langle \text{proof} \rangle$

lemma *sortnMTQ*: $p \neq [] \implies \text{qsort } p \ l \neq []$

$\langle \text{proof} \rangle$

lemmas *C-eq-Lemmas-sep* =

C-eq-Lemmas sortnMT sortnMTQ RS2-NMT NMPrd NMPRS3

lemma *C-eq-until-separated*:

$\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{all-in-list } (\text{policy2list } p) \ l;$
 $\text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $\text{Cp } (\text{list2FWpolicy } (\text{separate } (\text{sort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } \text{Cp } (\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \ l))) = \text{Cp } p$
 $\langle \text{proof} \rangle$

lemma *C-eq-until-separatedQ*:

$\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p); \text{all-in-list } (\text{policy2list } p) \ l;$
 $\text{allNetsDistinct } (\text{policy2list } p) \rrbracket \implies$
 $\text{Cp } (\text{list2FWpolicy } (\text{separate } (\text{qsort } (\text{removeShadowRules2 } (\text{remdups } (\text{removeShadowRules3 } \text{Cp } (\text{insertDeny } (\text{removeShadowRules1 } (\text{policy2list } p)))))) \ l))) = \text{Cp } p$
 $\langle \text{proof} \rangle$

lemma *domID[rule-format]*: $p \neq [] \wedge x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } p)) \longrightarrow$

$x \in \text{dom } (\text{Cp } (\text{list2FWpolicy } (\text{insertDenies } p)))$

$\langle \text{proof} \rangle$

lemma *DA-is-deny*:

$x \in \text{dom } (Cp \ (DenyAllFromTo \ a \ b \oplus \ DenyAllFromTo \ b \ a \oplus \ DenyAllFromTo \ a \ b)) \implies$
 $Cp \ (DenyAllFromTo \ a \ b \oplus \ DenyAllFromTo \ b \ a \oplus \ DenyAllFromTo \ a \ b) \ x = \text{Some } (\text{deny } ())$
 $\langle \text{proof} \rangle$

lemma *iDdomAux*[rule-format]:

$p \neq [] \longrightarrow x \notin \text{dom } (Cp \ (\text{list2FWpolicy } p)) \longrightarrow$
 $x \in \text{dom } (Cp \ (\text{list2FWpolicy } (\text{insertDenies } p))) \longrightarrow$
 $Cp \ (\text{list2FWpolicy } (\text{insertDenies } p)) \ x = \text{Some } (\text{deny } ())$
 $\langle \text{proof} \rangle$

lemma *iD-isD*[rule-format]: $p \neq [] \longrightarrow x \notin \text{dom } (Cp \ (\text{list2FWpolicy } p))$

$\longrightarrow Cp \ (DenyAll \oplus \ \text{list2FWpolicy } (\text{insertDenies } p)) \ x = Cp \ DenyAll \ x$
 $\langle \text{proof} \rangle$

lemma *inDomConc*: $\llbracket x \notin \text{dom } (Cp \ a); x \notin \text{dom } (Cp \ (\text{list2FWpolicy } p)) \rrbracket \implies$
 $x \notin \text{dom } (Cp \ (\text{list2FWpolicy } (a \# p)))$

$\langle \text{proof} \rangle$

lemma *domsdisj*[rule-format]: $p \neq [] \longrightarrow (\forall \ x \ s. \ s \in \text{set } p \wedge x \in \text{dom } (Cp \ A) \longrightarrow$
 $x \notin \text{dom } (Cp \ s)) \longrightarrow y \in \text{dom } (Cp \ A) \longrightarrow$
 $y \notin \text{dom } (Cp \ (\text{list2FWpolicy } p))$

$\langle \text{proof} \rangle$

lemma *isSepaux*:

$\llbracket p \neq []; \text{noDenyAll } (a \# p); \text{separated } (a \# p);$
 $x \in \text{dom } (Cp \ (DenyAllFromTo \ (\text{first-}srcNet \ a) \ (\text{first-}destNet \ a) \oplus$
 $DenyAllFromTo \ (\text{first-}destNet \ a) \ (\text{first-}srcNet \ a) \oplus \ a)) \rrbracket \implies$
 $x \notin \text{dom } (Cp \ (\text{list2FWpolicy } p))$
 $\langle \text{proof} \rangle$

lemma *noneMTsep*[rule-format]: $\text{noneMT } Cp \ p \longrightarrow \text{noneMT } Cp \ (\text{separate } p)$

$\langle \text{proof} \rangle$

lemma *dom-id*:

$\llbracket \text{noDenyAll } (a \# p); \text{separated } (a \# p); p \neq []; x \notin \text{dom } (Cp \ (\text{list2FWpolicy } p));$
 $x \in \text{dom } (Cp \ (a)) \rrbracket$
 $\implies x \notin \text{dom } (Cp \ (\text{list2FWpolicy } (\text{insertDenies } p)))$
 $\langle \text{proof} \rangle$

lemma *C-eq-iD-aux2*[rule-format]:

$noDenyAll1\ p \longrightarrow$
 $separated\ p \longrightarrow$
 $p \neq [] \longrightarrow$
 $x \in dom\ (Cp\ (list2FWpolicy\ p)) \longrightarrow$
 $Cp(list2FWpolicy\ (insertDenies\ p))\ x = Cp(list2FWpolicy\ p)\ x$
 $\langle proof \rangle$

lemma *C-eq-iD*: $\llbracket separated\ p; noDenyAll1\ p; wellformed-policy1-strong\ p \rrbracket \implies$
 $Cp\ (list2FWpolicy\ (insertDenies\ p)) = Cp\ (list2FWpolicy\ p)$
 $\langle proof \rangle$

lemma *noDAsortQ[rule-format]*: $noDenyAll1\ p \longrightarrow noDenyAll1\ (qsort\ p\ l)$
 $\langle proof \rangle$

lemma *NetsCollectedSortQ*: $distinct\ p \implies noDenyAll1\ p \implies all-in-list\ p\ l \implies$
 $singleCombinators\ p \implies NetsCollected\ (qsort\ p\ l)$
 $\langle proof \rangle$

lemmas *CLemmas* = *nMTSort nMTSortQ noneMTRS2 noneMTrd*
noDAsort noDAsortQ nDASC wp1-eq wp1ID
SCp2l ANDSep wp1n-RS2
OTNSEp OTNSC noDA1sep wp1-alternativesep wellformed-eq
wellformed1-alternative-sorted

lemmas *C-eqLemmas-id* = *CLemmas NC2Sep NetsCollectedSep*
NetsCollectedSort NetsCollectedSortQ separatedNC

lemma *C-eq-Until-InsertDenies*: $\llbracket DenyAll \in set\ (policy2list\ p); all-in-list$
 $(policy2list\ p)\ l; allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp\ (list2FWpolicy\ ((insertDenies\ (separate\ (sort\ (removeShadowRules2\ (remdups$
 $(removeShadowRules3\ Cp\ (insertDeny\ (removeShadowRules1\ (policy2list\ p))))))\ l)))))) =$
 $Cp\ p$
 $\langle proof \rangle$

lemma *C-eq-Until-InsertDeniesQ*: $\llbracket DenyAll \in set\ (policy2list\ p); all-in-list$
 $(policy2list\ p)\ l; allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp\ (list2FWpolicy\ ((insertDenies\ (separate\ (qsort\ (removeShadowRules2\ (remdups$
 $(removeShadowRules3\ Cp\ (insertDeny\ (removeShadowRules1\ (policy2list\ p))))))\ l)))))) =$
 $Cp\ p$
 $\langle proof \rangle$

lemma *C-eq-RD-aux*[rule-format]: $Cp\ (p)\ x = Cp\ (removeDuplicates\ p)\ x$
 <proof>

lemma *C-eq-RAD-aux*[rule-format]:
 $p \neq [] \longrightarrow Cp\ (list2FWpolicy\ p)\ x = Cp\ (list2FWpolicy\ (removeAllDuplicates\ p))\ x$
 <proof>

lemma *C-eq-RAD*:
 $p \neq [] \implies Cp\ (list2FWpolicy\ p) = Cp\ (list2FWpolicy\ (removeAllDuplicates\ p))$
 <proof>

lemma *C-eq-compile*:
 $\llbracket DenyAll \in set\ (policy2list\ p); all-in-list\ (policy2list\ p)\ l;$
 $allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp\ (list2FWpolicy\ (removeAllDuplicates\ (insertDenies\ (separate\ (sort$
 $(removeShadowRules2\ (remdups\ (removeShadowRules3\ Cp\ (insertDeny$
 $(removeShadowRules1\ (policy2list\ p))))))\ l)))) = Cp\ p$
 <proof>

lemma *C-eq-compileQ*:
 $\llbracket DenyAll \in set\ (policy2list\ p); all-in-list\ (policy2list\ p)\ l;$
 $allNetsDistinct\ (policy2list\ p) \rrbracket \implies$
 $Cp\ (list2FWpolicy\ (removeAllDuplicates\ (insertDenies\ (separate\ (qsort$
 $(removeShadowRules2\ (remdups\ (removeShadowRules3\ Cp\ (insertDeny$
 $(removeShadowRules1\ (policy2list\ p))))))\ l)))) = Cp\ p$
 <proof>

lemma *C-eq-normalizePr*:
 $\llbracket DenyAll \in set\ (policy2list\ p);$
 $allNetsDistinct\ (policy2list\ p);$
 $all-in-list\ (policy2list\ p)\ (Nets-List\ p) \rrbracket \implies$
 $Cp\ (list2FWpolicy\ (normalizePr\ p)) = Cp\ p$
 <proof>

lemma *C-eq-normalizePrQ*:
 $\llbracket DenyAll \in set\ (policy2list\ p);$
 $allNetsDistinct\ (policy2list\ p);$
 $all-in-list\ (policy2list\ p)\ (Nets-List\ p) \rrbracket \implies$
 $Cp\ (list2FWpolicy\ (normalizePrQ\ p)) = Cp\ p$
 <proof>

lemma *domSubset3*: $\text{dom } (Cp \ (DenyAll \oplus x)) = \text{dom } (Cp \ (DenyAll))$
 ⟨proof⟩

lemma *domSubset4*: $\text{dom } (Cp \ (DenyAllFromTo \ x \ y \oplus DenyAllFromTo \ y \ x \oplus AllowPortFromTo \ x \ y \ dn)) =$
 $\text{dom } (Cp \ (DenyAllFromTo \ x \ y \oplus DenyAllFromTo \ y \ x))$
 ⟨proof⟩

lemma *domSubset5*:
 $\text{dom } (Cp \ (DenyAllFromTo \ x \ y \oplus DenyAllFromTo \ y \ x \oplus AllowPortFromTo \ y \ x \ dn)) =$
 $\text{dom } (Cp \ (DenyAllFromTo \ x \ y \oplus DenyAllFromTo \ y \ x))$
 ⟨proof⟩

lemma *domSubset1*: $\text{dom } (Cp \ (DenyAllFromTo \ one \ two \oplus DenyAllFromTo \ two \ one \oplus AllowPortFromTo \ one \ two \ dn \oplus x)) =$
 $\text{dom } (Cp \ (DenyAllFromTo \ one \ two \oplus DenyAllFromTo \ two \ one \oplus x))$
 ⟨proof⟩

lemma *domSubset2*:
 $\text{dom } (Cp \ (DenyAllFromTo \ one \ two \oplus DenyAllFromTo \ two \ one \oplus AllowPortFromTo \ two \ one \ dn \oplus x)) =$
 $\text{dom } (Cp \ (DenyAllFromTo \ one \ two \oplus DenyAllFromTo \ two \ one \oplus x))$
 ⟨proof⟩

lemma *ConcAssoc2*: $Cp \ (X \oplus Y \oplus ((A \oplus B) \oplus D)) = Cp \ (X \oplus Y \oplus A \oplus B \oplus D)$
 ⟨proof⟩

lemma *ConcAssoc3*: $Cp \ (X \oplus ((Y \oplus A) \oplus D)) = Cp \ (X \oplus Y \oplus A \oplus D)$
 ⟨proof⟩

lemma *RS3-NMT[rule-format]*: $DenyAll \in \text{set } p \longrightarrow$
 $\text{removeShadowRules3 } Cp \ p \neq []$
 ⟨proof⟩

lemma *norm-notMT*: $DenyAll \in \text{set } (\text{policy2list } p) \implies \text{normalizePr } p \neq []$
 ⟨proof⟩

lemma *norm-notMTQ*: $\text{DenyAll} \in \text{set } (\text{policy2list } p) \implies \text{normalizePrQ } p \neq []$
 <proof>

lemma *domDA*: $\text{dom } (\text{Cp } (\text{DenyAll} \oplus A)) = \text{dom } (\text{Cp } (\text{DenyAll}))$
 <proof>

lemmas *domain-reasoningPr* = *domDA ConcAssoc2 domSubset1 domSubset2 domSubset3 domSubset4 domSubset5 domSubsetDistr1 domSubsetDistr2 domSubsetDistrA domSubsetDistrD coerc-assoc ConcAssoc ConcAssoc3*

The following lemmas help with the normalisation

lemma *list2policyR-Start*[rule-format]: $p \in \text{dom } (\text{Cp } a) \longrightarrow$
 $\text{Cp } (\text{list2policyR } (a \# \text{list})) \text{ } p = \text{Cp } a \text{ } p$
 <proof>

lemma *list2policyR-End*: $p \notin \text{dom } (\text{Cp } a) \implies$
 $\text{Cp } (\text{list2policyR } (a \# \text{list})) \text{ } p = (\text{Cp } a \oplus \text{list2policy } (\text{map } \text{Cp } \text{list})) \text{ } p$
 <proof>

lemma *l2polR-eq-el*[rule-format]: $N \neq [] \longrightarrow$
 $\text{Cp } (\text{list2policyR } N) \text{ } p = (\text{list2policy } (\text{map } \text{Cp } N)) \text{ } p$
 <proof>

lemma *l2polR-eq*: $N \neq [] \implies$
 $\text{Cp } (\text{list2policyR } N) = (\text{list2policy } (\text{map } \text{Cp } N))$
 <proof>

lemma *list2FWpolicys-eq-el*[rule-format]:
 $\text{Filter} \neq [] \longrightarrow \text{Cp } (\text{list2policyR } \text{Filter}) \text{ } p = \text{Cp } (\text{list2FWpolicy } (\text{rev } \text{Filter})) \text{ } p$
 <proof>

lemma *list2FWpolicys-eq*:
 $\text{Filter} \neq [] \implies$
 $\text{Cp } (\text{list2policyR } \text{Filter}) = \text{Cp } (\text{list2FWpolicy } (\text{rev } \text{Filter}))$
 <proof>

lemma *list2FWpolicys-eq-sym*:
 $\text{Filter} \neq [] \implies$

$Cp (list2policyR (rev Filter)) = Cp (list2FWpolicy Filter)$
 $\langle proof \rangle$

lemma $p\text{-eq}[rule\text{-format}]$: $p \neq [] \longrightarrow$
 $list2policy (map Cp (rev p)) = Cp (list2FWpolicy p)$
 $\langle proof \rangle$

lemma $p\text{-eq2}[rule\text{-format}]$: $normalizePr x \neq [] \longrightarrow$
 $Cp (list2FWpolicy (normalizePr x)) = Cp x \longrightarrow$
 $list2policy (map Cp (rev (normalizePr x))) = Cp x$
 $\langle proof \rangle$

lemma $p\text{-eq2Q}[rule\text{-format}]$: $normalizePrQ x \neq [] \longrightarrow$
 $Cp (list2FWpolicy (normalizePrQ x)) = Cp x \longrightarrow$
 $list2policy (map Cp (rev (normalizePrQ x))) = Cp x$
 $\langle proof \rangle$

lemma $list2listNMT[rule\text{-format}]$: $x \neq [] \longrightarrow map\ sem\ x \neq []$
 $\langle proof \rangle$

lemma $Norm\text{-}Distr2$:
 $r\ o\text{-}f\ ((P \otimes_2 (list2policy Q))\ o\ d) =$
 $(list2policy ((P \otimes_L Q)\ (op \otimes_2)\ r\ d))$
 $\langle proof \rangle$

lemma $NATDistr$: $\llbracket N \neq [] ; F = Cp (list2policyR N) \rrbracket \implies$
 $((\lambda (x,y). x)\ o\text{-}f\ ((NAT \otimes_2 F)\ o\ (\lambda x. (x,x)))) =$
 $(list2policy\ (\ (\ (NAT \otimes_L (map\ Cp\ N))\ (op \otimes_2)\$
 $(\lambda (x,y). x)\ (\lambda x. (x,x))))))$
 $\langle proof \rangle$

lemma $C\text{-eq-normalize-manual}$:
 $\llbracket DenyAll \in set\ (policy2list\ p);$
 $allNetsDistinct\ (policy2list\ p);$
 $all\text{-in-list}\ (policy2list\ p)\ l \rrbracket \implies$
 $Cp (list2FWpolicy (normalize-manual-orderPr p l)) = Cp p$
 $\langle proof \rangle$

lemma *p-eq2-manualQ[rule-format]: normalize-manual-orderPrQ x l ≠ [] →*
Cp (list2FWpolicy (normalize-manual-orderPrQ x l)) = Cp x →
list2policy (map Cp (rev (normalize-manual-orderPrQ x l))) = Cp x
 ⟨proof⟩

lemma *norm-notMT-manualQ: DenyAll ∈ set (policy2list p) ⇒ normalize-manual-orderPrQ*
p l ≠ []
 ⟨proof⟩

lemma *C-eq-normalizePr-manualQ:*
 $\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } p);$
 $\text{allNetsDistinct } (\text{policy2list } p);$
 $\text{all-in-list } (\text{policy2list } p) \text{ l} \rrbracket \Rightarrow$
 $Cp (\text{list2FWpolicy } (\text{normalize-manual-orderPrQ } p \text{ l})) = Cp p$
 ⟨proof⟩

lemma *p-eq2-manual[rule-format]: normalize-manual-orderPr x l ≠ [] →*
Cp (list2FWpolicy (normalize-manual-orderPr x l)) = Cp x →
list2policy (map Cp (rev (normalize-manual-orderPr x l))) = Cp x
 ⟨proof⟩

lemma *norm-notMT-manual: DenyAll ∈ set (policy2list p) ⇒ normalize-manual-orderPr p l*
≠ []
 ⟨proof⟩

As an example, how this theorems can be used for a concrete normalisation instantiation.

lemma *normalizePrNAT: $\llbracket \text{DenyAll} \in \text{set } (\text{policy2list } \text{Filter});$*
allNetsDistinct (policy2list Filter);
all-in-list (policy2list Filter) (Nets-List Filter) $\rrbracket \Rightarrow$
 $((\lambda (x,y). x) \circ f (((\text{NAT} \otimes_2 Cp \text{Filter}) \circ (\lambda x. (x,x)))))) =$
 $\text{list2policy } ((\text{NAT} \otimes_L (\text{map } Cp (\text{rev } (\text{normalizePr } \text{Filter})))) (\text{op } \otimes_2) (\lambda (x,y). x) (\lambda x.$
 $(x,x)))$
 ⟨proof⟩

lemma *dom.Simpl[simp]: dom (Cp (A ⊕ DenyAll)) = dom (Cp (DenyAll))*
 ⟨proof⟩

end

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